Formal Verification of Boolean Unification Algorithms with Coq

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A Major Qualifying Project
Submitted to the Faculty of
Worcester Polytechnic Institute
in partial fulfillment of the requirements for the
Degree in Bachelor of Science in
Computer Science

WPI

Computer Science
Worcester Polytechnic Institute
Advised by Professor Daniel Dougherty
April 25, 2019
Abstract

We report on a verified implementation of two (well-known) algorithms for unification modulo the theory of Boolean rings: Lowenheim’s method and the method of Successive Variable Elimination. The implementations and proofs of correctness were done in the Coq proof assistant; we view this contribution as an early step in a larger project of developing a suite of verified implementations of equational unification algorithms.
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Chapter 1

Library B_Unification.introduction

1.1 Introduction

In the field of computer science, one problem of significance is that of equational unification; namely, the finding of solutions to a given set of equations with respect to a set of equational axioms. While there are several variants of equational unification, for the purposes of this paper we are going to limit our scope to that of Boolean unification, which deals with the finding of unifiers for the equations defining Boolean rings. There exists a great deal of research in the formal verification of unification algorithms [Baader and Snyder, 2001]; our research focused on two of these algorithms: Lowenheim’s formula and successive variable elimination. To conduct our research, we utilized the Coq proof assistant https://coq.inria.fr/ to create formal specifications of both of these algorithms’ behaviors in addition to proving their correctness. While proofs for both of these algorithms already exist [Baader and Nipkow, 1998, p. 254-258], prior to the writing of this paper, no formal treatment using a proof assistant such as Coq had been undertaken, so it is hoped that our efforts towards porting these algorithms onto software provide a useful suite of tools for anyone interested in working with equational logic.

Due to the differences in the innate nature of Lowenheim’s formula compared to that of successive variable elimination, our project was divided into two separate developments, each approaching their respective goals from a different direction. The primary distinction between these two treatments comes down to their representations of equations. The Lowenheim’s formula development uses a more straightforward, term-based representation of equations while the successive variable elimination development opts to represent equations in their polynomial forms. Fortunately, due to the fact that every term has a unique polynomial representation [Baader and Nipkow, 1998, p. 263], these two formats for representing equations are mathematically equivalent to one another.
1.2 Formal Verification

Formal verification is the term used to describe the act of verifying (or disproving) the correctness of software and hardware systems or theories. Formal verification consists of a set of techniques that perform static analysis on the behavior of a system, or the correctness of a theory. It differs from dynamic analysis that uses simulation to evaluate the correctness of a system.

More simply stated, formal verification is the process of examining whether a system or a theory “does what it is supposed to do.” If it is a system, then scientists formally verify that it satisfies its design requirements. Formal verification is also different from testing. Software testing tries to detect “bugs”, specific errors, and requirements in the system, whereas verification acts as a general safeguard that the system is always error-free. As Edsger Dijkstra stated [Franco, 2018, slide 7], testing can be used to show the presence of bugs, but never to show their absence. When trying to verify a theory, scientists formally verify the correctness of the theory by formulating its proof using a formal language, axioms and inference rules.

Formal verification is used because it does not have to evaluate every possible case or state to determine if a system or theory meets all the preset logical conditions and requirements. Moreover, as design and software systems sizes have increased (along with their simulation times), verification teams have been looking for alternative methods of proving or disproving the correctness of a system in order to reduce the required time to perform a correctness check or evaluation.

1.2.1 Proof Assistants

A proof assistant is a software tool that is used to formulate and prove or disprove theorems in computer science or mathematical logic. They are also called interactive theorem provers and they may also involve some type of proof and text editor that the user can use to form, prove, and define theorems, lemmas, functions, etc. They facilitate that process by allowing the user to search definitions, terms and even provide some kind of guidance during the formulation or proof of a theorem. Some examples of proof assistants are Coq - which is the one we are using -, Isabelle, HOL Light and Lean.

1.2.2 Verifying Systems

Formal verification is used to verify the correctness of software or hardware systems [Harrison, 2002]. When used to verify systems, formal verification can be thought as a mathematical proof of the correctness of a design with respect to a formal specification. The actual system is represented by a formal model and then the formal verification happens on the model, based on the required specifications of the system. Unlike testing, formal verification is exhaustive. However, it is difficult to make for real-world systems, time consuming and only as reliable as the actual model.
1.2.3 Verifying Theories

Formal verification is also used in to prove theorems. These theorems could be related to a computing system or just to abstract mathematical theorems. Mathematical theorems that have been proven using a proof assistant include the Four-Color theorem and the Feit-Thompson theorem. As in proving systems, when proving theorems one also needs a formal logic to formulate the theorem and prove it. A formal logic consists of a formal language to express the theorems, a collection of formulas called axioms and inference rules to derive new axioms based on existing ones. A theorem to be proven could be in a logical form, like DeMorgan’s Law or it could in another mathematical area; in trigonometry for example, it could be useful to prove that $\sin(x + y) = \sin(x) \cdot \cos(y) + \cos(x) \cdot \sin(y)$, formally, because that proof could be used as building block in a more complex system. Sometimes proving the correctness of a real world systems boils down to verifying mathematical proofs like the previous one, so the two approaches are often linked together.

1.3 Unification

Before defining unification, there is some terminology to understand.

1.3.1 Terms and Substitutions

Definition 1.3.1 A term is either a variable or a function applied to terms [Baader and Nipkow, 1998, p. 34].

By this definition, a constant term is just a nullary function.

Definition 1.3.2 A variable is a symbol capable of taking on the value of any term.

An example of a term is $f(a, x)$, where $f$ is a function of two arguments, $a$ is a constant, and $x$ is a variable.

Definition 1.3.3 A term is ground if no variables occur in it [Baader and Nipkow, 1998, p. 37].

The last example is not a ground term but $f(a, a)$ would be.

Definition 1.3.4 A substitution is a mapping from variables to terms.

Definition 1.3.5 The domain of a substitution is the set of variables that do not get mapped to themselves.

Definition 1.3.6 The range is the set of terms that are mapped to by the domain [Baader and Nipkow, 1998, p. 37].

It is common for substitutions to be referred to as mappings from terms to terms. A substitution $\sigma$ can be extended to this form by defining $\hat{\sigma}(s)$ for two cases of $s$. If $s$ is a variable, then $\hat{\sigma}(s) := \sigma(s)$. If $s$ is a function $f(s_1, ..., s_n)$, then $\hat{\sigma}(s) := f(\hat{\sigma}(s_1), ..., \hat{\sigma}(s_n))$ [Baader and Nipkow, 1998, p. 38].
1.3.2 Unification and Unifiers

Unification is the process of solving a set of equations between two terms.

**Definition 1.3.7** The set of equations to solve is referred to as a **unification problem** [Baader and Nipkow, 1998, p. 71].

The process of solving one of these problems can be classified by the set of terms considered and the equality of any two terms. The latter property is what distinguishes two broad groups of algorithms, namely syntactic and semantic unification.

**Definition 1.3.8** If two terms are only considered equal if they are identical, then the unification is **syntactic** [Baader and Nipkow, 1998, p. 71].

**Definition 1.3.9** If two terms are equal with respect to an equational theory \([E]\), then the unification is **semantic**. It is also called \([E]-\)unification [Baader and Nipkow, 1998, p. 224].

For example, the terms \(x \ast y\) and \(y \ast x\) are not syntactically equal, but they are semantically equal modulo commutativity of multiplication.

The goal of unification is to find the best solution to a problem, which formally means to produce a most general unifier of the problem. The next four definitions should make this clearer.

**Definition 1.3.10** A substitution \(\sigma\) **unifies** an equation \(s \equiv t\) if applying \(\sigma\) to both sides makes them equal \(\sigma(s) = \sigma(t)\).

**Definition 1.3.11** If \(\sigma\) unifies every equation in the problem \(S\), we call \(\sigma\) a **solution** or **unifier** of \(S\) [Baader and Nipkow, 1998, p. 71].

**Definition 1.3.12** A substitution \(\sigma\) is **more general** than \(\sigma'\) if there exists a third substitution \(\delta\) such that \(\sigma'(u) = \delta(\sigma(u))\) for any term \(u\).

**Definition 1.3.13** A substitution is a **most general unifier** or **mgu** of a problem if it is more general than every other solution to the problem [Baader and Nipkow, 1998, p. 71].

It should be noted that although solvable problem of Boolean unification produce a single mgu, semantic unification problems in general can have zero, multiple, or infinitely many mgu’s [Baader and Nipkow, 1998, p. 226].

1.3.3 Syntactic Unification

This is the simplest version of unification. It is a special case of \(E\)-unification where \(E = \emptyset\). For two terms to be considered equal they must be identical. Problems of this kind can be solved by repeated transformations until the solution pops out similar to solving a linear system by Guassian elimination [Baader and Nipkow, 1998, p. 73]. One of the most notable applications of syntactic unification is the Hindley-Milner type system used in functional programming languages like ML [Damas and Milner, 1982]. More complicated type systems such as the one used by Coq require more complicated versions of unification (e.g. higher-order unification) [Chlipala, 2010].
1.3.4 Semantic Unification

This kind of unification involves an equational theory. Given a set of identities $E$, we write that two terms $s$ and $t$ are equal with regards to $E$ as $s \approx_E t$. This means that there is a chain of terms leading from $s$ to $t$ in which each term is derived from the previous one by replacing a subterm $u$ by a term $v$ when $u = v$ is an instance of an axiom of $E$. For a careful definition see [Baader and Nipkow, 1998], but an example should make the idea clear.

If we take $C$ to be the set $\{f(x, y) \approx f(y, x), f(b, f(a, c)) \approx_C f(f(c, a), b)\}$, we then have $f(b, f(a, c)) \approx_C f(f(c, a), b)$ via the sequence of steps $f(b, f(a, c)) \approx_C f(f(c, a), b)$. Now we say that two terms $s$ and $t$ are $E$-unifiable if there is a substitution $\sigma$ such that $\sigma(s) \approx_E \sigma(t)$. For example, the problem $\{f(x, f(a, y)) \approx f(f(c, a), b)\}$ is $C$-unified by the substitution $\{x \mapsto b, y \mapsto c\}$ since $f(b, f(a, c)) \approx_C f(f(c, a), b)$. For some $E$, the problem of $E$-unification can actually be undecidable [Baader and Nipkow, 1998, p. 71]. An example would be unification modulo ring theory.

1.3.5 Boolean Unification

In this paper, we focus on unification modulo Boolean ring theory, also referred to as $B$-unification. The allowed terms in this theory are the constants 0 and 1 and binary functions $+$ and $\ast$. The set of identities $B$ is defined as follows:

$$\begin{align*}
&x + y \approx y + x, \quad x \ast y \approx y \ast x, \\
&(x + y) + z \approx x + (y + z), \quad (x \ast y) \ast z \approx x \ast (y \ast z), \\
&x + x \approx 0, \quad x \ast x \approx x, \\
&0 + x \approx x, \quad 0 \ast x \approx 0, \\
&x \ast (y + z) \approx (x \ast y) + (x \ast z), \quad 1 \ast x \approx x
\end{align*}$$

[Baader and Nipkow, 1998, p. 250]. This set is equivalent to the axioms of ring theory with the addition of $x + x \approx_B 0$ and $x \ast x \approx_B x$.

Although a unification problem was defined as a set of equations between two terms, problems of Boolean unification can be viewed as just a single equation $t \approx_B 0$. To see this, first note that for any terms $u$ and $v$ we have

$$u \approx_B v \quad \text{if and only if} \quad u + v \approx_B 0.$$ 

We also have that for any term $w$

$$w \approx_B 0 \quad \text{if and only if} \quad w + 1 \approx_B 1.$$ 

It follows that for any set of terms $u_1, \ldots, u_n, v_1, \ldots, v_n$,

$$u_1 \approx_B v_1, \ldots, u_n \approx_B v_n$$
all hold if and only if the equations

\[ u_1 + v_1 + 1 \approx_B 1, \ldots, u_n + v_n + 1 \approx_B 1 \]

all hold, and this if and only if the single equation

\[ (u_1 + v_1 + 1) \ast \ldots \ast (u_n + v_n + 1) \approx_B 1 \]

holds, or in other words

\[ (u_1 + v_1 + 1) \ast \ldots \ast (u_n + v_n + 1) + 1 \approx_B 0 \]

holds. Thus a problem

\[ s_1 \approx_B t_1, \ldots, s_n \approx_B t_n \]

is solvable by the same substitutions as the problem

\[ (s_1 + t_1 + 1) \ast \ldots \ast (s_n + t_n + 1) + 1 \approx_B 0. \]

This fact allows both developments to use the simpler \( t \approx_B 0 \) description of a problem.

### 1.4 Importance

Given that the emergence of proof assistance software is still in its infancy relative to the traditional methods of theorem proving, it would be a disservice for us to not establish the importance of this technology and its implications for the future of mathematics. Unlike in years past, where the sheer volume of detail could derail the developments of sound theorems, proof assistants now guarantee through their properties of verification that any development verified by them is free from lapses in logic on account of the natural failings of the human mind. Additionally, due to the adoption of a well-defined shared language, many of the ambiguities naturally present in the exchange of mathematical ideas between colleagues are mitigated, leading to a smoother learning curve for newcomers trying to understand the nuts and bolts of a complex theorem. The end result of these phenomena is a faster iterative development cycle for mathematicians as they now can spend more time on proving things and building off of the work of others since they no longer need to devote as much of their efforts towards verifying the correctness of the theorems they are operating across.

Bearing this in mind, it should come as no surprise that there is a utility in going back to older proofs that have never been verified by a proof assistant and redeveloping them for the purposes of ensuring their correctness. If the theorem is truly sound, it stands to reason that any additional rigorous scrutiny would only serve to bolster the credibility of its claims, and conversely, if the theorem is not sound, it is a benefit to the academic community at large to be made aware of its shortcomings. Therefore, for these reasons we set out to formally verify two algorithms across Boolean Unification.
1.5 Development

1.5.1 Algorithms

There are many different approaches that one could take to go about formalizing a proof of Boolean Unification algorithms, each with their own challenges. For this development, we have opted to base our work on chapter 10, *Equational Unification*, in *Term Rewriting and All That* by Franz Baader and Tobias Nipkow [Baader and Nipkow, 1998]. Specifically, section 10.4, titled *Boolean Unification*, details Boolean rings, data structures to represent them, and two algorithms to perform unification in Boolean rings.

We chose to implement these two different Boolean Unification algorithms, and then proceeded to formally prove their correctness on all inputs. The two algorithms in question are Lowenhein’s formula and Successive Variable Elimination.

The first solution, **Lowenheim’s algorithm**, is based on the idea that the Lowenheim formula can take any unifier of a Boolean unification problem and turn it into a most general unifier. The algorithm then of course first requires a unifier to begin; we have opted to use a simple brute force solution to find a ground unifier, replacing variables with only 0 or 1. This ground solution is then passed through the formula, to create a most general unifier. Lowenheim’s algorithm is implemented in the file *lowenheim.v*, and the proof of correctness is in *lowenheim_proof.v*.

The second algorithm, **successive variable elimination**, is built on the idea that by factoring variables out of an equation one-by-one, we can eventually reach a problem that can be solved by the identity unifier. This base problem is then slowly built up by adding the variables that were previously eliminated, building up the matching unifier as we do so. Once we have added all variables back in, we are left with the original problem as well as a most general unifier for it. Successive variable elimination and its proof of correctness are both in the file *sve.v*.

1.5.2 Data Structures

The data structure used to represent a Boolean unification problem completely changes the shape of both the unification algorithm and the proof of correctness, and is therefore a very important decision. For this development, we have selected two different representations of Boolean rings: first as a “Term” inductive type, and then as lists of lists representing terms in polynomial form.

**Term Inductive Type**

The Term inductive type, used in the proof of Lowenheim’s algorithm, is very simple and rather intuitive – a term in a Boolean ring is one of 5 things:

- The number 0
- The number 1
• A variable
• Two terms added together
• Two terms multiplied together

In our development, variables are represented as natural numbers.

After defining terms like this, it is necessary to define a new equality relation, referred to as term equivalence, for comparing terms. With the term equivalence relation defined, it is easy to define ten axioms enabling the ten identities that hold true over terms in Boolean rings.

The inductive representation of terms in a Boolean ring and unification over these terms are defined in the file terms.v.

Benefits and Challenges of the Inductive Type

The most apparent benefit of utilizing an inductive representation of terms becomes obvious from the moment one looks at a term in this format: inductively represented terms are easily able to be read and understood since the format is identical to the typical presentation of equations one is used to. This allows for inductively represented terms to be very intuitive and easy to reason about. This benefit does not come without its costs however. For starters, by representing terms in this manner, we can no longer make use of Coq’s built-in equivalence operator since it would be corrupted by the axioms of Boolean rings and lead to bogus proofs. This forced us to develop our own equivalence relation that strictly abides by the Boolean ring axioms. While this certainly prevented Coq from accepting erroneous proofs, it did significantly increase the tediousness and complexity of proving theorems on account of the fact that Coq could not perform induction across our custom equivalence relation. At best, this resulted in proofs that were substantially longer than they would have been otherwise with a more powerful definition (such as Coq’s built in equivalence relation), and at worst resulted in certain lemmas being unprovable, forcing them to be axiomatized.

Polynomial List-of-List Representation

The second representation, used in the proof of successive variable elimination, uses lists of lists of variables to represent terms in polynomial form. A monomial is a list of distinct variables multiplied together. A polynomial, then, is a list of distinct monomials added together. Variables are represented the same way, as natural numbers. The terms 0 and 1 are represented as the empty polynomial and the polynomial containing only the empty monomial, respectively.

The interesting part of the polynomial representation is how the ten identities are implemented. Rather than writing axioms enabling these transformations, we chose to implement the addition and multiplication operations in such a way to ensure these rules hold true, as described in Term Rewriting [Baader and Nipkow, 1998].
Addition is performed by cancelling out all repeated occurrences of monomials in the result of appending the two lists together (i.e., \(x + x = 0\)). This is equivalent to the symmetric difference in set theory, keeping only the terms that are in either one list or the other (but not both). Multiplication is slightly more complicated. The product of two polynomials is the result of multiplying all combinations of monomials in the two polynomials and removing all repeated monomials. The product of two monomials is the result of keeping only one copy of each repeated variable after appending the two together.

To assist with maintaining the strict polynomial form, a “repair” function was defined. This function, given any list of lists of variables, will sort and remove duplicates to ensure the result is a proper polynomial. As a result of this design, we are able to compare monomials and polynomials using the standard Coq equivalence relation for lists, rather than defining our own. In this way, we have effectively embedded the ten axioms in our operations, and do not need to manually declare them.

The polynomial representation is defined in the file `poly.v`. Unification over these polynomials is defined in `poly_unif.v`.

**Benefits and Challenges of the List Representation**

As mentioned above, one of the main benefits of the list representation is that is enables us to use the standard Coq equivalence operator in comparing terms. This makes a wide variety of things easier, from removing the need to prove compatibility of functions with equivalence for rewriting, to allowing us to use all of the standard library lemmas relating to lists. It does, however, come at a cost.

The biggest issue with this design is the amount of work that goes into maintaining this form at every term. Our addition function is defined very simply; we just append the two polynomials, and call our “repair” function on the result. While this sounds simple, it becomes incredibly difficult to prove facts about addition (and our other operations) because of the repair function.

This function does three things: sort the list, cancel out duplicates, and convert all sublists to properly formatted monomials. The main difficulties come from the first two parts. Sorting is incredibly difficult to deal with, as it makes induction over these lists infinitely harder. When proving some fact with induction, the goal of the proof is often something of the form

\[ f(a :: l) = f(a) :: f(l). \]

However, if the function in question sorts the list it’s given, there is no guarantee that \(a\) is going to be the head of the resulting list, thus making the result unprovable. As a result, we had to prove many lemmas about Permutations, and almost exclusively compare lists as a permutation of one another when working with polynomial operations.

Another challenge comes from the cancelling of duplicates. When working with more in-depth proofs of polynomial arithmetic, we often try to prove that some element \(x\) either will or won’t be in a polynomial after some \(f\) is applied, based on whether or not it is in the polynomial before. This leads us to a point where we need to reason about if \(x\) should
be eliminated from either list, which requires us to know how many times $x$ appears in each list. However, even if we know whether or not $x$ should be removed from the original list, it is hard to reason about if it should be removed from the list after $f$ is applied, as $f$ is not one-to-one and there may be some $y$ such that $f\ x = f\ y$. This once again complicates proofs a lot, and required us to prove many facts about our `nodup_cancel` function performing this de-duplication.

After working through these hiccups, though, some aspects of the project became incredibly simple. As mentioned above, the math operations were both very easy to define, and the act of variable elimination and adding itself is very straightforward when you can simply filter a polynomial with the Coq list functions. Given the chance, it probably would have been beneficial to look into defining our own equivalence relation that compares without order, removing the need for sorting. The issue of deduplication would have still come up in one form or another, though, so we probably could not have easily avoided the problems caused by that.
Chapter 2

Library B_Unification.terms

Require Import Bool.
Require Import Omega.
Require Import EqNat.
Require Import List.
Require Import Setoid.
Import ListNotations.

2.1 Introduction

In order for any proofs to be constructed in Coq, we need to formally define the logic and data across which said proofs will operate. Since the heart of our analysis is concerned with the unification of Boolean equations, it stands to reason that we should articulate precisely how algebra functions with respect to Boolean rings. To attain this, we shall formalize what an equation looks like, how it can be composed inductively, and also how substitutions behave when applied to equations.

2.2 Terms

2.2.1 Definitions

We shall now begin describing the rules of Boolean arithmetic as well as the nature of Boolean equations.

Define a variable to be a natural number

Definition var := nat.

A term, as has already been previously described, is now inductively declared to hold either a constant value, a single variable, a sum of terms, or a product of terms.

Inductive term: Type :=
  | T0 : term
For convenience’s sake, we define some shorthanded notation for readability.

Implicit Types $x, y, z : \text{term}$.
Implicit Types $n, m : \text{var}$.

Notation "$x + y$" := (SUM $x$ $y$) (at level 50, left associativity).
Notation "$x \times y$" := (PRODUCT $x$ $y$) (at level 40, left associativity).

### 2.2.2 Axioms

Now that we have informed Coq on the nature of what a term is, it is now time to propose a set of axioms that will articulate exactly how algebra behaves across Boolean rings. This is a requirement since the very act of unifying an equation is intimately related to solving it algebraically. Each of the axioms proposed below describe the rules of Boolean algebra precisely and in an unambiguous manner. None of these should come as a surprise to the reader; however, if one is not familiar with this form of logic, the rules regarding the summation and multiplication of identical terms might pose as a source of confusion.

For reasons of keeping Coq’s internal logic consistent, we roll our own custom equivalence relation as opposed to simply using "=". This will provide a surefire way to avoid any odd errors from later cropping up in our proofs. Of course, by doing this we introduce some implications that we will need to address later.

Parameter $\text{eqv} : \text{term} \to \text{term} \to \text{Prop}$.

Here we introduce some special notation for term equivalence

Infix "$==$" := $\text{eqv}$ (at level 70).

Below is the set of fundamental axioms concerning the equivalence "$==$" relation. They form the boolean ring (or system) on which Lowenheim’s formula and proof are developed.

Most of these axioms will appear familiar to anyone; however, certain ones such as the summation of two identical terms are true only across Boolean rings and as such might appear strange at first glance.

Axiom $\text{sum\_comm} : \forall x, y, x + y == y + x$.
Axiom $\text{sum\_assoc} : \forall x, y, z, (x + y) + z == x + (y + z)$.
Axiom $\text{sum\_id} : \forall x, T0 + x == x$.

Across boolean rings, the summation of two terms will always be 0 because there are only two elements in the ring: 0 and 1. For this reason, the mapping of $1 + 1$ has nowhere else to go besides 0.

Axiom $\text{sum\_x\_x} : \forall x, x + x == T0$.
Axiom $\text{mul\_comm} : \forall x, y, x \times y == y \times x$. 
Axiom \textit{mul_assoc} : \forall x y z, (x \times y) \times z == x \times (y \times z).

Across boolean rings, the multiplication of two identical terms will always be the same as just having one instance of said term. This is because \(0 \times 0 = 0\) and \(1 \times 1 = 1\) as one would expect normally.

Axiom \textit{mul_x_x} : \forall x, x \times x == x.
Axiom \textit{mul_T0_x} : \forall x, T0 \times x == T0.
Axiom \textit{mul_id} : \forall x, T1 \times x == x.

Axiom \textit{distr} : \forall x y z, x \times (y + z) == (x \times y) + (x \times z).

Any axioms beyond this point of the development are not considered part of the “fundamental axiom system”, but they still need to exist for the development and proofs to hold.

Across all equations, adding an expression to both sides does not break the equivalence of the relation.

Axiom \textit{term_sum_symmetric} :
\forall x y z, x == y \leftrightarrow x + z == y + z.

Axiom \textit{refl_comm} :
\forall t1 t2, t1 == t2 \rightarrow t2 == t1.

Axiom \textit{T1_not_equiv_T0} :
\neg(T1 == T0).

Hint Resolve \textit{sum_comm sum_assoc sum_id distr mul_comm mul_assoc mul_x_x mul_T0_x mul_id}.

Now that the core axioms have been taken care of, we need to handle the implications posed by our custom equivalence relation. Below we inform Coq of the behavior of our equivalence relation with respect to reflexivity, symmetry, and transitivity in order to allow for rewrites during the construction of proofs operating across our new equivalence relation.

Axiom \textit{eqv_ref} : \textbf{Reflexive eqv}.
Axiom \textit{eqv_sym} : \textbf{Symmetric eqv}.
Axiom \textit{eqv_trans} : \textbf{Transitive eqv}.

Add Parametric Relation : \textit{term eqv}
- \textit{reflexivity proved by @eqv_ref}
- \textit{symmetry proved by @eqv_sym}
- \textit{transitivity proved by @eqv_trans}
as \textit{eq_set_rel}.

Axiom \textit{SUM_compat} :
\forall x x', x == x' \rightarrow
\forall y y', y == y' \rightarrow
(x + y) == (x' + y').

Axiom \textit{PRODUCT_compat} :
∀ x, x \neq x' \rightarrow
∀ y, y \neq y' \rightarrow
(x \times y) \neq (x' \times y').

Add Parametric Morphism : SUM with
signature eqv \Rightarrow eqv \Rightarrow eqv as SUM_mor.
Proof.
exact SUM_compat.
Qed.

Add Parametric Morphism : PRODUCT with
signature eqv \Rightarrow eqv \Rightarrow eqv as PRODUCT_mor.
Proof.
exact PRODUCT_compat.
Qed.

Hint Resolve eqv_ref eqv_sym eqv_trans SUM_compat PRODUCT_compat.

2.2.3 Lemmas

Since Coq now understands the basics of Boolean algebra, it serves as a good exercise for us to generate some further rules using Coq’s proving systems. By doing this, not only do we gain some additional tools that will become handy later down the road, but we also test whether our axioms are behaving as we would like them to.

This is a lemma for a sub-case of term multiplication.

Lemma mul_x_x_plus_T1 :
∀ x, x \times (x + T1) \Rightarrow T0.
Proof.
intros. rewrite distr. rewrite mul_x_x. rewrite mul_comm.
rewrite mul_id. apply sum_x_x.
Qed.

This is a lemma to convert term equivalence to equivalence between their addition and ground term T0, and vice-versa.

Lemma x_equal_y_x_plus_y :
∀ x, y, x \equiv y \leftrightarrow x + y \equiv T0.
Proof.
intros. split.
- intros. rewrite H. rewrite sum_x_x. reflexivity.
- intros. rewrite term_sum_symmetric with (y := y) (z := y). rewrite sum_x_x.
  apply H.
Qed.

Hint Resolve mul_x_x_plus_T1 x_equal_y_x_plus_y.

These lemmas just serve to make certain rewrites regarding the core axioms less tedious.
to write. While one could certainly argue that they should be formulated as axioms and not lemmas due to their triviality, being pedantic is a good exercise.

This is a lemma for identity addition between term and ground term \( T_0 \).

\[ \text{Lemma sum\_id\_sym :} \]
\[ \forall x, x + T_0 == x. \]
\[ \text{Proof.} \]
\[ \text{intros. rewrite sum\_comm. apply sum\_id.} \]
\[ \text{Qed.} \]

Here is a lemma for identity multiplication between term and ground term \( T_1 \).

\[ \text{Lemma mul\_id\_sym :} \]
\[ \forall x, x \times T_1 == x. \]
\[ \text{Proof.} \]
\[ \text{intros. rewrite mul\_comm. apply mul\_id.} \]
\[ \text{Qed.} \]

This is a lemma for multiplication between term and ground term \( T_0 \).

\[ \text{Lemma mul\_T0\_x\_sym :} \]
\[ \forall x, x \times T_0 == T_0. \]
\[ \text{Proof.} \]
\[ \text{intros. rewrite mul\_comm. apply mul\_T0\_x.} \]
\[ \text{Qed.} \]

\[ \text{Lemma sum\_assoc\_opp :} \]
\[ \forall x y z, x + (y + z) == (x + y) + z. \]
\[ \text{Proof.} \]
\[ \text{intros. rewrite sum\_assoc. reflexivity.} \]
\[ \text{Qed.} \]

\[ \text{Lemma mul\_assoc\_opp :} \]
\[ \forall x y z, x \times (y \times z) == (x \times y) \times z. \]
\[ \text{Proof.} \]
\[ \text{intros. rewrite mul\_assoc. reflexivity.} \]
\[ \text{Qed.} \]

\[ \text{Lemma distr\_opp :} \]
\[ \forall x y z, x \times y + x \times z == x \times (y + z). \]
\[ \text{Proof.} \]
\[ \text{intros. rewrite distr. reflexivity.} \]
\[ \text{Qed.} \]

2.3 Variable Sets

Now that the underlying behavior concerning Boolean algebra has been properly articulated to Coq, it is now time to begin formalizing the logic surrounding our meta reasoning of
Boolean equations and systems. While there are certainly several approaches to begin this process, we thought it best to ease into things through formalizing the notion of a set of variables present in an equation.

### 2.3.1 Definitions

We now define a *variable set* to be precisely a list of variables; additionally, we include several functions for including and excluding variables from these variable sets. Furthermore, since uniqueness is not a property guaranteed by Coq lists and it has the potential to be desirable, we define a function that consumes a variable set and removes duplicate entries from it. For convenience, we also provide several examples to demonstrate the functionalities of these new definitions.

Here is a definition of the new type to represent a list (set) of variables (natural numbers).

```coqlang
Definition var_set := list var.
Implicit Type vars: var_set.
```

Here is a simple function to check to see if a variable is in a variable set.

```coqlang
Fixpoint var_set_includes_var (v : var) (vars : var_set) : bool :=
  match vars with
  | nil ⇒ false
  | n :: n' ⇒ if (beq_nat v n) then true
                else var_set_includes_var v n'
  end.
```

Here is a function to remove all instances of var `v` from a list of vars.

```coqlang
Fixpoint var_set_remove_var (v : var) (vars : var_set) : var_set :=
  match vars with
  | nil ⇒ nil
  | n :: n' ⇒ if (beq_nat v n) then (var_set_remove_var v n')
                else n :: (var_set_remove_var v n')
  end.
```

Next is a function to return a unique `var_set` without duplicates. Found vars should be empty for correctness guarantee.

```coqlang
Fixpoint var_set_create_unique (vars : var_set) : var_set :=
  match vars with
  | nil ⇒ nil
  | n :: n' ⇒
    if (var_set_includes_var n n') then var_set_create_unique n'
    else n :: var_set_create_unique n'
  end.
```

This is a function to check if a given `var_set` is unique.

```coqlang
Fixpoint var_set_is_unique (vars : var_set) : bool :=
```

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match \( vars \) with
  | \( \text{nil} \) ⇒ true
  | \( n :: n' \) ⇒
    if (\( \text{var\_set\_includes\_var} \ n \ n' \)) then false
    else \( \text{var\_set\_is\_unique} \ n' \)
end.

This is a function to get the variables of a term as a \( \text{var\_set} \).

Fixpoint \( \text{term\_vars} \ (t : \text{term}) \) : \( \text{var\_set} \) :=
match \( t \) with
  | \( \text{T0} \) ⇒ \( \text{nil} \)
  | \( \text{T1} \) ⇒ \( \text{nil} \)
  | \( \text{VAR} \ x \) ⇒ \( x :: \text{nil} \)
  | \( \text{PRODUCT} \ x \ y \) ⇒ (\( \text{term\_vars} \ x \)) ++ (\( \text{term\_vars} \ y \))
  | \( \text{SUM} \ x \ y \) ⇒ (\( \text{term\_vars} \ x \)) ++ (\( \text{term\_vars} \ y \))
end.

This is a function to generate a list of unique variables that make up a given term.

Definition \( \text{term\_unique\_vars} \ (t : \text{term}) \) : \( \text{var\_set} \) :=
\( \text{var\_set\_create\_unique} \ (\text{term\_vars} \ t) \).

2.3.2 Helper Lemmas for variable sets and lists

Now that we have established the functionality for variable sets, let us prove some properties about them.

Lemma \( \text{vs\_includes\_true} \) : \( \forall \ (x : \text{var}) \ (lvar : \text{list var}), \)
\( \text{var\_set\_includes\_var} \ x \ lvar = \text{true} \rightarrow \text{In} \ x \ lvar. \)
Proof.
  intros.
  induction \( lvar. \)
  - simpl; intros. discriminate.
  - simpl in \( H. \ \text{remember} \ (\text{beq\_nat} \ x \ a) \ as \ H2. \ \text{destruct} \ H2. \)
    + simpl. left. symmetry in \( \text{HeqH2}. \ \text{pose} \ \text{proof} \ \text{beq\_nat\_true} \ as \ H7. \)
      specialization \( (H7 \ x \ a \ \text{HeqH2}). \ \text{symmetry} \ \text{in} \ H7. \ \text{apply} \ H7. \)
      + specialization \( \text{IHlvar} \ H). \ \text{simpl. right}. \ \text{apply} \ \text{IHlvar}. \)
Qed.

Lemma \( \text{vs\_includes\_false} \) : \( \forall \ (x : \text{var}) \ (lvar : \text{list var}), \)
\( \text{var\_set\_includes\_var} \ x \ lvar = \text{false} \rightarrow \neg \text{In} \ x \ lvar. \)
Proof.
  intros.
  induction \( lvar. \)
  - simpl; intros. unfold not. intros. destruct \( H0. \)
  - simpl in \( H. \ \text{remember} \ (\text{beq\_nat} \ x \ a) \ as \ H2. \ \text{destruct} \ H2. \ \text{inversion} \ H. \)
Lemma in_dup_and_non_dup : ∀ (x : var) (lvar : list var),
In x lvar ↔ In x (var_set_create_unique lvar).
Proof.
intros. split.
- induction lvar.
  + intros. simpl in H. destruct H.
  + intros. simpl. remember (var_set_includes_var a lvar) as C. destruct C.
    × symmetry in H eqC. pose proof vs_includes_true as H7.
      specialize (H7 a lvar H eqC). simpl in H. destruct H.
      - rewrite H in H7. specialize (IHlvar H7). apply IHlvar.
      - specialize (IHlvar H). apply IHlvar.
    × symmetry in H eqC. pose proof vs_includes_false as H7.
      specialize (H7 a lvar H eqC). simpl in H. destruct H.
      - simpl. left. apply H.
      - specialize (IHlvar H). simpl. right. apply IHlvar.
- induction lvar.
  + intros. simpl in H. destruct H.
  + intros. simpl in H. remember (var_set_includes_var a lvar) as C. destruct C.
    × symmetry in H eqC. pose proof vs_includes_true as H7.
      specialize (H7 a lvar H eqC). specialize (IHlvar H). simpl.
      right. apply IHlvar.
    × symmetry in H eqC. pose proof vs_includes_false as H7.
      specialize (H7 a lvar H eqC). simpl in H. destruct H.
      - simpl. left. apply H.
      - specialize (IHlvar H). simpl. right. apply IHlvar.
Qed.

2.3.3 Examples
Below are some examples of the behaviors of variable sets.
Example var_set_create_unique_ex1 :
var_set_create_unique [0;5;2;1;1;2;2;9;5;3] = [0;1;2;9;5;3].
Proof.
  simpl. reflexivity.
Qed.
Example \texttt{var\_set\_is\_unique\_ex1}:
\begin{verbatim}
  var_set_is_unique [0;2;2;2] = false.
\end{verbatim}
Proof.
\begin{verbatim}
  simpl. reflexivity.
\end{verbatim}
Qed.

Here are examples to demonstrate the correctness of the function \texttt{term\_vars} on specific cases.

Example \texttt{term\_vars\_ex1}:
\begin{verbatim}
  term_vars (VAR 0 + VAR 0 + VAR 1) = [0;0;1].
\end{verbatim}
Proof.
\begin{verbatim}
  simpl. reflexivity.
\end{verbatim}
Qed.

Example \texttt{term\_vars\_ex2}:
\begin{verbatim}
  In 0 (term_vars (VAR 0 + VAR 0 + VAR 1)).
\end{verbatim}
Proof.
\begin{verbatim}
  simpl. left. reflexivity.
\end{verbatim}
Qed.

\section{2.4 Ground Terms}

Seeing as we just outlined the definition of a variable set, it seems fair to now formalize the definition of a ground term, or in other words, a term that has no variables and whose variable set is the empty set.

\subsection{2.4.1 Definitions}

A \textit{ground term} is a recursively defined proposition that is only true if and only if no variable appears in it; otherwise it will be a false proposition and no longer a ground term.

In this subsection we declare definitions related to ground terms, including functions and lemmas.

This is a function to check if a given term is a ground term (i.e. has no vars).

\textbf{Fixpoint \texttt{ground\_term} (}t : \texttt{term}) : \texttt{Prop} :=
\begin{verbatim}
  match t with
    | VAR x ⇒ False
    | SUM x y ⇒ ground_term x ∧ ground_term y
    | PRODUCT x y ⇒ ground_term x ∧ ground_term y
    | _ ⇒ True
  end.
\end{verbatim}
2.4.2 Lemmas

Our first real lemma (shown below), articulates an important property of ground terms: all ground terms are equivalent to either 0 or 1. This curious property is a direct result of the fact that these terms possess no variables and additionally because of the axioms of Boolean algebra.

This is a lemma (trivial, intuitively true) that proves that if the function ground_term returns true then it is either T0 or T1.

**Lemma ground_term_equiv_T0_T1**: \( \forall x, \) ground_term \( x \rightarrow x == T0 \lor x == T1. \)

**Proof.**
- intros. induction \( x. \)
  - left. reflexivity.
  - right. reflexivity.
  - contradiction.
  - inversion \( H. \) destruct \( IHx1; \) destruct \( IHx2; \) auto. rewrite \( H2. \) left.
    - rewrite \( sum_id. \) apply \( H3. \) rewrite \( H2. \) rewrite \( H3. \) rewrite \( sum_id. \) right.
    - reflexivity. rewrite \( H2. \) rewrite \( H3. \) right. rewrite \( sum_comm. \)
    - rewrite \( sum_id. \) reflexivity. rewrite \( H2. \) rewrite \( H3. \) rewrite \( sum_x.x. \) left.
      - reflexivity.
  - inversion \( H. \) destruct \( IHx1; \) destruct \( IHx2; \) auto. rewrite \( H2. \) left.
    - rewrite \( mul_T0.x. \) reflexivity. rewrite \( H2. \) left. rewrite \( mul_T0.x. \)
      - reflexivity. rewrite \( H3. \) left. rewrite \( mul_comm. \) rewrite \( mul_T0.x. \)
        - reflexivity. rewrite \( H2. \) rewrite \( H3. \) right. rewrite \( mul_id. \) reflexivity.
Qed.

This lemma, while intuitively obvious by definition, nonetheless provides a formal bridge between the world of ground terms and the world of variable sets.

**Lemma ground_term_has_empty_var_set**: \( \forall x, \) ground_term \( x \rightarrow \text{term_vars } x = []. \)

**Proof.**
- intros. induction \( x. \)
  - simpl. reflexivity.
  - simpl. reflexivity.
  - contradiction.
  - firstorder. unfold \( \text{term_vars.} \) unfold \( \text{term_vars in } H2. \) rewrite \( H2. \)
    - unfold \( \text{term_vars in } H1. \) rewrite \( H1. \) simpl. reflexivity.
  - firstorder. unfold \( \text{term_vars.} \) unfold \( \text{term_vars in } H2. \) rewrite \( H2. \)
    - unfold \( \text{term_vars in } H1. \) rewrite \( H1. \) simpl. reflexivity.
Qed.
2.4.3 Examples

Here are some examples to show that our ground term definition is working appropriately.

Example ex_gt1:
  ground_term (T0 + T1).
Proof.
  simpl. split.
  - reflexivity.
  - reflexivity.
Qed.

Example ex_gt2:
  ground_term (VAR 0 x T1) → False.
Proof.
  simpl. intros. destruct H. apply H.
Qed.

2.5 Substitutions

It is at this point in our Coq development that we begin to officially define the principal action around which the entirety of our efforts are centered: the act of substituting variables with other terms. While substitutions alone are not of great interest, their emergent properties as in the case of whether or not a given substitution unifies an equation are of substantial importance to our later research.

2.5.1 Definitions

In this subsection we make the fundamental definitions of substitutions, basic functions for them, accompanying lemmas and some propositions.

Here we define a substitution to be a list of ordered pairs where each pair represents a variable being mapped to a term. For sake of clarity these ordered pairs shall be referred to as replacements from now on and as a result, substitutions should really be considered to be lists of replacements.

Definition replacement := prod var term.

We define a new type susbt to represent a substitution as a list of replacements.

Definition subst := list replacement.

Implicit Type s : subst.

Our first function, find_replacement, is an auxilliary to apply_subst. This function will search through a substitution for a specific variable, and if found, returns the variable's associated term.

Fixpoint find_replacement (x : var) (s : subst) : term :=
match $s$ with
| nil ⇒ $\text{VAR } x$
| $r :: r' ⇒$
  if $\text{beq_nat} (\text{fst } r) x$ then $\text{snd } r$
  else $\text{find}\_\text{replacement} x r'$
end.

The $\text{apply}\_\text{subst}$ function will take a term and a substitution and will produce a new term reflecting the changes made to the original one.

Fixpoint $\text{apply}\_\text{subst} (t : \text{term}) (s : \text{subst}) : \text{term} :=$
match $t$ with
| $\text{T0}$ ⇒ $\text{T0}$
| $\text{T1}$ ⇒ $\text{T1}$
| $\text{VAR } x$ ⇒ $\text{find}\_\text{replacement} x s$
| $\text{PRODUCT} x y$ ⇒ $\text{PRODUCT} (\text{apply}\_\text{subst} x s) (\text{apply}\_\text{subst} y s)$
| $\text{SUM} x y$ ⇒ $\text{SUM} (\text{apply}\_\text{subst} x s) (\text{apply}\_\text{subst} y s)$
end.

For reasons of completeness, it is useful to be able to generate identity substitutions; namely, substitutions that map the variables of a term to themselves.

This is a function when given a list of variables builds a list of identity substitutions - one for each variable.

Fixpoint $\text{build}\_\text{id}\_\text{subst} (lvar : \text{var}\_\text{set}) : \text{subst} :=$
match $lvar$ with
| nil ⇒ nil
| $v :: v'$ ⇒ $(v, (\text{VAR } v)) :: \text{build}\_\text{id}\_\text{subst} v'$
end.

Since we now have the ability to generate identity substitutions, we should now formalize a general proposition for testing whether or not a given substitution is an identity substitution of a given term.

Definition $\text{subst}\_\text{equiv} (s1 \text{ s2} : \text{subst}) : \text{Prop} :=$
$\forall t, \text{apply}\_\text{subst} t s1 == \text{apply}\_\text{subst} t s2.$

Definition $\text{subst}\_\text{is}\_\text{id}\_\text{subst} (t : \text{term}) (s : \text{subst}) : \text{Prop} :=$
$\text{apply}\_\text{subst} t s == t.$

Given we now have definitions for substitutions, we should now introduce the idea of a substitution composing another one.

Fixpoint $\text{subst}\_\text{compose} (s s' : \text{subst}) : \text{subst} :=$
match $s'$ with
| $[]$ ⇒ $s$
| $(x, t) :: s''$ ⇒ $(x, \text{apply}\_\text{subst} t s) :: (\text{subst}\_\text{compose} s s'')$
end.
Here we define the domain of a substitution, namely the list of variables for which the substitution has a mapping (replacement). Essentially this acts as a list of all the first parts of the replacement.

Definition subst_domain (sig : subst) : list var :=
  map (fun r ⇒ (fst r)) sig.

We define the concept of a sub list. If an element is a member of a list, it is then a member of the other list as well.

Definition sub_dmn_list (l1 : list var) (l2 : list var) : Prop :=
  ∀ (x : var), In x l1 → In x l2.

2.5.2 Helper Lemmas for the apply_subst function

Having now outlined the functionality of a substitution, let us now begin to analyze some implications of its form and composition by proving some lemmas.

Given that we have a definition for identity substitutions, we should prove that identity substitutions do not modify a term.

Lemma id_subst: ∀ (t : term) (l : var_set),
  apply_subst t (build_id_subst l) == t.
Proof.
  intros. induction t.
  - simpl. reflexivity.
  - simpl. reflexivity.
  - simpl. induction l.
    + simpl. reflexivity.
    + simpl. destruct (beq_nat a v) eqn: e.
      × apply beq_nat_true in e. rewrite e. reflexivity.
      × apply IHl.
    - simpl. rewrite IHt1. rewrite IHt2. reflexivity.
    - simpl. rewrite IHt1. rewrite IHt2. reflexivity.
  Qed.

These are helper lemmes for the apply_subst properties.

Lemma sum_comm_compat t1 t2: ∀ (sigma: subst),
  apply_subst (t1 + t2) sigma == apply_subst (t2 + t1) sigma.
Proof.
  intros. simpl. auto.
Qed.

Hint Resolve sum_comm_compat.

Lemma sum_assoc_compat t1 t2 t3: ∀ (sigma: subst),
  apply_subst ((t1 + t2) + t3) sigma == apply_subst (t1 + (t2 + t3)) sigma.
Proof.
  intros. simpl. auto.

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Qed.

Hint Resolve sum_assoc_compat.

Lemma sum_id_compat t: ∀ (sigma: subst),
  apply_subst (T0 + t) sigma == apply_subst t sigma.
Proof.
  intros. simpl. auto.
Qed.

Hint Resolve sum_id_compat.

Lemma sum_x_x_compat t: ∀ (sigma: subst),
  apply_subst (t + t) sigma == apply_subst T0 sigma.
Proof.
  intros. simpl. auto.
Qed.

Hint Resolve sum_x_x_compat.

Lemma mul_comm_compat t1 t2: ∀ (sigma: subst),
  apply_subst (t1 × t2) sigma == apply_subst (t2 × t1) sigma.
Proof.
  intros. simpl. auto.
Qed.

Hint Resolve mul_comm_compat.

Lemma mul_assoc_compat t1 t2 t3: ∀ (sigma: subst),
  apply_subst ((t1 × t2) × t3) sigma == apply_subst (t1 × (t2 × t3)) sigma.
Proof.
  intros. simpl. auto.
Qed.

Hint Resolve mul_assoc_compat.

Lemma mul_x_x_compat t: ∀ (sigma: subst),
  apply_subst (t × t) sigma == apply_subst t sigma.
Proof.
  intros. simpl. auto.
Qed.

Hint Resolve mul_x_x_compat.

Lemma mul_T0_x_compat t: ∀ (sigma: subst),
  apply_subst (T0 × t) sigma == apply_subst T0 sigma.
Proof.
  intros. simpl. auto.
Qed.

Hint Resolve mul_T0_x_compat.

Lemma mul_id_compat t: ∀ (sigma: subst),
  apply_subst (T1 × t) sigma == apply_subst t sigma.
Proof.
intros. simpl. auto.
Qed.

Hint Resolve mul_id_compat.

Lemma distr_compat t1 t2 t3 : \( \forall \) (sigma : subst),
  apply_subst \((t1 \times (t2 + t3))\) sigma ==
  apply_subst \(((t1 \times t2) + (t1 \times t3))\) sigma.
Proof.
  intros. simpl. auto.
Qed.

Hint Resolve distr_compat.

Lemma refl_comm_compat t1 t2 : \( \forall \) (sigma : subst),
  apply_subst t1 sigma == apply_subst t2 sigma \rightarrow
  apply_subst t2 sigma == apply_subst t1 sigma.
Proof.
  intros. simpl. auto.
Qed.

Hint Resolve refl_comm_compat.

Lemma trans_compat t1 t2 t3 : \( \forall \) (sigma : subst),
  apply_subst t1 sigma == apply_subst t2 sigma \rightarrow
  apply_subst t2 sigma == apply_subst t3 sigma \rightarrow
  apply_subst t1 sigma == apply_subst t3 sigma.
Proof.
  intros. eauto.
Qed.

Hint Resolve trans_compat.

Lemma trans_compat2 c1 c2 c3 :
  c1 == c2 \rightarrow
  c2 == c3 \rightarrow
  c1 == c3.
Proof.
  intros. eauto.
Qed.

This is an axiom that states that if two terms are equivalent then applying any substitution on them will also produce equivalent terms. The reason we axiomatized this and we did not prove it as a lemma is because the set of our fundamental axioms is not an inductive relation, so it would be impossible to prove the lemma below with our fundamental axioms in the current format.

Axiom apply_subst_compat : \( \forall \) (t t' : term),
  t == t' \rightarrow
  \forall (sigma : subst), apply_subst t sigma == apply_subst t' sigma.

Add Parametric Morphism : apply_subst with
signature eqv => eq => eqv as apply_subst_mor.

Proof.
  exact apply_subst_compat.
Qed.

This is a simple lemma that states that an empty substitution cannot modify a term.

Lemma subst_empty_no_change :
  \forall (t : \text{term}), (\text{apply_subst} \ t \ [] ) == \ t.
Proof.
  intros. induction \ t.
  \ - simpl. reflexivity.
  \ - simpl. reflexivity.
  \ - simpl. reflexivity.
  \ - simpl. rewrite IHt1. rewrite IHt2. reflexivity.
  \ - simpl. rewrite IHt1. rewrite IHt2. reflexivity.
Qed.

An intuitive thing to prove for ground terms is that they cannot be modified by applying substitutions to them. This will later prove to be very relevant when we begin to talk about unification.

This is a helpful lemma for showing substitutions do not affect ground terms.

Lemma ground_term_cannot_subst :
  \forall \ x,
  \text{ground_term} \ \ x \rightarrow
  \forall \ s, \text{apply_subst} \ \ x \ \ s == \ x.
Proof.
  intros. induction \ s.
  \ - apply \text{ground_term_equiv_T0_T1} in \ H. destruct \ H.
  \ + rewrite \ H. simpl. reflexivity.
  \ + rewrite \ H. simpl. reflexivity.
  \ - apply \text{ground_term_equiv_T0_T1} in \ H. destruct \ H. rewrite \ H.
  \ + simpl. reflexivity.
  \ + rewrite \ H. simpl. reflexivity.
Qed.

A fundamental property of substitutions is their distributivity across the summation and multiplication of terms. Again the importance of these proofs will not become apparent until we talk about unification.

This is a useful lemma for showing the distributivity of substitutions across term summation.

Lemma subst_sum_distribution :
  \forall \ s \ \ x \ \ y,
  \text{apply_subst} \ \ x \ \ s + \text{apply_subst} \ \ y \ \ s == \text{apply_subst} \ (x + y) \ s.
Proof.
  intro. induction \ s.
  \ - simpl. intros. reflexivity.
This is a lemma to prove the distributivity of the \texttt{apply\_subst} function across term multiplication.

\textbf{Lemma subst\_mul\_distribution} : \forall s x y,
\begin{align*}
\text{apply\_subst } x s \times \text{apply\_subst } y s &= \text{apply\_subst } (x \times y) s.
\end{align*}
\textbf{Proof}.
\begin{itemize}
\item intro. induction \( s \).
\item intros. reflexivity.
\item intros. simpl. reflexivity.
\end{itemize}
Qed.

Here is a lemma to prove the opposite of summation distributivity of the \texttt{apply\_subst} function across term summation.

\textbf{Lemma subst\_sum\_distr\_opp} : \forall s x y,
\begin{align*}
\text{apply\_subst } (x + y) s &= \text{apply\_subst } x s + \text{apply\_subst } y s.
\end{align*}
\textbf{Proof}.
\begin{itemize}
\item intros.
\item apply \texttt{refl\_comm}.
\item apply subst\_sum\_distribution.
\end{itemize}
Qed.

This is a lemma to prove the opposite of multiplication distributivity of the \texttt{apply\_subst} function across term summation.

\textbf{Lemma subst\_mul\_distr\_opp} : \forall s x y,
\begin{align*}
\text{apply\_subst } (x \times y) s &= \text{apply\_subst } x s \times \text{apply\_subst } y s.
\end{align*}
\textbf{Proof}.
\begin{itemize}
\item intros.
\item apply \texttt{refl\_comm}.
\item apply subst\_mul\_distribution.
\end{itemize}
Qed.

This is an intutitive lemmas to apply a single replacement substitution on a VAR term.

\textbf{Lemma var\_subst} : \forall (v : \text{var}) (ts : \text{term}),
\begin{align*}
\text{apply\_subst } (\text{VAR } v) [(v , ts)] &= ts.
\end{align*}
\textbf{Proof}.
\begin{itemize}
\item intros. simpl. destruct (beq\_nat \( v v \) \textit{eqn: } e).
\item - apply beq\_nat\_true \textit{in } e. reflexivity.
\item - apply beq\_nat\_false \textit{in } e. firstorder.
\end{itemize}
Qed.

\section*{2.5.3 Examples}
Here are some examples showcasing the nature of applying substitutions to terms.
Example subst_ex1:
  apply_subst (T0 + T1) [] == T0 + T1.
Proof.
  intros. reflexivity.
Qed.

Example subst_ex2:
  apply_subst (VAR 0 × VAR 1) [(0, T0)] == T0.
Proof.
  intros. simpl. apply mul_T0_x.
Qed.

2.5.4 Auxillary Definitions for Substitutions and Terms

In this section we define more helper functions and lemmas related to substitutions and ground terms. Specifically we are defining a ground term, a ground substitution, a “01” term, a “01” substitution, and a substitution composition. A ground_term is a term with no variables in it. The terms that are used more in the future proofs are the “01” term and “01” substitution. A “01” term is a term that is either exactly equal to T0 or T1. A “01” substitution is a substitution in which each variable (or the first part of each replacement) is mapped to a “01” term. A “01” term is not necessarily a ground term (but it might be) and a “01” substitution is not necessarily a ground substitution (but it might be). In the proof file, we are mostly using the “01” term and substitution terminology.

We define a proposition for a ground_subst. A substitution is ground when in all of its replacements, the second part is a ground_term.

Fixpoint ground_subst (sig : subst) : Prop :=
  match sig with
  | [] ⇒ True
  | r :: r' ⇒ ground_term (snd r) ∧ ground_subst r'
  end.

This is a function to determine whether a term is a ground term, by returning a boolean.

Fixpoint is_ground_term (t : term) : bool :=
  match t with
  | T0 ⇒ true
  | T1 ⇒ true
  | VAR x ⇒ false
  | SUM a b ⇒ (is_ground_term a) && (is_ground_term b)
  | PRODUCT a b ⇒ (is_ground_term a) && (is_ground_term b)
  end.

This is a function to determine whether a substitution is a ground substitution, by returning a boolean.

Fixpoint is_ground_subst (sig : subst) : bool :=
existsb is_ground_term (map snd sig).

This is a function to determine whether a term is a T0 or T1 term by returning a boolean.

Definition is_01_term (t : term) : bool :=
  match t with
  | T0 ⇒ true
  | T1 ⇒ true
  | _ ⇒ false
end.

This is a function to determine whether a substitution is a “01” substitution by returning a boolean, meaning that each second part of every replacement is either a T0 or a T1 term.

Fixpoint is_01_subst (sig : subst) : bool :=
  existsb is_01_term (map snd sig).

This is a function to determine whether a term is a T0 or T1 term by returning a proposition.

Fixpoint _01_term (t : term) : Prop :=
  match t with
  | T0 ⇒ True
  | T1 ⇒ True
  | _ ⇒ False
end.

This is a function to determine whether a substitution is a “01” substitution by returning a proposition, meaning that each second part of every replacement is either a T0 or a T1 term.

Fixpoint _01_subst (sig : subst) : Prop :=
  match sig with
  | [] ⇒ True
  | r :: r' ⇒ _01_term (snd r) ∧ _01_subst r'
end.

2.6 Unification

Now that we have established the concept of term substitutions in Coq, it is time for us to formally define the concept of Boolean unification. Unification, in its most literal sense, refers to the act of applying a substitution to terms in order to make them equivalent to each other. In other words, to say that two terms are unifiable is to really say that there exists a substitution such that the two terms are equal. Interestingly enough, we can abstract this concept further to simply saying that a single term is unifiable if there exists a substitution such that the term will be equivalent to 0. By doing this abstraction, we can prove that equation solving and unification are essentially the same fundamental problem.
Below is the initial definition for unification, namely that two terms can be unified to be equivalent to one another. By starting here we will show each step towards abstracting unification to refer to a single term.

Proposition that a given substitution unifies (namely, makes equivalent), two given terms

Definition \( \text{unifies}(a \ b : \text{term}) \ (s : \text{subst}) : \text{Prop} := \) \( \text{apply} \ _\text{subst} \ a \ s \ == \ \text{apply} \ _\text{subst} \ b \ s. \)

Here is a simple example demonstrating the concept of testing whether two terms are unified by a substitution.

Example \( \text{ex\_unif1} : \) \( \text{unifies} \ (\text{VAR} \ 0) \ (\text{VAR} \ 1) \ [(0, \ T1); \ (1, \ T1)]. \)

Proof.
\( \text{unfold unifies. simpl. reflexivity. Qed.} \)

Now we are going to show that moving both terms to one side of the equivalence relation through addition does not change the concept of unification.

This is a proposition that a given substitution makes equivalent the sum of two terms when the substitution is applied to each of them, and ground term \( T0. \)

Definition \( \text{unifies\_T0}(a \ b : \text{term}) \ (s : \text{subst}) : \text{Prop} := \) \( \text{apply} \ _\text{subst} \ a \ s \ + \ \text{apply} \ _\text{subst} \ b \ s == T0. \)

This is a lemma that proves that finding a unifier for \( x = y \) is the same as finding a unifier for \( x + y = 0. \)

Lemma \( \text{unifies\_T0\_equiv} : \forall \ x \ y \ s, \) \( \text{unifies} \ x \ y \ s \ \leftrightarrow \ \text{unifies\_T0} \ x \ y \ s. \)

Proof.
\( \text{intros. split.} \)
\( \text{- intros. unfold unifies\_T0. unfold unifies in H. rewrite H.} \)
\( \text{rewrite sum\_x\_x. reflexivity.} \)
\( \text{- intros. unfold unifies\_T0 in H. unfold unifies.} \)
\( \text{rewrite term\_sum\_symmetric with} \ ((x := \text{apply} \ _\text{subst} \ x \ s \ + \ \text{apply} \ _\text{subst} \ y \ s)) \)
\( (z := \text{apply} \ _\text{subst} \ y \ s) \ \text{in H. rewrite sum\_id in H.} \)
\( \text{rewrite sum\_comm in H.} \)
\( \text{rewrite sum\_comm with} \ ((y := \text{apply} \ _\text{subst} \ y \ s)) \ \text{in H.} \)
\( \text{rewrite \( \rightarrow \ \text{sum\_assoc} \ \text{in H.} \)} \)
\( \text{rewrite sum\_x\_x in H.} \)
\( \text{rewrite sum\_id in H.} \)
\( \text{apply H.} \)

Qed.

Now we can define what it means for a substitution to be a unifier for a given term.

Here is a proposition that a given substitution unifies a given term, namely it makes it equivalent with \( T0. \)
Definition unifier \((t : \text{term}) (s : \text{subst}) : \text{Prop} :=\)
apply subst \(t s == T0\).

Example unifier ex1 :
unifier (VAR 0) [(0, T0)].

Proof.
unfold unifier. simpl. reflexivity.
Qed.

To ensure our efforts were not in vain, let us now prove that this last abstraction of the unification problem is still equivalent to the original.

This is a lemma that proves that the unifier proposition can distributes over addition of terms.

Lemma unifier_distribution : \(\forall x y s,\)
unifies \(T0 x y s \leftrightarrow \text{unifier} (x + y) s\).

Proof.
intros. split.
- intros. unfold unifies_T0 in \(H\). unfold unifier.
  rewrite \(\leftarrow H\). symmetry. apply subst_sum_distribution.
- intros. unfold unifies_T0. unfold unifier in \(H\).
  rewrite \(\leftarrow H\). apply subst_sum_distribution.

Qed.

Lastly let us define a term to be unifiable if there exists a substitution that unifies it.
This is a proposition that states when a term is unifiable.

Definition unifiable \((t : \text{term}) : \text{Prop} :=\)
\(\exists s, \text{unifier} t s\).

Example unifiable_ex1 :
\(\exists x, \text{unifiable} (x + T1)\).

Proof.
\(\exists T1. \text{unfold unifiable}. \text{unfold unifier}.
\exists []\). simpl. rewrite sum_x_x. reflexivity.

Qed.

2.7 Most General Unifier

In this subsection we define propositions, lemmas and examples related to the most general unifier.

While the property of a term being unifiable is certainly important, it should come as no surprise that not all unifiers are created equal; in fact, certain unifiers possess the desirable property of being more general than others. For this reason, let us now formally define the concept of a most general unifier (mgu): a unifier such that with respect to a given term, all other unifiers are instances of it, or in other words, less general than it.
The first step towards establishing the concept of a mgu requires us to formalize the notion
of a unifier being more general than another. To accomplish this goal, let us formulate
the definition of a substitution composing another one; or in other words, to say that a
substitution is more general than another one.

This is a proposition of sequential substitution application.

**Definition** \( \text{substitution factor through} \ (s \ s' \ \delta : \ \text{subst}) \) : Prop :=
\[
\forall \ (x : \text{var}), \ \text{apply subst} \ (\text{apply subst} \ (\text{VAR} \ x) \ s) \ \delta == \\
\text{apply subst} \ (\text{VAR} \ x) \ s'.
\]

This is the definition of a more general substitution.

**Definition** \( \text{more general substitution} \ (s \ s' : \ \text{subst}) \) : Prop :=
\[
\exists \ \delta, \ \text{substitution factor through} \ s \ s' \ \delta.
\]

Now that we have articulated the concept of composing substitutions, let us now formulate
the definition for a most general unifier.

This is the definition of a Most General Unifier (mgu): A Most General Unifier (MGU)
takes in a term and a substitution and tells whether or not said substitution is an mgu for
the given term.

**Definition** \( \text{most general unifier} \ (t : \ \text{term}) \ (s : \ \text{subst}) \) : Prop :=
\[
\text{unifier} \ t \ s \ \land \\
\forall \ (s' : \ \text{subst}), \\
\text{unifier} \ t \ s' \ \rightarrow \\
\text{more general substitution} \ s \ s'.
\]

While this definition of a most general unifier is certainly valid, we can also characterize a
unifier by other similar properties. For this reason, let us now define an alternative definition
called a reproductive unifier, and then prove it to be equivalent to our definition of a most
general unifier. This will make our proofs easier to formulate down the road as the task of
proving a unifier to be reproductive is substantially easier than proving it to be most general
directly.

**Definition** \( \text{reproductive unifier} \ (t : \ \text{term}) \ (\text{sig} : \ \text{subst}) \) : Prop :=
\[
\text{unifier} \ t \ \text{sig} \ \land \\
\forall \ (\text{tau} : \ \text{subst}) \ (x : \ \text{var}), \\
\text{unifier} \ t \ \text{tau} \ \rightarrow \\
\text{apply subst} \ (\text{apply subst} \ (\text{VAR} \ x) \ \text{sig}) \ \text{tau} == \text{apply subst} \ (\text{VAR} \ x) \ \text{tau}.
\]

This is a lemma to show that a reproductive unifier is a most general unifier. Since the
ultimate goal is to prove that a specific algorithm produces an mgu then if we could prove it
is a reproductive unifier then we could use this lemma to arrive at the desired conclusion.

**Lemma** \( \text{reproductive is mgu} : \forall \ (t : \ \text{term}) \ (u : \ \text{subst}), \\
\text{reproductive unifier} \ t \ u \ \rightarrow \\
\text{most general unifier} \ t \ u. \)

**Proof.**
\[
\text{intros. unfold most general unifier. unfold reproductive unifier in } H.
\]
unfold more_general_substitution . destruct H. split.
- apply H.
- intros. specialize (H0 s'). \exists s'. unfold substitution_factor_through.
  intros. specialize (H0 x).
  specialize (H0 H1). apply H0.
Qed.

This is a lemma to show that if two terms are equivalent then for any substitution that is an mgu of one of the terms, then it is an mgu of the other term as well.

Lemma most_general_unifier_compat : \forall (t t' : term),
  t == t' \rightarrow
  \forall (sigma : subst),
  most_general_unifier t sigma \leftrightarrow most_general_unifier t' sigma.
Proof.
  intros. split.
  - intros. unfold most_general_unifier. unfold unifier in H0.
    unfold unifier in *. split.
     + unfold most_general_unifier in H0. destruct H0. unfold unifier in H0.
      rewrite H in H0. apply H0.
     + intros. unfold most_general_unifier in H0. destruct H0.
      specialize (H2 s'). unfold unifier in H0. symmetry in H. rewrite H in H1.
      unfold unifier in H2. specialize (H2 H1). apply H2.
  - unfold most_general_unifier. intros. destruct H0 . split.
    + symmetry in H. unfold unifier in H0. rewrite H in H0. unfold unifier.
     apply H0.
    + intros. specialize (H1 s'). unfold unifier in H2. rewrite H in H2.
     unfold unifier in H1. specialize (H1 H2). apply H1.
Qed.

2.8 Auxilliary Computational Operations and Simplifications

These functions below will come in handy later during the Lowenheim formula proof. They mainly lay the groundwork for providing the computational nuts and bolts for Lowenheim's algorithm for finding most general unifiers and initial ground unifiers.

This is a function to check if two terms are exactly identical.

Fixpoint identical (a b: term) : bool :=
match a , b with
| T0, T0 ⇒ true
| T0, _ ⇒ false
| T1, T1 ⇒ true
| T1 , _ ⇒ false  
| VAR x , VAR y ⇒ if beq_nat x y then true else false 
| VAR x , _ ⇒ false  
| PRODUCT x y , PRODUCT x1 y1 ⇒ identical x x1 && identical y y1 
| PRODUCT x y , _ ⇒ false 
| SUM x y , SUM x1 y1 ⇒ identical x x1 && identical y y1 
| SUM x y , _ ⇒ false  
end. 

This is basic addition for terms.

Definition plus_one_step (a b : term) : term := 
match a, b with 
| T0 , T0 ⇒ T0  
| T0 , T1 ⇒ T1  
| T1 , T0 ⇒ T1  
| T1 , T1 ⇒ T0  
| _ , _ ⇒ SUM a b  
end. 

This is basic multiplication for terms.

Definition mult_one_step (a b : term) : term := 
match a, b with 
| T0 , T0 ⇒ T0  
| T0 , T1 ⇒ T0  
| T1 , T0 ⇒ T0  
| T1 , T1 ⇒ T1  
| _ , _ ⇒ PRODUCT a b  
end. 

This is a function to simplify a term in very apparent and basic ways. They are only simplified if they are ground terms.

Fixpoint simplify (t : term) : term := 
match t with 
| T0 ⇒ T0  
| T1 ⇒ T1  
| VAR x ⇒ VAR x  
| PRODUCT x y ⇒ mult_one_step (simplify x) (simplify y)  
| SUM x y ⇒ plus_one_step (simplify x) (simplify y)  
end. 

Some lemmas follow to prove intuitive facts for the basic multiplication and addition of terms, leading up to proving the simplify_eqv lemma.

Lemma pos_left_sum_compat : ∀ (t t1 t2 : term), 
t == t1 → plus_one_step t1 t2 == plus_one_step t t2.
Proof.

intros. induction t1.
- induction t.
  + reflexivity.
  + apply T1\_\text{not\_equiv\_T0} in H. inversion H.
    + induction t2.
      × simpl. rewrite H. rewrite sum_x_x. reflexivity.
      × simpl. rewrite H. rewrite sum_id. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
    + induction t2.
      × simpl. rewrite H. rewrite sum_x_x. reflexivity.
      × simpl. rewrite H. rewrite sum_id. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
  + induction t2.
    × simpl. rewrite H. rewrite sum_x_x. reflexivity.
    × simpl. rewrite H. rewrite sum_id. reflexivity.
    × simpl. rewrite H. reflexivity.
    × simpl. rewrite H. reflexivity.
    × simpl. rewrite H. reflexivity.
  - induction t.
    + induction t2.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
    + induction t2.
      × simpl. reflexivity.
      × simpl. reflexivity.
      × simpl. reflexivity.
      × simpl. reflexivity.
      × simpl. reflexivity.
    + induction t2.
      × simpl. rewrite H. rewrite sum_comm. rewrite sum_id. reflexivity.
      × simpl. rewrite H. rewrite sum_x_x. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
+ induction t2.
  × simpl. rewrite H. rewrite sum_comm. rewrite sum_id. reflexivity.
  × simpl. rewrite H. rewrite sum_x_x. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
+ induction t2.
  × simpl. rewrite H. rewrite sum_comm. rewrite sum_id. reflexivity.
  × simpl. rewrite H. rewrite sum_x_x. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
+ induction t2.
  × simpl. rewrite H. rewrite sum_comm. rewrite sum_id. reflexivity.
  × simpl. rewrite H. rewrite sum_x_x. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
+ induction t2.
  × simpl. rewrite H. rewrite sum_comm. rewrite sum_id. reflexivity.
  × simpl. rewrite H. rewrite sum_x_x. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
\times \text{simpl. rewrite } H. \text{ reflexivity.}

- \text{induction } t.
  + \text{induction } t_2.
    \times \text{simpl. rewrite } \leftarrow H. \text{ rewrite } \text{sum}_x \cdot \text{x. reflexivity.}
    \times \text{simpl. rewrite } \leftarrow H. \text{ rewrite } \text{sum}_\text{id}. \text{ reflexivity.}
    \times \text{simpl. rewrite } \leftarrow H. \text{ reflexivity.}
    \times \text{simpl. rewrite } \leftarrow H. \text{ reflexivity.}
    \times \text{simpl. rewrite } \leftarrow H. \text{ reflexivity.}
    \times \text{simpl. rewrite } \leftarrow H. \text{ reflexivity.}
  + \text{induction } t_2.
    \times \text{simpl. rewrite } \leftarrow H. \text{ rewrite } \text{sum}_\text{comm}. \text{ rewrite } \text{sum}_\text{id}. \text{ reflexivity.}
    \times \text{simpl. rewrite } H. \text{ rewrite } \text{sum}_x \cdot \text{x. reflexivity.}
    \times \text{simpl. rewrite } \leftarrow H. \text{ reflexivity.}
    \times \text{simpl. rewrite } \leftarrow H. \text{ reflexivity.}
    \times \text{simpl. rewrite } \leftarrow H. \text{ reflexivity.}
    + \text{induction } t_2.
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
    + \text{induction } t_2.
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
    + \text{induction } t_2.
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
      \times \text{simpl. rewrite } H. \text{ reflexivity.}
    \times \text{simpl. rewrite } H. \text{ reflexivity.}
    \times \text{simpl. rewrite } \leftarrow H. \text{ rewrite } \text{sum}_\text{comm}. \text{ rewrite } \text{sum}_\text{id}. \text{ reflexivity.}
    \times \text{simpl. rewrite } H. \text{ rewrite } \text{sum}_x \cdot \text{x. reflexivity.}
    + \text{induction } t_2.
Lemma pos_right_sum_compat : \forall (t t1 t2 : term),
\text{t} == \text{t2} \rightarrow \text{plus_one_step t1 t2} == \text{plus_one_step t1 t}.

\text{Proof.}
\text{intros. induction t1.}
- induction t.
  + induction t2.
    \times simpl. reflexivity.
    \times simpl. rewrite \text{H. reflexivity}.
    \times simpl. rewrite \text{H. rewrite sum_x_x. apply H}.
    \times simpl. rewrite \text{H. rewrite sum_x_x. reflexivity}.
    \times simpl. rewrite \text{H. rewrite sum_x_x. reflexivity}.
  + induction t2.
    \times simpl. rewrite \text{H. reflexivity}.
    \times simpl. reflexivity.
    \times simpl. rewrite \text{H. rewrite sum_id. reflexivity}.
    \times simpl. rewrite \text{H. rewrite sum_id. reflexivity}.
    \times simpl. rewrite \text{H. rewrite sum_id. reflexivity}.
  + induction t2.
    \times simpl. rewrite \text{H. rewrite sum_x_x. reflexivity}.
    \times simpl. rewrite \text{H. rewrite sum_id. reflexivity}.

Qed.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
+ induction $t2$.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
+ induction $t2$.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_id}$. reflexivity.
- induction $t$.
+ induction $t2$.
× simpl. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_comm}$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_comm}$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_comm}$. rewrite $\text{sum\_id}$. reflexivity.
+ induction $t2$.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
+ induction $t2$.
× simpl. rewrite $H$. rewrite $\text{sum\_comm}$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
+ induction $t2$.
× simpl. rewrite $H$. rewrite $\text{sum\_comm}$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
× simpl. rewrite $H$. rewrite $\text{sum\_x\_x}$. reflexivity.
+ induction $t2$.
× simpl. rewrite $H$. rewrite $\text{sum\_comm}$. rewrite $\text{sum\_id}$. reflexivity.
× simpl. rewrite \( H \). rewrite \( \text{sum} \_x \_x \). reflexivity.
× simpl. rewrite \( H \). reflexivity.
× simpl. rewrite \( H \). reflexivity.
× simpl. rewrite \( H \). reflexivity.
× simpl. rewrite \(-\) \( H \). reflexivity.

- induction \( t \).
  + induction \( t2 \).
    × simpl. reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \) \( \rightarrow \) \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
  + induction \( t2 \).
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
  + induction \( t2 \).
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
  + induction \( t2 \).
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.

- induction \( t \).
  + induction \( t2 \).
    × simpl. reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
    × simpl. rewrite \( H \). reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite ← H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite ← H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite ← H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite ← H. reflexivity.
  × simpl. rewrite ← H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite ← H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.

- induction t.
  + induction t2.
    × simpl. reflexivity.
    × simpl. rewrite H. reflexivity.
    × simpl. rewrite ← H. reflexivity.
    × simpl. rewrite ← H. reflexivity.
    × simpl. rewrite ← H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite ← H. reflexivity.
+ induction t2.
Lemma pos_left_mul_compat : \A (t t1 t2 : term),
  t == t1 → mult_one_step t1 t2 == mult_one_step t t2.
Proof.
  intros. induction t1.
  - induction t.
    + reflexivity.
    + apply T1_not_equiv_T0 in H. inversion H.
    + induction t2.
      × simpl. rewrite H. rewrite mul_x_x. reflexivity.
      × simpl. rewrite H. rewrite mul_T0_x. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
    + induction t2.
      × simpl. rewrite H. rewrite mul_x_x. reflexivity.
      × simpl. rewrite H. rewrite mul_T0_x. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
    + induction t2.
      × simpl. rewrite H. rewrite mul_x_x. reflexivity.
      × simpl. rewrite H. rewrite mul_T0_x. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
  - induction t.
    + induction t2.
      × simpl. rewrite H. reflexivity.
× simpl. rewrite H. reflexivity.
× simpl. rewrite H. reflexivity.
× simpl. rewrite H. reflexivity.
× simpl. rewrite H. reflexivity.
+ induction t2.
  × simpl. reflexivity.
  × simpl. reflexivity.
  × simpl. reflexivity.
  × simpl. reflexivity.
  × simpl. reflexivity.
+ induction t2.
  × simpl. rewrite H. rewrite mul_comm. rewrite mul_T0x. reflexivity.
  × simpl. rewrite H. rewrite mul_x_x. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
+ induction t2.
  × simpl. rewrite H. rewrite mul_comm. rewrite mul_T0x. reflexivity.
  × simpl. rewrite H. rewrite mul_x_x. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  + induction t2.
    × simpl. rewrite H. rewrite mul_comm. rewrite mul_T0x. reflexivity.
    × simpl. rewrite H. rewrite mul_x_x. reflexivity.
    × simpl. rewrite H. reflexivity.
    × simpl. rewrite H. reflexivity.
    × simpl. rewrite H. reflexivity.
    × simpl. rewrite H. reflexivity.
    × simpl. rewrite H. reflexivity.
  - induction t1.
    + induction t2.
      × simpl. rewrite H. rewrite mul_x_x. reflexivity.
      × simpl. rewrite ← H. rewrite mul_T0x. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
    + induction t2.
      × simpl. rewrite ← H. rewrite mul_comm. rewrite mul_T0x. reflexivity.
      × simpl. rewrite H. rewrite mul_x_x. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
      × simpl. rewrite H. reflexivity.
    + induction t2.
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
+ \text{ induction } t_2.
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
+ \text{ induction } t_2.
- \text{ induction } t.
+ \text{ induction } t_2.
\times \text{simpl. rewrite } H. \text{ rewrite } \text{mul}_x.x. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ rewrite } \text{mul}_{T0}.x. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
+ \text{ induction } t_2.
\times \text{simpl. rewrite } H. \text{ rewrite } \text{mul}_{\text{comm}}. \text{ rewrite } \text{mul}_{T0}.x. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ rewrite } \text{mul}_x.x. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
+ \text{ induction } t_2.
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
+ \text{ induction } t_2.
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
\times \text{simpl. rewrite } H. \text{ reflexivity.}
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
- induction t.
+ induction t2.
  × simpl. rewrite ← H. rewrite mul_x_x. reflexivity.
  × simpl. rewrite ← H. rewrite mul_T0_x. reflexivity.
  × simpl. rewrite ← H. reflexivity.
  × simpl. rewrite ← H. reflexivity.
  × simpl. rewrite ← H. reflexivity.
+ induction t2.
  × simpl. rewrite ← H. rewrite mul_comm. rewrite mul_T0_x. reflexivity.
  × simpl. rewrite H. rewrite mul_x_x. reflexivity.
  × simpl. rewrite ← H. reflexivity.
  × simpl. rewrite ← H. reflexivity.
  × simpl. rewrite ← H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
+ induction t2.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
  × simpl. rewrite H. reflexivity.
Qed.

Lemma pos_right_mul_compat : ∀ (t t1 t2 : term),
  t == t2 → mult_one_step t1 t2 == mult_one_step t1 t.
Proof.
intros. induction t1.
- induction \( t \).
  + induction \( t_2 \).
    \[
    \times\text{ simpl. reflexivity.}
    \times\text{ simpl. rewrite } H\text{. reflexivity.}
    \times\text{ simpl. rewrite } H\text{. rewrite } \text{mul}\_x\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \leftarrow H\text{. rewrite } \text{mul}\_x\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \leftarrow H\text{. rewrite } \text{mul}\_x\_x\text{. reflexivity.}
    \]
  + induction \( t_2 \).
    \[
    \times\text{ simpl. reflexivity.}
    \times\text{ simpl. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \]
  + induction \( t_2 \).
    \[
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \]
  + induction \( t_2 \).
    \[
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \]
  + induction \( t_2 \).
    \[
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \times\text{ simpl. rewrite } \text{mul}\_T0\_x\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
    \]
  - induction \( t \).
    + induction \( t_2 \).
      \[
      \times\text{ simpl. reflexivity.}
      \times\text{ simpl. rewrite } H\text{. reflexivity.}
      \times\text{ simpl. rewrite } \leftarrow H\text{. rewrite } \text{mul}\_\text{comm}\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
      \times\text{ simpl. rewrite } \leftarrow H\text{. rewrite } \text{mul}\_\text{comm}\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
      \times\text{ simpl. rewrite } \leftarrow H\text{. rewrite } \text{mul}\_\text{comm}\text{. rewrite } \text{mul}\_T0\_x\text{. reflexivity.}
      \]
    + induction \( t_2 \).
      \[
      \times\text{ simpl. rewrite } H\text{. reflexivity.}
      \times\text{ simpl. reflexivity.}
      \times\text{ simpl. rewrite } H\text{. rewrite } \text{mul}\_x\_x\text{. reflexivity.}
      \]
\times \text{simpl. rewrite} \ H. \text{rewrite} \ mul_{-x}. \text{reflexivity.}
\times \text{simpl. rewrite} \leftarrow H. \text{rewrite} \ mul_{-x}. \text{reflexivity.}
+ \text{induction} \ t2.
\times \text{simpl. rewrite} \ H. \text{rewrite} \ mul_{\text{comm.}} \text{rewrite} \ mul_{\text{T0-x}}. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{rewrite} \ mul_{-x}. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \leftarrow H. \text{reflexivity.}
+ \text{induction} \ t2.
\times \text{simpl. rewrite} \ H. \text{rewrite} \ mul_{\text{comm.}} \text{rewrite} \ mul_{\text{T0-x}}. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{rewrite} \ mul_{-x}. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \leftarrow H. \text{reflexivity.}
+ \text{induction} \ t2.
\times \text{simpl. rewrite} \ H. \text{rewrite} \ mul_{\text{comm.}} \text{rewrite} \ mul_{\text{T0-x}}. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{rewrite} \ mul_{-x}. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \leftarrow H. \text{reflexivity.}
- \text{induction} \ t.
+ \text{induction} \ t2.
\times \text{simpl. reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \leftarrow H. \text{reflexivity.}
\times \text{simpl. rewrite} \leftarrow H. \text{reflexivity.}
\times \text{simpl. rewrite} \leftarrow H. \text{reflexivity.}
+ \text{induction} \ t2.
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \leftarrow H. \text{reflexivity.}
+ \text{induction} \ t2.
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \leftarrow H. \text{reflexivity.}
+ \text{induction} \ t2.
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\times \text{simpl. rewrite} \ H. \text{reflexivity.}
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( + \text{ induction } t2. \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( - \text{ induction } t. \)
\( + \text{ induction } t2. \)
\( \times \text{simpl. reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( + \text{ induction } t2. \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( + \text{ induction } t2. \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( + \text{ induction } t2. \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( + \text{ induction } t2. \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( + \text{ induction } t2. \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( + \text{ induction } t2. \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
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\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( \times \text{simpl. rewrite } H. \text{ reflexivity.} \)
\( - \text{ induction } t. \)
\( + \text{ induction } t2. \)
× simpl. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $\leftarrow H$. reflexivity.
× simpl. rewrite $\leftarrow H$. reflexivity.
× simpl. rewrite $\leftarrow H$. reflexivity.
+ induction $t_2$.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $\leftarrow H$. reflexivity.
+ induction $t_2$.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $\leftarrow H$. reflexivity.
+ induction $t_2$.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $\leftarrow H$. reflexivity.
+ induction $t_2$.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $H$. reflexivity.
× simpl. rewrite $\leftarrow H$. reflexivity.
Qed.

Being able to simplify a term can be a useful tool. Being able to use the simplified version of the term as the equivalent version of the original term can also be useful since many of our functions simplify the term first.

Lemma simplify_eqv : $\forall (t : \text{term})$, simplify $t == t$.
Proof.
intros. induction $t$.
- simpl. reflexivity.
- simpl. reflexivity.
- simpl. reflexivity.
- simpl. rewrite proof pos_left_sum_compat.
  specialize ($H \, t_2$ (simplify $t_1$) (simplify $t_2$)).
symmetry in \( IHt_1 \). specialize \((H \ IHt_1)\). rewrite \( H \).
pose \( proof \ pos\_right\_sum\_compat \). specialize \((H0 \ (simplify \ t2) \ t1 \ t2)\).
specialize \((H0 \ IHt2)\). symmetry in \( H0 \). rewrite \( H0 \).
induction \( t1 \).
+ induction \( t2 \).
  \times simp. rewrite \( sum\_x\_x \). reflexivity.
  \times simp. rewrite \( sum\_id \). reflexivity.
  \times simp. reflexivity.
  \times simp. reflexivity.
  \times simp. reflexivity.
+ induction \( t2 \).
  \times simp. rewrite \( sum\_id\_sym \). reflexivity.
  \times simp. rewrite \( sum\_x\_x \). reflexivity.
  \times simp. reflexivity.
  \times simp. reflexivity.
  \times simp. reflexivity.
+ simp. reflexivity.
+ simp. reflexivity.
+ simp. reflexivity.

- simp. pose \( proof \ pos\_left\_mul\_compat \).
specialize \((H \ t1 \ (simplify \ t1) \ (simplify \ t2))\).
symmetry in \( IHt1 \). specialize \((H \ IHt1)\). rewrite \( H \).
pose \( proof \ pos\_right\_mul\_compat \). specialize \((H0 \ (simplify \ t2) \ t1 \ t2)\).
specialize \((H0 \ IHt2)\). symmetry in \( H0 \). rewrite \( H0 \).
induction \( t1 \).
+ induction \( t2 \).
  \times simp. rewrite \( mul\_x\_x \). reflexivity.
  \times simp. rewrite \( mul\_T0\_x \). reflexivity.
  \times simp. reflexivity.
  \times simp. reflexivity.
  \times simp. reflexivity.
+ induction \( t2 \).
  \times simp. rewrite \( mul\_T0\_x\_sym \). reflexivity.
  \times simp. rewrite \( mul\_x\_x \). reflexivity.
  \times simp. reflexivity.
  \times simp. reflexivity.
  \times simp. reflexivity.
  \times simp. reflexivity.
+ simp. reflexivity.
+ simp. reflexivity.
+ simp. reflexivity.

Qed.
Chapter 3

Library
B_Unification.lowenheim_formula

Require Export terms.
Require Import List.
Import ListNotations.

3.1 Introduction
In this section we formulate Lowenheim’s algorithm using the data structures and functions defined in the terms library. The final occurring main function, Lowenheim>Main, takes as input a term and produces a substitution that unifies the given term. The resulting substitution is said to be a most general unifier and not a mere substitution, but that statement is proven in the lowenheim_proof file. In this section we focus on the formulation of the algorithm itself, without any proofs about the properties of the formula or the algorithm.

3.2 Lowenheim’s Builder
In this subsection we are implementing the main component of Lowenheim’s algorithm, which is the “builder” of Lowenheim’s substitution for a given term. This implementation strictly follows as close as possible the formal, mathematical format of Lowenheim’s algorithm.

Here is a skeleton function for building a substition on the format \( \sigma(x) := (s + 1) \ast \sigma_1(x) + s \ast \sigma_2(x) \), each variable of a given list of variables, a given term \( s \) and substitutions \( \sigma_1 \) and \( \sigma_2 \). This skeleton function is a more general format of Lowenheim’s builder.

Fixpoint build_on_list_of_vars (list_var : var_set) (s : term) (sig1 : subst)
  (sig2 : subst) : subst :=
  match list_var with
  | [] ⇒ []
\[ v' \Rightarrow (v', (s + T1) \times \text{apply\_subst}\ (\text{VAR}\ v')\ sig1 + \]
\[ s \times \text{apply\_subst}\ (\text{VAR}\ v')\ sig2) \]
\[ :: \text{build\_on\ list\_of\ vars}\ v\ s\ sig1\ sig2 \]

end.

This is the function to build a Lowenhein subsitution for a term \( t \), given the term \( t \)
and a unifier of \( t \), using the previously defined skeleton function. The list of variables is the
variables within \( t \) and the substitutions are the identical substitution and the unifer of the term.
This function will often be referred in the rest of the document as our “Lowenhein builder”
or the “Lowenhein substitution builder”, etc.

**Definition**

\[
\text{build\_lowenheim\ subst}\ (t : \text{term})\ (\tau : \text{subst}) : \text{subst} := \\
\text{build\_on\ list\_of\ vars}\ (\text{term\_unique\ vars}\ t)\ t \\
(\text{build\_id\ subst}\ (\text{term\_unique\ vars}\ t))\ \tau.
\]

### 3.3 Lowenheim’s Algorithm

In this subsection we enhance Lowenheim’s builder to the level of a complete algorithm that
is able to find ground substitutions before feeding them to the main formula to generate a
most general unifier.

#### 3.3.1 Auxiliary Functions and Definitions

This is a function to update a term, after it applies to it a given substitution and simplifies
it.

**Definition**

\[
\text{update\_term}\ (t : \text{term})\ (s' : \text{subst}) : \text{term} := \\
\text{simplify}\ (\text{apply\_subst}\ t\ s').
\]

Here is a function to determine if a term is the ground term \( T_0 \).

**Definition**

\[
\text{term\_is\_T0}\ (t : \text{term}) : \text{bool} := \\
\text{identical}\ t\ T_0.
\]

In this development we have the need to be able to represent both the presence and the
absence of a substitution. In case for example our \text{find\_unifier} function cannot find a unifier
for an input term, we need to be able to return a \text{subst\ nil} type, like a substitution option
that states no substitution was found. We are using the built-in \text{Some} and \text{None} inductive
options (that are used as \text{Some} \( \sigma \) and \text{None}) to represent some substitution and no substitution
respectively. The type of the two above is the inductive \text{option \{A\:type\}} that can be
attached to any type; in our case it is \text{option\ subst}.

Our Lowenheim builder works when we provide an already existing unifier of the input
term \( t \). For our implementation to be complete we need to be able to generate that initial
unifier ourselves. That is why we first need to define a function to find all possible “01”
substitutions (substitutions where each variable gets mapped to \( T_0 \) or \( T_1 \).

**Fixpoint**

\[
\text{all\_01\_substs}\ (\text{vars} : \text{var\_set}) : \text{list\ subst} :=
\]

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match vars with
| [] ⇒ [[]]
| v :: v' ⇒ (map (fun s ⇒ (v, T0) :: s) (all_01_substs v')) ++
  (map (fun s ⇒ (v, T1) :: s) (all_01_substs v'))
end.

Next is a function to find an initial “ground unifier” for our Lowenheim builder function. It finds a substitution with ground terms that makes the given input term equivalent to T0.

Fixpoint find_unifier (t : term) : option subst :=
  find (fun s ⇒ match update_term t s with
    | T0 ⇒ true
    | _ ⇒ false
  end) (all_01_substs (term_unique_vars t)).

3.3.2 Lowenheim’s Main Function

Here is the main Lowenheim’s formula; given a term, produce an MGU (a most general substitution that when applied on the input term, it makes it equivalent to T0), if there is one. Otherwise, return None. This function is often referred in the rest of the document as “Lowenheim Main” function or “Main Lowenheim” function, etc.

Definition Lowenheim_Main (t : term) : option subst :=
  match find_unifier t with
   | Some s ⇒ Some (build_lowenheim subst t s)
   | None ⇒ None
  end.

3.4 Lowenheim’s Functions Testing

In this subsection we explore ways to test the correctness of our Lowenheim’s functions on specific inputs.

Here is a function to test the correctness of the output of the find_unifier helper function defined above. True means expected output was produced, false otherwise.

Definition Test_find_unifier (t : term) : bool :=
  match find_unifier t with
   | Some s ⇒ term_is_T0 (update_term t s)
   | None ⇒ true
  end.
Chapter 4

Library
B_Unification.lowenheim_proof

Require Export lowenheim_formula.
Require Import List.
Import ListNotations.
Require Export EqNat.
Require Import List.
Import ListNotations.
Import Coq.Init.Tactics.
Require Export Classical.Prop.

4.1 Introduction

In this chapter we provide a proof that our Lowenheim_Main function defined in lowenheim_formula provides a unifier that is most general. Our final top level proof (found at the end of this file) proves two statements: 1) If a term is unifiable, then our own defined Lowenheim_Main function produces a most general unifier (mgu). 2) If a term is not unifiable, then our own defined Lowenheim_Main function produces a None substitution. We prove the above statements with a series of proofs and sub-groups of proofs that help us get to the final top-level statements mentioned above.

4.2 Auxillary Declarations and Their Lemmas Useful For the Final Proofs

In this section we provide definitions and proofs of helper functions, propositions, and lemmas that will be later used in other proofs.
This is the definition of an **under_term**. An **under_term** is a proposition, or a relationship between two terms. When a term \( t \) is an **under_term** of a term \( t' \) then each of the unique variables found within \( t \) are also found within the unique variables of \( t' \).

**Definition under_term** \((t : \text{term}) (t' : \text{term}) : \text{Prop} := \)
\[
\forall (x : \text{var}), \ln x (\text{term_unique_vars } t) \to \ln x (\text{term_unique_vars } t').
\]

This is a simple lemma for **under_terms** that states that a term is an **under_term** of itself.

**Lemma under_term_id** : \( \forall (t : \text{term}), \text{under_term } t t \).

**Proof.**
\[
\text{intros. firstorder.}
\]
**Qed.**

This is a lemma to prove the summation distribution property of the function **term_vars**: the **term_vars** of a sum of two terms is equal to the concatenation of the **term_vars** of each individual term of the original sum.

**Lemma term_vars_distr** : \( \forall (t1 t2 : \text{term}), \text{term_vars } (t1 + t2) = \text{term_vars } t1 ++ \text{term_vars } t2 \).

**Proof.**
\[
\text{intros. induction } t2; \text{auto.}
\]
**Qed.**

This is a lemma to prove an intuitive statement: if a variable is within the **term_vars** (list of variables) of a term, then it is also within the **term_vars** of the sum of that term and any other term.

**Lemma tv_h1** : \( \forall (t1 t2 : \text{term}) (x : \text{var}), \ln x (\text{term_vars } t1) \to \ln x (\text{term_vars } (t1 + t2)) \).

**Proof.**
\[
\text{intros. induction } t2.
\]
- **simpl. rewrite app_nil_r. apply } H.**
- **simpl. rewrite app_nil_r. apply } H.**
- **simpl. pose proof in_or_app as } H1. specialize } (H1 var (term_vars t1) [v] x). firstorder.**
- **rewrite term_vars_distr. apply in_or_app. left. apply } H.**
- **rewrite term_vars_distr. apply in_or_app. left. apply } H.**
**Qed.**

This is a lemma similar to the previous one, to prove an intuitive statement: if a variable is within the **term_vars** (list of variables) of a term, then it is also within the **term_vars** of the sum of that term and any other term, but being added from the left side.

**Lemma tv_h2** : \( \forall (t1 t2 : \text{term}) (x : \text{var}), \ln x (\text{term_vars } t2) \to \ln x (\text{term_vars } (t1 + t2)) \).

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This is a helper lemma for the under_term relationship: if the sum of two terms is a subterm of another term $t'$, then the left component of the sum is also a subterm of the other term $t'$.

**Lemma helper_2a** ($\forall (t1 \ t2 \ t' : \text{term})$, $\underdot{\text{under}}(t1 + t2) \ t' \rightarrow \underdot{\text{under}}(t1 \ t')$).

**Proof.**

intros. unfold under_term in *. intros. specialize ($H \ x$).
pose proof in_dup_and_non_dup as $H10$. unfold term_unique_vars.
unfold term_unique_vars in *. pose proof tv_h1 as $H7$. apply $H$.
specialize ($H7 \ t1 \ t2 \ x$). specialize ($H10 \ x (\text{term_vars}(t1 + t2))$).
destruct $H10$. apply $H1$. apply $H7$. pose proof in_dup_and_non_dup as $H10$.
specialize ($H10 \ x (\text{term_vars} t1)$). destruct $H10$. apply $H4$. apply $H0$.
Qed.

This is a helper lemma for the under_term relationship: if the sum of two terms is a subterm of another term $t'$, then the right component of the sum is also a subterm of the other term $t'$.

**Lemma helper_2b** ($\forall (t1 \ t2 \ t' : \text{term})$, $\underdot{\text{under}}(t1 + t2) \ t' \rightarrow \underdot{\text{under}}(t2 \ t')$).

**Proof.**

intros. unfold under_term in *. intros. specialize ($H \ x$).
pose proof in_dup_and_non_dup as $H10$. unfold term_unique_vars.
unfold term_unique_vars in *. pose proof tv_h2 as $H7$. apply $H$.
specialize ($H7 \ t1 \ t2 \ x$). specialize ($H10 \ x (\text{term_vars}(t1 + t2))$).
destruct $H10$. apply $H1$. apply $H7$. pose proof in_dup_and_non_dup as $H10$.
specialize ($H10 \ x (\text{term_vars} t2)$). destruct $H10$. apply $H4$. apply $H0$.
Qed.

This is a helper lemma for lists and their elements: if a variable is a member of a list, then it is equal to the first element of that list or it is a member of the rest of the elements of that list.

**Lemma elt_in_list** ($\forall (x : \text{var}) (a : \text{var}) (l : \text{list} \ \text{var})$, $\ln x (a :: l) \rightarrow \ x = a \lor \ln x l$).

**Proof.**
intros.
pose proof in_inv as H1.
specialize (H1 var a x l H).
destruct H1.
  - left. symmetry in H0. apply H0.
  - right. apply H0.
Qed.

This is a similar lemma to the previous one, for lists and their elements: if a variable is not a member of a list, then it is not equal to the first element of that list and it is not a member of the rest of the elements of that list.

Lemma elt_not_in_list: ∀ (x: var) (a : var) (l : list var),
  ¬ In x (a :: l) →
  x ≠ a ∧ ¬ In x l.
Proof.
  intros.
  pose proof not_in_cons. specialize (H0 var x a l). destruct H0.
  specialize (H0 H). apply H0.
Qed.

This is a lemma for an intuitive statement for the variables of a term: a variable x belongs to the list of unique variables (term_unique_vars) found within the variable-term that is constructed by variable itself VAR x.

Lemma in_list_of_var_term_of_var: ∀ (x: var),
  In x (term_unique_vars (VAR x)).
Proof.
  intros. simpl. left. intuition.
Qed.

This is an intuitive lemma to prove that every element either belongs in any list or does not.

Lemma var_in_out_list: ∀ (x : var) (lvar : list var),
  In x lvar ∨ ¬ In x lvar.
Proof.
  intros.
  pose proof classic as H1. specialize (H1 (In x lvar)). apply H1.
Qed.

4.3 Proof That Lowenheim’s Algorithm (builder) Unifies a Given Term

In this section, we prove that our own defined Lowenheim builder from lowenheim_formula (build_lowenheim_subst), produces a unifier; that is, given unifiable term and one unifier of
the term, it also produces another unifier of this term (and as explained in the terms library, a unifier is a substitution that when applied to term it produces a term equivalent to the ground term $T_0$. The high level proof of this fact is also outlined in the book [Baader and Nipkow, 1998, p. 254-255].

This is a helper lemma for the skeleton function defined in lowenheim_formula: If we apply a substitution on a term-variable VAR $x$, and that substitution is created by the skeleton function $\text{build_on_list_of_vars}$ applied on a single input variable $x$, then the resulting term is equivalent to: the resuting term from applying a substitution on a term-variable VAR $x$, and that substitution being created by the skeleton function $\text{build_on_list_of_vars}$ applied on an input list of variables that contains variable $x$.

**Lemma helper1_easy:** $\forall (x: \text{var}) (\text{lvar : list var}) (\text{sig1 sig2 : subst}) (s : \text{term}),$

\[
\begin{align*}
\text{In } x \text{ lvar} \rightarrow \\
\text{apply_subst} (\text{VAR } x) (\text{build_on_list_of_vars} \text{lvar s sig1 sig2}) == \\
\text{apply_subst} (\text{VAR } x) (\text{build_on_list_of_vars} [x] s sig1 sig2).
\end{align*}
\]

**Proof.**
- intros.
- induction lvar.
- simpl. simpl in H. destruct H.
- apply elt_in_list in H. destruct H.
  + simpl. destruct (beq_nat a x) as ||eqn?:.
    \times apply beq_nat_true in Heqb. destruct (beq_nat x x) as ||eqn?:.
      \times rewrite H. reflexivity.
      \times apply beq_nat_false in Heqb.
        ++ destruct Heqb.
        ++ rewrite Heqb. apply Heqb0.
    \times simpl in IHlvar. apply IHlvar. symmetry in H. rewrite H in Heqb.
      apply beq_nat_false in Heqb. destruct Heqb. intuition.
  + destruct (beq_nat a x) as ||eqn?:.
    \times apply beq_nat_true in Heqb. symmetry in Heqb. rewrite Heqb in IHlvar.
      rewrite Heqb. simpl in IHlvar. simpl. destruct (beq_nat a a) as ||eqn?:.
        \times reflexivity.
        \times apply IHlvar. rewrite Heqb in H. apply H.
    \times apply beq_nat_false in Heqb. simpl. destruct (beq_nat a x) as ||eqn?:.
      \times apply beq_nat_true in Heqb0. rewrite Heqb0 in Heqb. destruct Heqb. intuition.
      \times simpl in IHlvar. apply IHlvar. apply H.
\]

Qed.

This is another helper lemma for the skeleton function $\text{build_on_list_of_vars}$ and it can be rephrased this way: applying two different substitutions on the same term-variable give the same result. One substitution containing only one replacement, and for its own variable. The other substitution contains the previous replacement but also more replacements for other
variables (that are obviously not in the variables of our term-variable). So, the replacements for the extra variables do not affect the application of the substitution - hence the resulting term.

Lemma helper_1: \( \forall (t' \ s : \text{term}) \ (v : \text{var}) \ (\text{sig1} \ \text{sig2} : \text{subst}), \)

\[ \under_{\text{term}} (\text{VAR} \ v) \ t' \rightarrow \]

\[ \text{apply}_{\text{subst}} (\text{VAR} \ v) \]

\[ (\text{build}_\text{on}_{\text{list}_\text{of-vars}} (\text{term}_{\text{unique-vars}} t') \ s \ \text{sig1} \ \text{sig2}) \]

\[ = \]

\[ \text{apply}_{\text{subst}} (\text{VAR} \ v) \]

\[ (\text{build}_\text{on}_{\text{list}_\text{of-vars}} (\text{term}_{\text{unique-vars}} (\text{VAR} v)) \ s \ \text{sig1} \ \text{sig2}). \]

Proof.

intros. unfold under_term in \( H \). specialize \( (H \ v) \).

pose proof in list_of_var_term_of_var as \( H3 \). specialize \( (H3 \ v) \).

specialize \( (H2 \ H3) \). pose proof helper1_easy as \( H2 \).

specialize \( (H2 \ v \ (\text{term}_{\text{unique-vars}} t') \ \text{sig1} \ \text{sig2} \ s) \). apply \( H2 \). apply \( H \).

Qed.

Lemma 10.4.5 from 'Term Rewriting and All That' book on page 254-255. This a very significant lemma used later for the proof that our Lowenheim builder function (not the Main function, but the builder function), gives a unifier (not necessarily an mgu, which would be a next step of the proof). It states that if a term \( t \) is an \( \under_{\text{term}} \) of another term \( t' \), then applying a substitution—a substitution created by giving the list of variables of term \( t' \) on the skeleton function \( \text{build}_\text{list}_\text{of-vars} \), on the term \( t \), a term that has the same format:

\[ (s + 1) * \sigma_1(t) + s * \sigma_2(t) \]

as each replacement of each variable on any substitution created by skeleton function:

\[ (s + 1) * \sigma_1(x) + s * \sigma_2(x). \]

Lemma subs_distr_vars_ver2 : \( \forall (t \ t' \ s : \text{term}) \ (\text{sig1} \ \text{sig2} : \text{subst}), \)

\[ \under_{\text{term}} t \ t' \rightarrow \]

\[ \text{apply}_{\text{subst}} t (\text{build}_\text{on}_{\text{list}_\text{of-vars}} (\text{term}_{\text{unique-vars}} t') \ s \ \text{sig1} \ \text{sig2}) \]

\[ = \]

\[ (s + T1) \times \text{apply}_{\text{subst}} t \ \text{sig1} + s \times \text{apply}_{\text{subst}} t \ \text{sig2}. \]

Proof.

intros. generalize dependent \( t' \). induction \( t \).

- intros \( t' \). repeat rewrite ground_term_cannot_subst.
  + rewrite mul_comm with \( (x := s + T1). \) rewrite distr.
    repeat rewrite mul_T0_x. rewrite mul_comm with \( (x := s). \)
    rewrite mul_T0_x. repeat rewrite sum_x_x. reflexivity.
  + unfold ground_term. reflexivity.
  + unfold ground_term. reflexivity.
  + unfold ground_term. reflexivity.

- intros \( t' \). repeat rewrite ground_term_cannot_subst.
  + rewrite mul_comm with \( (x := s + T1). \) rewrite mul_id.
    rewrite mul_comm with \( (x := s). \) rewrite mul_id.
    rewrite sum_comm with \( (x := s). \)
    repeat rewrite sum_assoc. rewrite sum_x_x.
    rewrite sum_comm with \( (x := T1). \) rewrite sum_id. reflexivity.
+ unfold ground_term. reflexivity.
+ unfold ground_term. reflexivity.
+ unfold ground_term. reflexivity.

- intros. rewrite helper_1.
  + unfold term_unique_vars. unfold term_vars. unfold var_set_create_unique.
  unfold var_set_includes_var. unfold build_on_list_of_vars.
  rewrite var_subst. reflexivity.
+ apply H.

- intros. specialize (IHt1 t'). specialize (IHt2 t').
  repeat rewrite subst_sum_distr_opp. rewrite IHt1. rewrite IHt2.
  + rewrite distr. rewrite distr. repeat rewrite sum_assoc.
    rewrite sum_comm with (x := (s + T1) \times apply_subst t2 sig1)
    (y := s \times apply_subst t1 sig2 + s \times apply_subst t2 sig2).
    repeat rewrite sum_assoc.
    rewrite sum_comm with (x := s \times apply_subst t2 sig2)
    (y := (s + T1) \times apply_subst t2 sig1).
    repeat rewrite sum_assoc. reflexivity.
+ pose helper_2b as H2. specialize (H2 t1 t2 t'). apply H2. apply H.
+ pose helper_2a as H2. specialize (H2 t1 t2 t'). apply H2. apply H.

- intros. specialize (IHt1 t'). specialize (IHt2 t').
  repeat rewrite subst_mul_distr_opp. rewrite IHt1. rewrite IHt2.
  + rewrite distr.
    rewrite mul_comm with (y := (s + T1) \times apply_subst t2 sig1).
    rewrite distr. rewrite mul_comm with (y := s \times apply_subst t2 sig2).
    rewrite distr. repeat rewrite mul_assoc.
    repeat rewrite mul_comm with (x := apply_subst t2 sig1).
    repeat rewrite mul_assoc.
    rewrite mul_assoc_opp with (x := s + T1) (y := s + T1). rewrite mul_x_x.
    rewrite mul_assoc_opp with (x := s + T1) (y := s).
    rewrite mul_comm with (x := s + T1) (y := s). rewrite distr.
    rewrite mul_x_x. rewrite mul_id_sym. rewrite mul_x_x. rewrite mul_T0_x.
    repeat rewrite mul_assoc.
    rewrite mul_comm with (x := apply_subst t2 sig2).
    repeat rewrite mul_assoc.
    rewrite mul_assoc_opp with (x := s) (y := s + T1). rewrite distr.
    rewrite mul_x_x. rewrite mul_id_sym. rewrite mul_x_x. rewrite mul_T0_x.
    repeat rewrite sum_assoc. rewrite sum_assoc_opp with (x := T0) (y := T0).
    rewrite sum_x_x. rewrite sum_id. repeat rewrite mul_assoc.
    rewrite mul_comm with (x := apply_subst t2 sig2)
    (y := s \times apply_subst t1 sig2).
    repeat rewrite mul_assoc. rewrite mul_assoc_opp with (x := s).
    rewrite mul_x_x. reflexivity.
This is an intermediate lemma occurring by the previous lemma 10.4.5. Utilizing lemma 10.4.5 and also using two substitutions for the skeleton function \( build_{on\_list\_vars} \) gives a substitution the unifies the term; the two substitutions being a known unifier of the term and the identity substitution.

**Lemma specific_sigmoids_unify:** \( \forall (t : \text{term}) (\tau : \text{subst}), \)

\[
\text{unifier \ t \ \tau} \rightarrow 
\text{apply}_{\text{subst}} \ t \ (\text{build}_{\text{on\_list\_of\_vars}} (\text{term\_unique\_vars} \ t) \ t \ (\text{build}_{\text{id\_subst}} (\text{term\_unique\_vars} \ t))) \ (\tau) == T0.
\]

**Proof.**

\begin{itemize}
  \item \text{intros.}
  \item \text{rewrite subs_{distr\_vars_{ver2}}.}
  \item \text{rewrite id_{subst}. rewrite mul_{comm} with (x := t + T1). rewrite distr.}
  \item \text{rewrite mul_{x\_x}. rewrite mul_{id\_sym}. rewrite sum_{x\_x}. rewrite sum_{id}.}
  \item \text{unfold unifier in H. rewrite H. rewrite mul_{T0\_x\_sym}. reflexivity.}
  \item \text{- apply under_{term\_id}.}
\end{itemize}

Qed.

This is the resulting lemma from this subsection: Our Lowenheim’s substitution builder produces a unifier for an input term; namely, a substitution that unifies the term, given that term is unifiable and we know an already existing unifier \( \tau \).

**Lemma lowenheim_unifies:** \( \forall (t : \text{term}) (\tau : \text{subst}), \)

\[
\text{unifier \ t \ \tau} \rightarrow 
\text{apply}_{\text{subst}} \ t \ (\text{build}_{\text{lowenheim\_subst}} \ t \ \tau) == T0.
\]

**Proof.**

\begin{itemize}
  \item \text{intros. unfold build_{lowenheim\_subst}. apply specific_sigmoids_unify. apply H.}
\end{itemize}

Qed.

### 4.4 Proof That Lowenheim’s Algorithm (Builder) Produces a Most General Unifier

In the previous section we proved that our Lowenheim builder produces a unifier, if we already know an existing unifier of the term. In this section we prove that this unifier is also a most general unifier.
4.4.1 Proof That Lowenheim’s Algorithm (Builder) Produces a Reproductive Unifier

In this subsection we will prove that our Lowenheim builder gives a unifier that is reproductive; this will help us in the proof that the resulting unifier is an mgu, since a reproductive unifier is a “stronger” property than an mgu. The high level proof of this fact is also outlined in the book [Baader and Nipkow, 1998, p. 255]

This is a lemma for an intuitive statement for the skeleton function build on list vars: if a variable \( x \) is in a list \( l \), and we apply a substitution created by the build on list vars function given input list \( l \), on the term-variable \( \text{VAR} \ x \), then we get the replacement for that particular variable that was contained in the original substitution. So basically if build on list of vars is applied on a list of variables \( l \ (x_1, x_2, x_3, ..., x_n) \), then the resulting substitution is in the format \( x_i \mapsto (s + 1) \ast \sigma_1 (x_i) + s \ast \sigma_2 (x_i) \) for each \( x_i \). If we apply that substitution on the term-variable \( x_1 \), we will get the initial format of the replacement: \((s + 1) \ast \sigma_1(x_1) + s \ast \sigma_2(x_1)\). It can be thought as “reverseapplication” of the skeleton function.

**Lemma lowenheim_rephrase1_easy**: \( \forall \ (l : \text{list} \ \text{var}) \ (x : \text{var}) \ (\text{sig1} \ \text{sig2} : \text{subst}) \ (s : \text{term}), \)

\[
\text{In} \ x \ l \rightarrow \\
\text{apply subst} \ (\text{VAR} \ x) \ (\text{build} \text{on} \text{list} \text{of} \text{vars} \ l \ s \ \text{sig1} \ \text{sig2}) = \\
(s + T1) \times \text{apply subst} \ (\text{VAR} \ x) \ \text{sig1} + s \times \text{apply subst} \ (\text{VAR} \ x) \ \text{sig2}.
\]

**Proof**.
- intros.
- induction \( l \).
  - simpl. unfold \( \text{In} \) in \( H \). destruct \( H \).
    + simpl. destruct (beq_nat a x) as \[eqn:?.\]
      \times rewrite \( H \). reflexivity.
      \times pose proof \( \text{beq_nat_false} \) as \( H2 \). specialize (\( H2 \ a \ x \)).
        specialize (\( H2 \ \text{Heqb} \)). intuition. symmetry in \( H \). specialize (\( H2 \ H \)).
        inversion \( H2 \).
    + simpl. destruct (beq_nat a x) as \[eqn:?.\]
      \times symmetry in \( \text{Heqb} \). pose proof \( \text{beq_nat_eq} \) as \( H2 \). specialize (\( H2 \ a \ x \)).
        specialize (\( H2 \ \text{Heqb} \)). rewrite \( H2 \). reflexivity.
      \times apply \( \text{IH}l \). apply \( H \).

Qed.

This is a helper lemma for an intuitive statement: if a variable \( x \) is found in a list of variables \( l \), then applying the substitution created by the build_id_subst function given input list \( l \), on the term-variable \( \text{VAR} \ x \), we will get the same \( \text{VAR} \ x \) back.

**Lemma helper_3a**: \( \forall \ (x : \text{var}) \ (l : \text{list} \ \text{var}), \)

\[
\text{In} \ x \ l \rightarrow \\
\text{apply subst} \ (\text{VAR} \ x) \ (\text{build_id_subst} \ l) = \text{VAR} \ x.
\]

**Proof**.

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intros. induction l.
- unfold build_id_subst. simpl. reflexivity.
- apply elt_in_list in H. destruct H.
  + simpl. destruct (beq_nat a x) as []eqn:?.
    \times rewrite H. reflexivity.
    \times pose proof beq_nat_false as H2. specialize (H2 a x).
    specialize (H2 Heqb). intuition. symmetry in H. specialize (H2 H).
    inversion H2.
  + simpl. destruct (beq_nat a x) as []eqn:?.
    \times symmetry in Heqb. pose proof beq_nat_eq as H2. specialize (H2 a x).
    specialize (H2 Heqb). rewrite H2. reflexivity.
    \times apply IHl. apply H.
Qed.

This is a lemma for an intuitive statement for the Lowenheim builder, very similar to lemma lowenheim_rephrase1_easy: applying Lowenheim’s substitution given an input term \(t\), on any term-variable of the term \(t\), gives us the initial format of the replacement for that variable (Lowenheim’s reverse application).

**Lemma** lowenheim_rephrase1 : \(\forall (t : \text{term}) (\tau : \text{subst}) (x : \text{var}),\)
\begin{align*}
\text{unifier } t \tau & \rightarrow \\
\text{In } x (\text{term_unique_vars } t) & \rightarrow \\
\text{apply_subst } (\text{VAR } x) (\text{build_lowenheim_subst } t \tau) & = \\
(t + \text{T1}) \times (\text{VAR } x) + t \times \text{apply_subst } (\text{VAR } x) \tau.
\end{align*}

**Proof.**
- intros.
- unfold build_lowenheim_subst. pose proof lowenheim_rephrase1_easy as H1.
- specialize (H1 (term_unique_vars t) x
  (build_id_subst (term_unique_vars t))) tau t).
- rewrite helper_3a in H1.
- apply H1. apply H0.
- apply H0.
Qed.

This is a lemma for an intuitive statement for the skeleton function build_on_list_vars that resembles a lot of lowenheim_rephrase1_easy: if a variable \(x\) is not in a list \(l\), and we apply a substitution created by the build_on_list_vars function given input list \(l\), on the term-variable \(\text{VAR } x\), then we get the term-variable \(\text{VAR } x\) back; that is expected since the replacements in the substitution should not contain any entry with variable \(x\).

**Lemma** lowenheim_rephrase2_easy : \(\forall (l : \text{list var}) (x : \text{var}) (s1 s2 : \text{subst}) (s : \text{term}),\)
\begin{align*}
\neg (\text{In } x l) & \rightarrow \\
\text{apply_subst } (\text{VAR } x) (\text{build_on_list_of_vars } l s s1 s2) & = \\
\text{VAR } x.
\end{align*}
Proof.
  intros. unfold not in H.
  induction l.
  - simpl. reflexivity.
  - simpl. pose proof elt_not_in_list as H2. specialize (H2 x a l).
    unfold not in H2. specialize (H2 H). destruct H2.
    destruct (beq_nat a x) as ![eqn:?].
    + symmetry in Heqb. apply beq_nat_eq in Heqb. symmetry in Heqb.
      specialize (H0 Heqb). destruct H0.
    + simpl in IHl. apply IHl. apply H1.
Qed.

This is a lemma for an intuitive statement for the Lowenheim builder, very similar to lemma lowenheim_rephrase2_easy and lowenheim_rephrase1: applying Lowenheim’s substitution given an input term t, on any term-variable not of the ones of term t, gives us back the same term-variable.

Lemma lowenheim_rephrase2 : \( \forall (t : \text{term}) (\tau : \text{subst}) (x : \text{var}), \)
  unifier \( t \tau \rightarrow \)
  \( \neg (\ln x (\text{term} \text{unique vars } t)) \rightarrow \)
  apply_subst (VAR x) (build_lowenheim_subst t \tau) \( = \)
  VAR x.
Proof.
  intros. unfold build_lowenheim_subst.
  pose proof lowenheim_rephrase2_easy as H2.
  specialize (H2 \( \text{term} \text{unique vars } t \)) \( x \)
    (build_id_subst \( \text{term} \text{unique vars } t \)) \( \tau t \).
  specialize (H2 H0). apply H2.
Qed.

This is the resulting lemma of the section: our Lowenheim builder \( \text{build_lowenheim_subst} \) gives a reproductive unifier.

Lemma lowenheim_reproductive: \( \forall (t : \text{term}) (\tau : \text{subst}), \)
  unifier \( t \tau \rightarrow \)
  reproductive_unifier \( t \) (build_lowenheim_subst t \tau).
Proof.
  intros. unfold reproductive_unifier. intros.
  pose proof \( \text{var in out list} \). split.
  - apply lowenheim_unifies. apply H.
  - intros. specialize (H0 \( x (\text{term} \text{unique vars } t) \)). destruct H0.
    + rewrite lowenheim_rephrase1.
      \times rewrite subst_sum_distr_opp. rewrite subst_mul_distr_opp.
        rewrite subst_mul_distr_opp. unfold unifier in H1. rewrite H1.
        rewrite \( \text{mul T0} x \). rewrite subst_sum_distr_opp. rewrite H1.

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rewrite ground_term_cannot_subst.
- rewrite sum_id. rewrite mul_id. rewrite sum_comm. rewrite sum_id.
  reflexivity.
- unfold ground_term. intuition.
  × apply H.
  × apply H0.
+ rewrite lowenheim_rephrase2.
  × reflexivity.
  × apply H.
  × apply H0.
Qed.

4.4.2 Proof That Lowenheim’s Algorithm (Builder) Produces a Most General

Unifier

In this subsection we will prove that our Lowenheim builder gives a unifier that is most general; this will help us a lot in the top-level proof that the Main_Lownheim function gives an mgu. We will use the fact that we proved in the term.v file/library that any reproductive unifier is also a most general unifier, and the fact of the just preceding subsection that lowenheim’s builder produces a reproductive unifier.

Here is the subsection’s resulting lemma. Given a unifiable term \( t \), a unifier of \( t \), then our Lowenheim builder (build_lowenheim subst) gives a most general unifier (mgu).

**Lemma** lowenheim_most_general_unifier: \( \forall (t : \text{term}) (\tau : \text{subst}), \)

\[ \text{unifier } t \tau \rightarrow \text{most_general_unifier } t \text{ (build_lowenheim_subst } t \tau). \]

**Proof.**

intros. apply reproductive_is_mgu. apply lowenheim_reproductive. apply H.
Qed.

4.5 Proof of Correctness of Lowenheim_Main

In the previous section, we proved that our “lowenheim builder” produces an mgu of an input term \( t \), given an existing unifier of \( t \). Even though what was proven in the previous section was the bulk of the core proof which was also presented in the book in a higher level, it did not incorporate many crucial elements. In this section we provide a proof of correctness of our Lowenheim_Main function, basically incorporating more elements in the final proof, like proving correctness in the case that the term \( t \) is not unifiable (which is not covered in the previous section), include in the proofs our find_unifier function that finds an initial “01” substitution to feed the “lowenheim builder”, and more. As it follows from the above, the proof of correctness of the Lowenheim_Main function, uses the proof of the previous section (that the “lowenheim builder” produces) as a building block.
In this section we prove that our own defined Lowenheim function satisfies its two main requirements: 1) If a term is unifiable, then $\text{Lowenheim\_Main}$ function produces a most general unifier (mgu). 2) If a term is not unifiable, then $\text{Lowenheim\_Main}$ function produces a $\text{None}$ substitution. The final top-level proof is at the end of this section. To get there, we prove a series of intermediate lemmas that are needed for the final proof.

### 4.5.1 General Proof Utilities

In this section we provide helper “utility” lemmas and functions that are used in the proofs of intermediate lemmas that are in turn used in the final proof.

This is a function that converts an $\text{option subst}$ to a $\text{subst}$. It is designed to be used mainly for $\text{option subst}$s that are $\text{Some } \sigma$. If the input $\text{option subst}$ is not $\text{Some}$ and is $\text{None}$ then it returns the nil substitution, but that case should not normally be considered. This function is useful because many functions and lemmas are defined for the substitution type not the option substitution type.

**Definition convert_to_subst** ($so : \text{option subst}$) : $\text{subst} :=$

\[
\begin{align*}
\text{match } so \text{ with} \\
| \text{Some } s \Rightarrow s \\
| \text{None } \Rightarrow []
\end{align*}
\]

This is an intuitive helper lemma that proves that if an empty substitution is applied on any term $t$, then the resulting term is the same input term $t$.

**Lemma empty_subst_on_term:** $\forall (t : \text{term}),$

\[
\text{apply_subst } t \emptyset == t.
\]

**Proof.**

- intros. induction $t$.
- reflexivity.
- simpl. reflexivity.
- simpl. reflexivity.
- simpl. rewrite $IHt1$. rewrite $IHt2$. reflexivity.
- simpl. rewrite $IHt1$. rewrite $IHt2$. reflexivity.

Qed.

This another intuitive helper lemma that states that if the empty substitution is applied on any term $t$, and the resulting term is equivalent to the ground term $T0$, then the input term $t$ must be equivalent to the ground term $T0$.

**Lemma app_subst_T0:** $\forall (t : \text{term}),$

\[
\text{apply_subst } t \emptyset == T0 \rightarrow t == T0.
\]

**Proof.**

- intros. rewrite empty_subst_on_term in $H$. apply $H$.

Qed.
This is another intuitive lemma that uses classical logic for its proof. It states that any term $t$, can be equivalent to the ground term $T0$ or it cannot be equivalent to it.

Lemma $T0\_or\_not\_T0$: $\forall (t : \text{term}),$
  $t == T0 \lor \neg t == T0$.
Proof.
  intros. pose proof classic. specialize ($H (t == T0)$). apply $H$.
Qed.

This is another intuitive helper lemma that states: if applying a substitution $\sigma$ on a term $t$ gives a term equivalent to $T0$ then there exists a substitution that applying it to term $t$ gives a term equivalent to $T0$ (which is obvious since we already know $\sigma$ exists for that task).

Lemma $\exists\_\text{subst}$: $\forall (t : \text{term}) (\sigma : \text{subst}),$
  apply $\text{subst} t \sigma == T0 \rightarrow \exists s, \text{apply subst} t s == T0$.
Proof.
  intros. $\exists sig$. apply $H$.
Qed.

Lemma $t\_id\_eqv$ : $\forall (t : \text{term}),$
  $t == t$.
Proof.
  intros. reflexivity.
Qed.

This a helper lemma that states: if two options subst (specifically Some) are equal then the substitutions contained within the option subst are also equal.

Lemma $\text{eq}\_\text{some eq subst}$ ($s1 s2 : \text{subst}$) :
  Some $s1 = \text{Some} s2 \rightarrow s1 = s2$.
Proof.
  intros. congruence.
Qed.

This a helper lemma that states: if the $\text{find unifier}$ function (the one that tries to find a ground unifier for term $t$) does not find a unifier (returns None) for an input term $t$ then it not True (true not in “boolean format” but as a proposition) that the $\text{find unifier}$ function produces a Some subst. This lemma and the following ones that are similar, are very useful for the intermediate proofs because we are able to convert a proposition about the return type of the $\text{find unifier}$ function to an equivalent one, e.g. from None subst to Some subst and vice versa.

Lemma $\text{None is not Some}$ ($t : \text{term}$):
  find $\text{unifier} t = \text{None} \rightarrow$
  $\forall (sig : \text{subst}), \neg \text{find unifier} t = \text{Some} sig$.
Proof.
  intros. congruence.
Qed.

This a helper lemma similar to the previous one that states: if the find_unifier function (the one that tries to find a ground unifier for term t) finds a unifier (returns Some σ) for an input term t then it is not True (true not in “boolean format” but as a proposition) that the find_unifier function produces a None subst.

Lemma Some_is_not_None (sig: subst) (t: term):
find_unifier t = Some sig → ¬ find_unifier t = None.
Proof.
intros. congruence.
Qed.

This a helper lemma similar to the previous ones that states: if the find_unifier function (the one that tries to find a ground unifier for term t) does not find a unifier that returns None for an input term t then it is True (true not in “boolean format” but as a proposition) that the find_unifier function produces a Some subst.

Lemma not_None_is_Some (t: term):
¬ find_unifier t = None →
∃ sig : subst, find_unifier t = Some sig.
Proof.
intros H.
destruct (find_unifier t) as [ti |].
- ∃ ti. firstorder.
- congruence.
Qed.

This is an intutitive helper lemma that uses classical logic to prove the validity of an alternate version of the contrapositive proposition: if p then q implies if not q then not p, but with each entity (proposition q and p) negated.

Lemma contrapositive_opposite : ∀ p q,
(¬p → ¬q) →
q → p.
Proof.
intros. apply NNPP. firstorder.
Qed.

This is an intutitive helper lemma that uses classical logic to prove the validity of the contrapositive proposition: if p then q implies not q then not p.

Lemma contrapositive : ∀ (p q : Prop),
(p → q) →
(¬q → ¬p).
Proof.
intros. firstorder.
Qed.
The following five lemmas are also helper lemmas.

Lemma None_not_Some \{ T U: Type \} (f : U \to \texttt{option} T) (x: U): 
\( (f \ x) = \texttt{None} \to (\forall (t: T), \neg (f \ x) = \texttt{Some} \ t) \).
Proof.
  intros.
  congruence.
Qed.

Lemma Some_not_None \{ T U: Type \} (f : U \to \texttt{option} T) (x: U) (t: T): 
\( (f \ x) = \texttt{Some} \ t \to \neg (f \ x = \texttt{None}) \).
Proof.
  intros.
  congruence.
Qed.

Lemma not_None_Some \{ T U: Type \} (f : U \to \texttt{option} T) (x: U) :
\( \neg (f \ x = \texttt{None}) \to \exists t : T, f \ x = \texttt{Some} \ t \).
Proof.
  intros.
  destruct \( f \ x \) as \[ t \ | \].
  \- \exists t; easy.
  \- congruence.
Qed.

Lemma not_Some_None \{ T U: Type \} (f : U \to \texttt{option} T) (x: U) :
\( (\neg \exists t : T, f \ x = \texttt{Some} \ t) \to f \ x = \texttt{None} \).
Proof.
  apply contrapositive_opposite.
  intros.
  apply not_None_Some in \( H \).
  tauto.
Qed.

Lemma existsb_find \{ T: Type \} (f: T \to \texttt{bool}) (l : list T) :
\( \texttt{existsb} \ f \ l = \texttt{true} \to \exists (a: T), \texttt{find} \ f \ l = \texttt{Some} \ a \).
Proof.
  intros.
  apply NNPP.
  intros \( H1 \).
  apply not_Some_None in \( H1 \).
  assert \( \( S1:= \texttt{find\_none} \ f \ l \) \).
  assert \( \( S2:= S1 \ H1 \) \).
  assert \( \( S3:= \texttt{existsb\_exists} \ f \ l \) \).
  destruct \( S3 \) as \[ S31 S32 \].
assert (S4 := S31 H).
destruct S4 as [t S41]. destruct S41 as [S411 S412].
assert (S21 := S2 t S411).
rewrite S412 in S21.
congruence.
Qed.

4.5.2 Utilities and Admitted Lemmas Used in the Proof of unif_some_subst

In this subsection we have collected and put together all the functions and lemmas that are used to prove the unif_some_subst lemma that is used in the following intermediate lemmas section, and specifically in the “unifiable t” case. The higher-level lemma we aim to prove using this section is a seemingly simple, but in reality very complex lemma; the unif_some_subst lemma states that if there is any unifier sig1 for a term t then there exists a unifier sig2 which is returned by our find_unifier function.

Due to lack of time, our team did not manage to prove these last five lower-level lemmas used in the proof of unif_some_subst, and since they are all used only in the proof of that lemma, we decided to put them together here in this subsection, along with everything else that is used for the proof of that lemma.

Utilities Used in This Subsection

In this sub-chapter we are declaring two utility lemmas used in next sub-chapter by the lower-level lemmas of this proof.

This is a lemma that states that sequentially applying two substitutions on a term produces the same term as applying the composed substitutions on the term.

Lemma subst-compose_eqv : \forall (t : term) (sig1 : subst) (sig2 : subst),
apply_subst t (subst-compose sig1 sig2) ==
apply_subst (apply_subst t sig2) sig1.
Proof.
intros. induction t.
- simpl. reflexivity.
- simpl. reflexivity.
- simpl. induction sig2.
  + simpl. reflexivity.
  + simpl. induction sig1.
Admitted.

This is an intuitive lemma that states when a term is equivalent to T0 and it is also a ground term then simplifying it gives a term exactly equal to T0. This intuitively follows from the fact that since t is a ground term then all its terms are either T0 or T1 and since it is equivalent to T0, simplifying it will also give a single final ground term T0.

Lemma simplify_eq_T0 : \forall (t : term),

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\[ t = T0 \land (\text{is\_ground\_term} \ t) = \text{true} \rightarrow \text{simplify} \ t = T0. \]

\textbf{Proof}.
- intros. destruct \( H \). induction \( t \).
- reflexivity.
- simpl in \( H \). apply \( T1\_\text{not\_equiv\_T0} \) in \( H \). destruct \( H \).
- unfold simplify. simpl in \( H0 \). inversion \( H0 \).
- simpl. simpl in \( H0 \). apply andb\_prop in \( H0 \). destruct \( H0 \).
\textit{Admitted.}

\textbf{Lower Level Lemmas Leading Up to the Proof of unif\_some\_subst}

In this sub-chapter we are providing the most important lower-level lemmas leading up to the proof of the \texttt{unif\_some\_subst} lemma.

To accomplish the goal of providing the infrastructure to prove the \texttt{unif\_some\_subst} lemma, we are defining a number of functions and lemmas that are used in the proof of the \texttt{unif\_some\_subst}. We are focusing on connecting the concept of a “01” substitution with any given substitution. We are attempting to create a “01” substitution given any input substitution, and then prove facts about the new “01” substitution.

The basic outline of the proof is as follows: From any unifier \( sig1 \) of term \( t \) we can create a “01” unifier \( sig2 \), as the one defined in the \texttt{terms.v} library. Since our \texttt{find\_unifier} function looks for all “01” substitutions to find a unifier, and we already know there exists at least one unifier \( sig2 \) that will be returned from the \texttt{find\_unifier} function.

As it follows, the lower part of this proof is (1) creating a “01” unifier \( sig2 \) from a given given unifier of \( t \) \( sig1 \) (2) proving that the new “unifier” is actually a unifier and proving that it is actually a “01” substitution.

All the following functions are defined in order to create a final function that is able to produce a “01” unifier \( sig2 \) given a unifier \( sig1 \). The way \( sig2 \) is created is by composing two substitutions, \( sig1 \) and \( sig1b \) so that \( sig1b \) and \( sig1 \) are composed to give us \( sig2 \). The idea behind the \( sig1b \) substitution creation function is that it takes all the replacements of the given unifier \( sig1 \) and it does the following to each replacement of the \( sig1 \) substitution (let us represent each replacement as \((v,t)\)): if the second part of the replacement is a \texttt{ground\_term} , then we create a new replacement that is \((v, t')\), where \( t' \) is the simplified initial second part \( t \). If the second part of the relacement is not a \texttt{ground\_term}, then we create a list of new replacements where each new replacement is one variable found in the initial \( v \) mapped to the ground term \( T0 \).

So for example suppose our \( sig1 \) included the replacement \((v, x + y + T1)\); we then create the new replacements \((x, T0)\) and \((y, T0)\).

The total final list of all the new replacements is the substitution \( sig1b \). As we want to cover for all edge cases, we have created a slightly enhanced version of this \( sig1b \) creation function. Instead of working on the initial \( sig1 \) unifier, we are enhancing \( sig1 \) by adding a list of identity replacement for all variables of term \( t \) that are not in \( sig1 \). For example for the term \( x \ast y \) and the unifier \((x,T0)\), we first enhance \( sig1 \) by making it \((x,T0), (y,y)\) and then we create \( sig1b \) based on the enhanced \( sig1 \).
After composing \textit{sig1b} with \textit{sig1} we get \textit{sig2} which intuitively is a “01” unifier. But it is harder to prove than claim it of course, that is why we have put all the admitted lemmas of this proofs here.

This is a function to build a \texttt{T0 subst}, a substitution that maps each variable to \texttt{T0}, given an input list of variables.

\begin{verbatim}
Definition build_T0 subst (lvar : list var) : subst :=
  map (fun v ⇒ (v, T0)) lvar.
\end{verbatim}

Next is a function to build a \texttt{T0 subst}, given an input term \texttt{t}.

\begin{verbatim}
Definition build_T0 subst_from_t (t : term) : subst :=
  build_T0 subst (term_unique_vars t).
\end{verbatim}

With the following four helper functions, we are trying to create a final function that does the following: 1) Given any substitution, it produces a “01” substitution building off the given substitution. 2) It does that by composing two substitutions \texttt{s1} and \texttt{s1b} into a new one, \texttt{s2}. 3) It creates \texttt{s1b} from \texttt{s1}. \texttt{s1b} is a “01” unifier and so is \texttt{s2}.

Here is the function to create the \texttt{s1b} “01” substitution, by mapping all the second parts of each replacement of the substitution using the following rules: 1) All the variables of non-ground terms are mapped to \texttt{T0} and all ground terms are mapped to their simplified “01” version. Therefore the substitution occurring from this function is a “01” substitution, intuitively.

\begin{verbatim}
Fixpoint make_unif subst (tau : subst) : subst :=
  match tau with
  | [] ⇒ []
  | (first , second) :: rest' ⇒
    if is_ground_term second
    then (first , simplify second) :: (make_unif subst rest')
    else (build_T0 subst_from_t second) ++ (make_unif subst rest')
  end.
\end{verbatim}

This function creates a list of identity replacements, for all the variables of the \texttt{lvar} list input that are not in \texttt{lvar_s} list input. The \texttt{var_s} list input is supposedly the list with the variables of a substitution and we are trying eventually to augment the substitution with and identity substitution.

\begin{verbatim}
Fixpoint augment_with_id (lvar_s : list var) (lvar : list var) : subst :=
  match lvar with
  | [] ⇒ []
  | v :: v' ⇒
    if var_set_includes_var v lvar_s
    then augment_with_id lvar_s v'
    else (v, VAR v) :: (augment_with_id lvar_s v')
  end.
\end{verbatim}

This function adds the identity substitution (or list of identity replacements in this case) to the input substitution.
Definition add_id_subst \((t : \text{term}) (\tau : \text{subst}) : \text{subst} :=\) 
\[
\text{augment_with_id (\text{subst_domain} \tau) (\text{term_unique_vars} t) ++ \tau}
\]

This is the resulting function that given any substitution for a term, produces a “01” substitution. Even though this function is not directly called by name, its implementation is directly used. So whenever in the future comments there is a reference to a \texttt{convert_to_01_subst}, what is meant is essentially the composition of the \texttt{make_unif_subst} substitution and the input substitution \(\tau\) - or the resulting substitution \(s2\), by composing \(s1\) and \(s1b\).

In this function, \(s1b\) is the \((\text{make_unif_subst} (\text{add_id_subst} t \tau))\), \(s1\) is the \((\text{add_id_subst} t \tau)\), \(\tau\) is the original input unifier and \(s2\) is the result of this function which basically composes \(s1b\) with \(s1\).

Definition convert_to_01_subst \((\tau : \text{subst}) (t : \text{term}) : \text{subst} :=\) 
\[
\text{subst_compose (make_unif_subst (add_id_subst t \tau)) (add_id_subst t \tau)}
\]

The following lemmas are about facts for the “01” substitutions and our \texttt{convert_to_01_subst} function which gives \(s2\). These lemmas are the ones that prove the facts that the new \(s2\) supposedly has: that it is a unifier, and also a “01” unifier. As stated at the very beggining of this section, these lemmas are very important for the intermediate lemmas section where in the \texttt{unifiable t} case we are trying to prove that when there exists any substitution for a term \(t\), then there exists a “01” substitution; the \texttt{unif_some_subst} lemma.

This is an intuitive lemma that states that adding an identity substitution to an existing unifier of a term gives also a unifier.

Lemma add_id_unf : \(\forall (t : \text{term}) (s1 : \text{subst}),\) 
\[
\text{unifier t \, s1} \rightarrow \text{unifier t (add_id_subst t s1)}.
\]

Proof.
\begin{itemize}
\item intros. induction \(s1\).
\item induction \(t\).
\item unfold unifier in *. simpl in *. apply \(H\).
\item unfold unifier in *. simpl in *. apply \(H\).
\item unfold unifier in *. simpl in *. destruct PeanoNat.Nat.eqb. apply \(H\).
\item apply \(H\).
\item unfold unifier in *. simpl in *. unfold add_id_subst. simpl.
\end{itemize}
Admitted.

This lemma states two facts, given a term \(t\) and a unifier \(s1\) of \(t\): 1) The \texttt{convert_to_01_subst} substitution is also a unifier. 2) Applying the \texttt{convert_to_01_subst} substitution on the term results in a term that is ground.

Lemma unif_grnd_unif : \(\forall (t : \text{term}) (s1 : \text{subst}),\) 
\[
\text{unifier t \, s1} \rightarrow \text{(unifier t (subst_compose (make_unif_subst (add_id_subst t s1))) (add_id_subst t s1))} \wedge \text{(is_ground_term (apply_subst t (subst_compose (make_unif_subst (add_id_subst t s1)))))}
\]
Proof.
intros. split.
- unfold unifier. unfold unifier in H. rewrite subst-compose_eqv.
  pose proof add_id_unf. specialize (H0 t sig1). unfold unifier in H0.
  specialize (H0 H). rewrite H0. simpl. reflexivity.
- admit.
Admitted.

If a substitution \( \text{sig1} \) is a “01” substitution and the domain of the substitution is a subset of a list of variable \( \text{l1} \) then the substitution \( \text{sig1} \) is an element of the set of all “01” substitutions of that list \( \text{l1} \).

**Lemma 01_in_all**: \( \forall (\text{l1} : \text{list var}) (\text{sig} : \text{subst}), \)
\( \text{is_01_subst ~ \text{sig} = \text{true} ~ \land ~ \text{sub_dmn_list ~ \text{l1} ~ (subst_domain ~ \text{sig}) ~ \rightarrow ~ ln ~ \text{sig} ~ (all_01_substs ~ \text{l1})} \).

Proof.
intros. destruct H. unfold is_01_subst in H.
Admitted.

Here is a specialized format of the _01_in_all lemma. Instead of \( \text{l1} \) we have \text{term_unique_vars t}.

**Lemma 01_in_rec**: \( \forall (\text{t} : \text{term}) (\text{sig} : \text{subst}), \)
\( \text{is_01_subst ~ \text{sig} = \text{true} ~ \land ~} \)
\( \text{sub_dmn_list (term_unique_vars t) (subst_domain ~ \text{sig}) ~ \rightarrow ~ ln ~ \text{sig} ~ (all_01_substs (term_unique_vars t))} \).

Proof.
intros.
pose proof _01_in_all.
specialize (H0 (term_unique_vars t) sig).
apply H0. apply H.
Qed.

Here is a lemma to show that given a unifier \( \text{sig1} \) of \text{t}, then the \text{convert_to_01_subst} substitution is a “01” subst and also the variables of \text{term t} are a subset of the domain of the \text{convert_to_01_subst} substitution.

**Lemma make_unif_is_01**: \( \forall (\text{t} : \text{term}) (\text{sig1} : \text{subst}), \)
\( \text{unifier ~ t ~ sig1 ~ \rightarrow ~} \)
\( \text{is_01_subst} ~ \text{(subst-compose (make_unif_subst (add_id_subst t sig1)))} \)
\( \text{(add_id_subst t sig1)) = \text{true ~ \land ~} \)
\( \text{sub_dmn_list} \)
\( \text{(term_unique_vars t)} \)
\( \text{(subst_domain ~ (subst-compose (make_unif_subst (add_id_subst t sig1)))} \)
\( \text{(add_id_subst t sig1)))}. \)

Proof.
This is a lemma to show that given a unifier of term $t$, then there exists a substitution sig2 that 1) belongs to all the “01” substitutions of term $t$ and it also unifies $t$, by making $t$ equal to T0 when applied on it (it is equal, not just equivalent because we want sig2 to be a ground substitution too).

**Lemma unif_exists_grnd_unif**: $\forall (t : \text{term}) (\text{sig1 : subst}),$

$\exists \text{sig2 : subst},$

ln sig2 (all_01_substs (term_unique_vars t)) \&

match update_term t sig2 with

| T0 ⇒ true
| _ ⇒ false

end = true.

**Proof.**

intros. $\exists (\text{subst-compose (make_unif_subst (add_id_subst t sig1))})$

(add_id_subst t sig1)). split.

- pose proof _01_in_rec as H1.

  specialize (H1 t (subst-compose (make_unif_subst (add_id_subst t sig1))))

  (add_id_subst t sig1))).

  pose proof make_unif_is_01 as H2. specialize (H2 t sig1).

  specialize (H2 H).

  specialize (H1 H2). apply H1.

- pose proof unif_grnd_unif. specialize (H0 t sig1 H). destruct H0.

  unfold unifier in H0. unfold update_term. pose proof simplify_eqv.

  specialize (H2 (apply_subst t (subst-compose (make_unif_subst

  (add_id_subst t sig1))) (add_id_subst t sig1))))).

  symmetry in H2. pose proof trans_compat2. symmetry in H0.

  specialize (H3 T0 (apply_subst t (subst-compose (make_unif_subst

  (add_id_subst t sig1))) (add_id_subst t sig1))) (simplify (apply_subst t

  (subst-compose (make_unif_subst (add_id_subst t sig1))

  (add_id_subst t sig1))))).

  specialize (H3 H0 H2). symmetry in H3. pose proof simplify_eq_T0.

  specialize (H4 (apply_subst t (subst-compose (make_unif_subst

  (add_id_subst t sig1))) (add_id_subst t sig1))))).

  symmetry in H0. rewrite H4.

  + reflexivity.

  + split.

  × apply H0.

  × apply H1.

Qed.
4.5.3 Intermediate Lemmas

In this subsection we prove a series of lemmas for each of the two statements of the final proof, which were: 1) if a term is unifiable, then the Lowenheim_Main function produces a most general unifier (mgu). 2) if a term is not unifiable, then Lowenheim_Main function produces a None substitution.

Not unifiable $t$ case

In this section we prove intermediate lemmas useful for the second statement of the final proof: if a term is not unifiable, then Lowenheim_Main function produces a None substitution.

This is a lemma to show that if find_unifier returns Some subst, the term is unifiable.

Lemma some_subst_unifiable: $\forall (t : \text{term}),$

$(\exists \text{sig}, \text{find\_unifier } t = \text{Some } \text{sig}) \rightarrow \text{unifiable } t.$

Proof.

intros.
destruct $H$ as [sig1 H1].
induction $t$.
- unfold unifiable. $\exists {}[\text{].}$ unfold unifier. simpl. reflexivity.
- simpl in H1. inversion H1.
- unfold unifiable. $\exists \text{sig1.}$ unfold find_unifier in H1.
  apply find_some in H1. destruct H1.
  remember (update_term (VAR $v$) sig1) in H0.
  destruct t.
  + unfold update_term in Heqt. pose proof simplify_eqv.
    specialize ($H1 \ (\text{apply\_subst } (VAR \ v) \ \text{sig1})$). unfold unifier.
    symmetry in Heqt. rewrite Heqt in H1. rewrite H1. reflexivity.
    + inversion H0.
    + inversion H0.
    + inversion H0.
    + inversion H0.
  - unfold unifiable. $\exists \text{sig1.}$ unfold find_unifier in H1.
    apply find_some in H1. destruct H1.
    remember (update_term (t1 + t2) sig1) in H0.
    destruct t.
    + unfold update_term in Heqt. pose proof simplify_eqv.
      specialize ($H1 \ (\text{apply\_subst } (t1 + t2) \ \text{sig1})$).
      symmetry in Heqt. unfold unifier. rewrite Heqt in H1. rewrite H1.
      reflexivity.
      + inversion H0.
      + inversion H0.
      + inversion H0.
+ inversion H0.
- unfold unifiable. \(\exists sig1.\) unfold find_unifier in H1.
  apply find_some in H1. destruct H1.
  remember (update_term (t1 × t2) sig1) in H0.
  destruct t.
+ unfold update_term in Heqt. pose proof simplify_eqv.
  specialize (H1 (apply_subst (t1 × t2) sig1)).
  symmetry in Heqt. unfold unifier. rewrite Heqt in H1. rewrite H1.
  reflexivity.
+ inversion H0.
+ inversion H0.
+ inversion H0.
+ inversion H0.
Qed.

This lemma shows that if no substitution makes find_unifier to return Some subst, the it returns None.

**Lemma not_Some_is_None** \((t:\ term)\):

\(\neg (\exists sig,\ find\_unifier\ t = Some\ sig) \rightarrow find\_unifier\ t = None.\)

**Proof.**

- apply contrapositive_opposite.
- intros H.
- apply not_None_is_Some in H.
- tauto.
Qed.

This is a lemma to show that if a term \(t\) is not unifiable, the find_unifier function returns None with \(t\) as input.

**Lemma not_unifiable_find_unifier_none_subst** : \(\forall (t:\ term),\neg unifiable\ t \rightarrow find\_unifier\ t = None.\)

**Proof.**

- intros.
- pose proof some_subst_unifiable.
- specialize (H0 t).
- pose proof contrapositive.
- specialize (H1 (\(\exists sig :\ subst,\ find\_unifier\ t = Some\ sig\) (unifiable \(t\)))).
- specialize (H1 H0). specialize (H1 H).
- pose proof not_Some_is_None.
- specialize (H2 t H1).
- apply H2.
Qed.

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Unifiable $t$ Case

In this section we prove intermediate lemmas useful for the first statement of the final proof: if a term is unifiable, then Lowenheim\_Main function produces a most general unifier (mgu).

Lemma to show that if find\_unifier on an input term $t$ returns Some $\sigma$, then $\sigma$ is a unifier of $t$.

Lemma Some\_subst\_unifiable : $\forall (t : \text{term}) (sig : \text{subst}),$

\[
\text{find\_unifier } t = \text{Some } sig \rightarrow \text{unifier } t sig.
\]

Proof.

intros. unfold find\_unifier in $H$.
induction $t$.
- simpl in $H$. apply eq\_some\_eq\_subst in $H$. symmetry in $H$. rewrite $H$.
  unfold unifier. simpl. reflexivity.
- simpl in $H$. inversion $H$.
- unfold find\_unifier in $H$. apply find\_some in $H$. destruct $H$.
  remember (update\_term (VAR v) sig) in $H0$.
  destruct $t$.
  + unfold unifier. unfold update\_term in $Hegt$. pose proof simplify\_eqv.
    specialize ($H1$ (apply\_subst (VAR v) sig)). symmetry in $Hegt$.
    rewrite $Hegt$ in $H1$. rewrite $H1$. reflexivity.
  + inversion $H0$.
  + inversion $H0$.
  + inversion $H0$.
- unfold find\_unifier in $H$. apply find\_some in $H$. destruct $H$.
  remember (update\_term (t1 + t2) sig) in $H0$.
  destruct $t$.
  + unfold unifier. unfold update\_term in $Hegt$. pose proof simplify\_eqv.
    specialize ($H1$ (apply\_subst (t1 + t2) sig)). symmetry in $Hegt$.
    rewrite $Hegt$ in $H1$. rewrite $H1$. reflexivity.
  + inversion $H0$.
  + inversion $H0$.
  + inversion $H0$.
- unfold find\_unifier in $H$. apply find\_some in $H$. destruct $H$.
  remember (update\_term (t1 \times t2) sig) in $H0$.
  destruct $t$.
  + unfold unifier. unfold update\_term in $Hegt$. pose proof simplify\_eqv.
    specialize ($H1$ (apply\_subst (t1 \times t2) sig)). symmetry in $Hegt$.
    rewrite $Hegt$ in $H1$. rewrite $H1$. reflexivity.
  + inversion $H0$.
  + inversion $H0$.
  + inversion $H0$. 

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This lemma is the one using all the utilities defined in the “utilities and admitted lemmas...” section for the unifiable t case. It states that if there is a unifier $\sigma_1$ for term $t$ then there exists some substitution $\sigma_2$ for which the $\text{find\_unifier}$ function returns $\text{Some} \ \sigma_2$. Here is the main outline of the proof: As done in the utilities section, given any unifier $\sigma_1$ of a term $t$, we can find a “01” unifier. Since our $\text{find\_unifier}$ function also finds a “01” unifier by going through the list of available “01” unifiers, there must exist a “01” unifier $\sigma_2$ returned by our $\text{find\_unifier}$ function under the $\text{Some}$ wrapper.

Lemma unif_some_subst : $\forall (t : \text{term})$,  
($\exists \ \sigma_1, \ \text{unifier} \ t \ \sigma_1$) $\rightarrow$  
$\exists \ \sigma_2, \ \text{find\_unifier} \ t = \text{Some} \ \sigma_2$.

Proof.
intros $t \ H$. induction $t$.
- simpl. $\exists \ [\].$ reflexivity.
- simpl. destruct $H$. unfold unifier in $H$. simpl in $H$.
  - apply $T1\_not\_equiv\_T0$ in $H$. inversion $H$.
- unfold $\text{find\_unifier}$. 
  apply existsb_find.
  apply existsb_exists. destruct $H$. pose proof unif_exists_grnd_unif.
  specialize ($H0 \ (\text{VAR} \ v) \ x$).
  apply $H0$. apply $H$.
- unfold $\text{find\_unifier}$. 
  apply existsb_find.
  apply existsb_exists. destruct $H$. pose proof unif_exists_grnd_unif.
  specialize ($H0 \ (t1 + t2) \ x$).
  apply $H0$. apply $H$.
- unfold $\text{find\_unifier}$. 
  apply existsb_find.
  apply existsb_exists. destruct $H$. pose proof unif_exists_grnd_unif.
  specialize ($H0 \ (t1 \times t2) \ x$).
  apply $H0$. apply $H$.
Qed.

Here is a lemma to show that if no substituion makes $\text{find\_unifier}$ return $\text{Some} \ \sigma$, then it returns $\text{None}$.

Lemma not_Some_not_unifiable ($t : \text{term}$) :
($\neg \ \exists \ \sigma, \ \text{find\_unifier} \ t = \text{Some} \ \sigma$) $\rightarrow$  
$\neg \ \text{unifiable} \ t$.

Proof.
intros.
pose proof not_Some_is_None.
specialize (H0 t H).
unfold unifiable.
intro.
unfold not in H.
pose proof unif_some_subst.
specialize (H2 t H1).
specialize (H H2).
apply H.
Qed.

This lemma shows that if a term is unifiable then find_unifier returns Some σ.

Lemma unifiable_find_unifier_some_subst : ∀ (t : term),
unifiable t →
(∃ sig, find_unifier t = Some sig).
Proof.
intros.
pose proof contrapositive.
specialize (H0 (¬ ∃ sig, find_unifier t = Some sig) (¬ unifiable t)).
pose proof not_Some_not_unifiable.
specialize (H1 t). specialize (H0 H1). apply NNPP in H0.
- apply H0.
- firstorder.
Qed.

This lemma shows that if a term is unifiable, then find_unifier returns a unifier.

Lemma find_unifier_is_unifier: ∀ (t : term),
unifiable t → unifier t (convert_to_subst (find_unifier t)).
Proof.
intros.
pose proof unifiable_find_unifier_some_subst.
specialize (H0 t H).
unfold unifier. unfold unifiable in H. simpl. unfold convert_to_subst.
destruct H0 as [sig H0]. rewrite H0.
pose proof Some_subst_unifiable.
specialize (H1 t sig). specialize (H1 H0).
unfold unifier in H1.
apply H1.
Qed.

4.5.4 Gluing Everything Together For the Final Proof

In this subsection we prove the two top-level final proof lemmas. Both of these proofs use the intermediate lemmas proved in the previous subsections.
The first one states that given a unifiable term \( t \) and the fact that our Lowenheim builder produces an mgu, then the Lowenheim Main function also produces an mgu. This is the part of the final proof for Lowenheim Main that uses the building block that was provided by the previous section where we had proved that our “lowenheim’s builder” produces an mgu given a unifiable term \( t \).

**Lemma builder_to_main:** \( \forall (t : \text{term}), \)

unifiable \( t \to \)

most_general_unifier \( t \) (build_lowenheim_subst

\( t \) (convert_to_subst (find_unifier \( t \)))) \to

most_general_unifier \( t \) (convert_to_subst (Lowenheim_Main \( t \))).

**Proof.**

intros.

pose proof lowenheim_most_general_unifier as \( H1 \).

pose proof find_unifier_is_unifier as \( H2 \).

specialize \((H2 t H)\). specialize \((H1 t (convert_to_subst (find_unifier \( t \))))\).

specialize \((H1 H2)\). unfold Lowenheim_Main. destruct (find_unifier \( t \)).

- simpl. simpl in \( H1 \). apply \( H1 \).

- simpl in \( H2 \). unfold unifier in \( H2 \). apply app_subst_T0 in \( H2 \). simpl.

  repeat simpl in \( H1 \). pose proof most_general_unifier_compat.

specialize \((H3 t T0 H2)\). specialize \((H3 [])\).

rewrite \( H3 \). unfold most_general_unifier. intros.

unfold more_general_substitution. split.

+ apply empty_subst_on_term.

+ intros. \( \exists s' \). unfold substitution_factor_through.

  intros. simpl. reflexivity.

Qed.

This is the final top-level lemma that encapsulates all our efforts so far. It proves the two main statements required for the final proof. The two statements, as phrased in the beginning of the chapter are: 1) If a term is unifiable, then our own defined Lowenheim Main function produces a most general unifier (mgu). 2) If a term is not unifiable, then our own defined Lowenheim Main function produces a None substitution. The two propositions are related with the “\( \land \)” symbol (namely, the propositional “and”) and each is proven separately using the intermediate lemmas proven in the previous section. This is why the final top-level proof is relatively short, because a lot of the significant components of the proof have already been proven as intermediate lemmas and in previous helper sections.

**Lemma lowenheim_main_most_general_unifier:** \( \forall (t : \text{term}), \)

(unifiable \( t \to \) most_general_unifier \( t \) (convert_to_subst (Lowenheim_Main \( t \))))

\( \land \)

(\( \neg \) unifiable \( t \to \) Lowenheim_Main \( t \) = None).

**Proof.**

intros.

split.
- intros. apply builder_to_main.
  + apply \( H \).
  + apply lowenheim_most_general_unifier. apply find_unifier_is_unifier.
    apply \( H \).
- intros. pose proof not_unifiable_find_unifier_none_subst.
  specialize \((H0 \; t \; H)\). unfold Lowenheim_Main. rewrite \( H0 \). reflexivity.
Qed.
Chapter 5

Library B_Unification.list_util

Require Import List.
Import ListNotations.
Require Import Arith.
Import Nat.
Require Import Sorting.
Require Import Permutation.
Require Import Omega.

5.1 Introduction

The second half of the project revolves around the successive variable elimination algorithm for solving unification problems. While we could implement this algorithm with the same data structures used for Lowenheim’s, this algorithm lends itself well to a new representation of terms as polynomials.

A polynomial is a list of monomials being added together, where a monomial is a list of variables being multiplied together. Since one of the rules is that $x \times x \approx_B x$, we can guarantee that there are no repeated variables in any given monomial. Similarly, because $x + x \approx_B 0$, we can guarantee that there are no repeated monomials in a polynomial.

Because of these properties, as well as the commutativity of addition and multiplication, we can represent both monomials and polynomials as unordered sets of variables and monomials, respectively. For simplicity when implementing and comparing these polynomials in Coq, we have opted to use the standard list structure, instead maintaining that the lists are maintained in our polynomial form after each stage.

In order to effectively implement polynomial lists in this way, a set of utilities are needed to allow us to easily perform operations on these lists. This file serves to implement and prove facts about these functions, as well as to expand upon the standard library when necessary.
5.2 Comparisons Between Lists

Checking if a list of natural numbers is sorted is easy enough. Comparing lists of lists of nats is slightly harder, and requires the use of a new function, called \texttt{lex}. \texttt{lex} simply takes in a comparison and applies the comparison across the list until it finds a point where the elements are not equal.

In all cases throughout this project, the comparator used will be the standard nat \texttt{compare} function.

For example, [1;2;3] is less than [1;2;4], and [1;2] is greater than [1].

\begin{verbatim}
Fixpoint lex {T} (cmp:T → T → comparison) (l1 l2:list T) : comparison :=
  match l1, l2 with
  | [], [] ⇒ Eq
  | [], _ ⇒ Lt
  | _, [] ⇒ Gt
  | h1 :: t1, h2 :: t2 ⇒
    match cmp h1 h2 with
    | Eq ⇒ lex cmp t1 t2
    | c ⇒ c
  end
end.
\end{verbatim}

There are some important but relatively straightforward properties of this function that are useful to prove. First, \textit{reflexivity}:

\begin{verbatim}
Lemma lex_nat_refl : ∀ l, lex compare l l = Eq.
Proof.
  intros.
  induction l; auto.
  simpl. rewrite compare_refl. apply IHl.
Qed.
\end{verbatim}

Next, \textit{antisymmetry}. This allows us to take a predicate or hypothesis about the comparison of two polynomials and reverse it.

For example, \(l < m\) implies \(m > l\).

\begin{verbatim}
Lemma lex_nat_antisym : ∀ l m,
  lex compare l m = CompOpp (lex compare m l).
Proof.
  intros l. induction l.
  - intros. simpl. destruct m; reflexivity.
  - intros. simpl. destruct m; auto. simpl.
    destruct (a ?= n) eqn:H; rewrite compare_antisym in H;
    rewrite CompOpp_iff in H; simpl in H; rewrite H; auto.
Qed.
\end{verbatim}

It is also useful to convert from the result of \texttt{lex compare} to a hypothesis about equality
in Coq. Clearly, if `lex compare` returns `Eq`, the lists are exactly equal, and if it returns `Lt` or `Gt` they are not.

**Lemma lex_eq :** \( \forall l m, \)
\[
\text{lex compare } l \ m = \text{Eq } \leftrightarrow \ l = m.
\]

**Proof.**
\[
\text{intros } l. \ \text{induction } l; \ \text{induction } m; \ \text{intros.}
\]
- split; reflexivity.
- split; intros; inversion \( H \).
- split; intros; inversion \( H \).
- split; intros; simpl in \( H \).
  + destruct \((a \ ?= a0) \ e\eqn: Hcomp\); try inversion \( H \).\ f\_equal.
    × apply \( \text{compare\_eq\_iff } Hcomp \); auto.
    × apply \( \text{IHl} \).\ auto.
  + inversion \( H \).\ simpl.\ rewrite \( \text{compare\_refl} \).
    rewrite \(\leftarrow H2 \).\ apply \( \text{IHl} \).\ reflexivity.

Qed.

**Lemma lex_neq :** \( \forall l m, \)
\[
\text{lex compare } l \ m = \text{Lt } \lor \ \text{lex compare } l \ m = \text{Gt } \leftrightarrow \ l \neq m.
\]

**Proof.**
\[
\text{intros } l. \ \text{induction } l; \ \text{induction } m.
\]
- simpl.\ split; intro.\ inversion \( H \); inversion \( H0 \).\ contradiction.
- simpl.\ split; intro.\ intro.\ inversion \( H0 \).\ auto.
- simpl.\ split; intro.\ intro.\ inversion \( H0 \).\ auto.
- clear \( \text{IHm}.\ split; \ \text{intros.} \)
  + destruct \( H \); intro; apply \( \text{lex\_eq } H0 \); rewrite \( H \) in \( H0 \); inversion \( H0 \).
  + destruct \((a \ ?= a0) \ e\eqn: Hcomp\).
    × simpl.\ rewrite \( Hcomp \).\ apply \( \text{IHl} \).\ apply \( \text{compare\_eq\_iff } Hcomp \).
    rewrite \( Hcomp \) in \( H \).\ intro.\ apply \( H \).\ rewrite \( H0 \).\ reflexivity.
    × left.\ simpl.\ rewrite \( Hcomp \).\ reflexivity.
    × right.\ simpl.\ rewrite \( Hcomp \).\ reflexivity.

Qed.

**Lemma lex_neq' :** \( \forall l m, \)
\[
(\text{lex compare } l \ m = \text{Lt } \rightarrow \ l \neq m) \ \land
(\text{lex compare } l \ m = \text{Gt } \rightarrow \ l \neq m).
\]

**Proof.**
\[
\text{intros } l. \ \text{split; repeat (intros; apply \lex\_neq; auto).}
\]

Qed.

It is also useful to be able to flip the arguments of a call to `lex compare`, since these two comparisons impact each other directly.

If `lex compare` returns that \( l = m \), then this also means that \( m = l \). More interesting is that if \( l < m \), then \( m > l \).
Lemma lex_rev_eq : \forall l m,  
    lex compare l m = Eq ↔ lex compare m l = Eq.
Proof.
    intros l m. split; intro; rewrite lex_nat antisym in H; unfold CompOpp in H.
    - destruct (lex compare m l) eqn:H0; inversion H. reflexivity.
    - destruct (lex compare l m) eqn:H0; inversion H. reflexivity.
Qed.

Lemma lex_rev.lt.gt : \forall l m,  
    lex compare l m = Lt ↔ lex compare m l = Gt.
Proof.
    intros l m. split; intro; rewrite lex_nat antisym in H; unfold CompOpp in H.
    - destruct (lex compare m l) eqn:H0; inversion H. reflexivity.
    - destruct (lex compare l m) eqn:H0; inversion H. reflexivity.
Qed.

Lastly is a property over lists. The comparison of two lists stays the same if the same
new element is added onto the front of each list. Similarly, if the item at the front of two
lists is equal, removing it from both does not change the lists’ comparison.

Lemma lex.nat.cons : \forall l m n,  
    lex compare l m = lex compare (n :: l) (n :: m).
Proof.
    intros. simpl. rewrite compare_refl. reflexivity.
Qed.

Hint Resolve lex_nat_refl lex_nat_antisym lex_nat_cons.

5.3 Extensions to the Standard Library

There were some facts about the standard library list functions that we found useful to prove,
as they repeatedly came up in proofs of our more complex custom list functions.

Specifically, because we are comparing sorted lists, it is often easier to disregard the
sortedness of the lists and instead compare them as permutations of one another. As a
result, many of the lemmas in the rest of this file revolve around proving that two lists are
permutations of one another.

5.3.1 Facts about In

First, a very simple fact about In. This mostly follows from the standard library lemma
Permutation_in, but is more convenient for some of our proofs when formalized like this.

Lemma Permutation_not_In : \forall A (a:A) l l',  
    Permutation l l' →  
    ¬ In a l →  
    ¬ In a l'.
5.3.2 Facts about incl

Next are some useful lemmas about incl. First is that if one list is included in another, but one element of the second list is not in the first, then the first list is still included in the second with that element removed.

Lemma incl_not_in : ∀ {A} (l:l: list A),
    incl l (a :: m) →
    ¬ ln a l →
    incl l m.

Proof.
intros A a l m Hincl Hnin. unfold incl in *. intros a0 Hin.
simpl in Hincl destruct (Hincl a0); auto. rewrite H in Hnin. contradiction.
Qed.

We also found it useful to relate Permutation to incl; if two lists are permutations of each other, then they must be set equivalent, or contain all of the same elements.

Lemma Permutation_incl : ∀ {A} (l m:list A),
    Permutation l m → incl l m ∧ incl m l.

Proof.
intros A l m H. apply Permutation_sym in H as H0. split.
+ unfold incl. intros a. apply (Permutation_in a H).
+ unfold incl. intros a. apply (Permutation_in a H0).
Qed.

Unfortunately, the definition above cannot be changed into an iff relation, as incl proves nothing about the lengths of the lists. We can, however, prove that if some list m includes a list l, then m includes all permutations of l.

Lemma incl_Permutation : ∀ {A} (l l’ m:list A),
    Permutation l l’ →
incl \ l \ m \rightarrow
incl \ l' \ m.

Proof.

intros \ A \ l \ l' \ m \ H \ H0. apply Permutation_incl in \ H as [].
apply incl_tran with \ (m:=l); auto.
Qed.

A really simple lemma is that if some list \ l \ is included in the empty list, then \ l \ must also
be empty.

Lemma incl_nil : \ \ \ \forall \ \ \ \{X\} \ (l:list \ X),
incl \ l \ [] \leftrightarrow \ l = [].

Proof.

intros \ X \ l. unfold incl. split; intro \ H.
- destruct \ l;[auto | destruct (H \ x); intuition].
- intros \ a Hin. destruct \ l;[auto | rewrite \ H \ in Hin; auto].
Qed.

The last fact about incl is simply a new way of formalizing the definition that is convenient
for some proofs.

Lemma incl_cons_inv : \ \ \ \forall \ A \ (a:A) \ l \ m,
incl \ (a :: l) \ m \rightarrow \ In \ a \ m \land incl \ l \ m.

Proof.

intros \ A \ a \ l \ m \ H. split.
- unfold incl in \ H. apply \ H. intuition.
- unfold incl in *.
intros \ b Hin. apply \ H. intuition.
Qed.

5.3.3 Facts about count_occ

Next is some facts about count_occ. Firstly, if two lists are permutations of each other, than
every element in the first list has the same number of occurrences in the second list.

Lemma count_occ_Permutation : \ \ \ \forall \ A \ Aeq_dec \ (a:A) \ l \ l',
Permutation \ l \ l' \rightarrow
count_occ \ Aeq_dec \ l \ a = count_occ \ Aeq_dec \ l' \ a.

Proof.

intros \ A \ Aeq_dec \ a \ l \ l' \ H. induction \ H.
- auto.
- simpl. destruct (Aeq_dec \ x \ a); auto.
- simpl. destruct (Aeq_dec \ y \ a); destruct (Aeq_dec \ x \ a); auto.
- rewrite <- IHPermutation2. rewrite IHPermutation1. auto.
Qed.

The function count_occ also distributes over list concatenation, instead becoming addi-
tion. This is useful especially when dealing with count occurrences of lists during induction.
Lemma count_occ_app : \( \forall \{ A \} (a:A) \) \( l m \) \( \text{Aeq\_dec} \),
\[
\text{count\_occ Aeq\_dec} \ l \ m \ a = \text{add (count\_occ Aeq\_dec} \ l \ a) \ (\text{count\_occ Aeq\_dec} \ m \ a).
\]
Proof.
intros \( A \) \( a \) \( l \) \( m \) \( \text{Aeq\_dec} \). induction \( l \); auto.
simpl. destruct (\( \text{Aeq\_dec} \ a0 \ a \)); simpl; auto.
Qed.

It is also convenient to reason about the relation between \text{count\_occ} and \text{remove}. If the element being removed is the same as the one being counted, then the count is obviously 0. If the elements are different, then the count is the same with or without the remove.

Lemma count_occ_remove : \( \forall \{ A \} \) \( \text{Aeq\_dec} \ (a:A) \) \( l \),
\[
\text{count\_occ Aeq\_dec} \ (\text{remove Aeq\_dec} \ a \ l) \ a = 0.
\]
Proof.
intros \( A \) \( \text{Aeq\_dec} \ a \) \( l \). induction \( l \); auto.
- destruct (\( \text{Aeq\_dec} \ a0 \ a \)) eqn:Haa0; auto. simpl.
destructor (\( \text{Aeq\_dec} \ a0 \ a \)); try (symmetry in \( e \); contradiction). apply IHl.
Qed.

Lemma count_occ_neq_remove : \( \forall \{ A \} \) \( \text{Aeq\_dec} \ (a b:A) \) \( l \),
a \( \neq \) b \( \rightarrow \)
\[
\text{count\_occ Aeq\_dec} \ (\text{remove Aeq\_dec} \ a \ l) \ b = \text{count\_occ Aeq\_dec} \ l \ b.
\]
Proof.
intros \( A \) \( \text{Aeq\_dec} \ a \) \( b \) \( l \) \( H \). induction \( l \); simpl; auto. destruct (\( \text{Aeq\_dec} \ a0 \ a \)).
- destruct (\( \text{Aeq\_dec} \ a0 \ b \)); auto.
  - rewrite \( \leftarrow e0 \) in \( H \). rewrite \( e \) in \( H \). contradiction.
- simpl. destruct (\( \text{Aeq\_dec} \ a0 \ b \)); auto.
Qed.

5.3.4 Facts about \text{concat}

Similarly to the lemma \text{Permutation\_map}, \text{Permutation\_concat} shows that if two lists are permutations of each other then the flattening of each list are also permutations.

Lemma Permutation_concat : \( \forall \{ A \} \) \( (l m:\text{list (list A)}) \),
\[
\text{Permutation} \ l \ m \rightarrow \text{Permutation} \ (\text{concat} \ l) \ (\text{concat} \ m).
\]
Proof.
intros \( A \) \( l \) \( m \) \( H \). induction \( H \).
- auto.
- simpl. apply Permutation\_app\_head. auto.
- simpl. apply Permutation\_trans with (\( l' := \text{concat} \ l \ + \ y \ + \ x \)).
  + rewrite \( \text{app\_assoc} \). apply Permutation\_app\_comm.
+ apply Permutation_trans with (l':=concat l ++ x ++ y).
  \times apply Permutation_app_head. apply Permutation_app_comm.
  \times rewrite (app_assoc x y). apply Permutation_app_comm.
- apply Permutation_trans with (l':=concat l'); auto.
Qed.

Before the creation of this next lemma, it was relatively hard to reason about whether elements are in the flattening of a list of lists. This lemma states that if there is a list in the list of lists that contains the desired element, then that element will be in the flattened version.

Lemma ln_concat_exists : \forall A ll (a:A),
  (\exists l, ln l ll \land a l) \leftrightarrow ln a (concat ll).
Proof.
  intros A ll a. split; intros H.
  - destruct H as [l ||]. apply ln_split in H. destruct H as [ll [l2 H]].
    rewrite H. apply Permutation_in with (l:=concat (l :: ll ++ l2)).
    + apply Permutation_concat. apply Permutation_middle.
    + simpl. apply ln_app_iff. auto.
  - induction ll.
    + inversion H.
    + simpl in. apply ln_app_iff in H. destruct H.
      \times \exists a0. split; intuition.
      \times destruct IHll; auto. \exists x. intuition.
Qed.

This particular lemma is useful if the function being mapped returns a list of its input type. If the resulting lists are flattened after, then the result is the same as mapping the function without converting the output to lists.

Lemma concat_map : \forall \{A B\} (f:A\rightarrow B) l,
  concat (map (fun a \Rightarrow [f a]) l) = map f l.
Proof.
  intros A B f l. induction l; auto. simpl. f_equal. apply IHl.
Qed.

Another fact similar to the last is that if you flatten the result of mapping a function that maps a function over a list, we can rearrange the order of the concat and the maps.

Lemma concat_map_map : \forall A B C l (f:B\rightarrow C) (g:A\rightarrow list B),
  concat (map (fun a \Rightarrow map f (g a)) l) =
  map f (concat (map g l)).
Proof.
  intros. induction l; auto.
  simpl. rewrite map_app. f_equal. auto.
Qed.
Lastly, if you map a function that converts every element of a list to nil, and then concat the list of nils, you end with nil.

Lemma `concat_map_nil` : \( \forall \{A\} \ (l:list \ A) \),
\[ \text{concat} \ (\text{map} \ (\text{fun} \ x \Rightarrow \ []) \ l) = (@\text{nil} \ A). \]
Proof.
induction \( l \); auto.
Qed.

5.3.5 Facts about `Forall` and `existsb`

This is similar to the inverse of `Forall`; any element in a list \( l \) must hold for predicate \( p \) if `Forall \ p` is true of \( l \).

Lemma `Forall_In` : \( \forall \ A \ (l:list \ A) \ a \ p, \)
\[ \text{In} \ a \ l \rightarrow \text{Forall} \ p \ l \rightarrow p \ a. \]
Proof.
intros \( A \ a \ p \ Hin \ Hfor. \) apply `(Forall_forall \ p \ l)`; auto.
Qed.

In Coq, `existsb` is effectively the “or” to `Forall`’s “and” when reasoning about lists. If there does not exist a single element in a list \( l \) where the predicate \( p \) holds, then \( p \ a \) must be false for any element \( a \) of \( l \).

Lemma `existsb_false_forall` : \( \forall \{A\} \ p \ (l:list \ A), \)
\[ \exists b \ p \ l = \text{false} \rightarrow \]
\[ (\forall a, \text{In} \ a \ l \rightarrow p \ a = \text{false}). \]
Proof.
intros \( A \ p \ l \ H a \ Hin. \) destruct \((p \ a) \ eqn:Hpa\); auto.
exfalso. rewrite ← `Bool.negb_true_iff` in \( H \).
apply `(Bool.eq_true_false_abs _ H)`. rewrite `Bool.negb_false_iff`.
apply `existsb_exists`. \exists \( a \). split; auto.
Qed.

Similarly to `Forall_In`, this lemma is just another way of formalizing the definition of `Forall` that proves useful when dealing with `StronglySorted` lists.

Lemma `Forall_cons_iff` : \( \forall \ A \ p \ (a:A) \ l, \)
\[ \text{Forall} \ p \ (a::l) \leftrightarrow \text{Forall} \ p \ l \land p \ a. \]
Proof.
intros \( A \ p \ a \ l \). split.
- intro \( H \). split.
  + rewrite `Forall_forall` in \( H \). apply `Forall_forall`. intros \( x \ Hin. \)
    apply \( H \). intuition.
  + apply `Forall_inv` in \( H \). auto.
- intros \( [] \). apply `Forall_cons`; auto.
Qed.
If a predicate \( p \) holds for all elements of a list \( l \), then \( p \) still holds if some elements are removed from \( l \).

**Lemma Forall_remove**: \( \forall A \text{Aeq_dec} \ p \ (a:A) \ l, \)  
\[ \text{Forall } p \ l \rightarrow \text{Forall } p \ (\text{remove Aeq_dec a} \ l). \]

**Proof**.
intros \( A \text{Aeq_dec} p \ a \ l \ H \). induction \( l \). auto. simpl.
apply Forall_cons_iff in \( H \). destruct \((\text{Aeq_dec a} \ a0)\).
- apply IHl. apply \( H \).
- apply Forall_cons_iff. split.
  + apply IHl. apply \( H \).
  + apply \( H \).
Qed.

This next lemma is particularly useful for relating **StronglySorted** lists to **Sorted** lists; if some relation holds between all elements of a list, then this can be converted to the \( \text{HdRel} \) proposition used by **Sorted**.

**Lemma Forall_HdRel**: \( \forall \{X\} \ r \ (x:X) \ l, \)
\[ \text{Forall } (r \ x) \ l \rightarrow \text{HdRel } r \ x \ l. \]

**Proof**.
intros \( X \ r \ x \ l \ H \). destruct \( l \).
- apply HdRel_nil.
- apply HdRel_cons. apply Forall_inv in \( H \). auto.
Qed.

Lastly, if some predicate \( p \) holds for all elements in a list \( l \), and the elements of a second list \( m \) are all included in \( l \), then \( p \) holds for all the elements in \( m \).

**Lemma Forall_incl**: \( \forall \{X\} \ p \ (l \ m: \text{list X}), \)
\[ \text{Forall } p \ l \rightarrow \text{incl } m \ l \rightarrow \text{Forall } p \ m. \]

**Proof**.
intros \( X \ p \ l \ m \ H \ H0. \) induction \( m \).
- apply Forall_nil.
- rewrite Forall_forall in \( H \). apply Forall_forall. intros \( x \ Hin \).
  apply \( H \). unfold incl in \( H0 \). intros \( x \ Hin \).
Qed.

### 5.3.6 Facts about remove

There are surprisingly few lemmas about **remove** in the standard library, so in addition to those proven in other places, we opted to add quite a few simple facts about **remove**. First is that if an element is in a list after something has been removed, then clearly it was in the list before as well.

**Lemma In_remove**: \( \forall \{A\} \text{Aeq_dec} \ (a:b:A) \ l, \)  
\[ \ln a (\text{remove Aeq_dec b} \ l) \rightarrow \ln a \ l. \]
Proof.
intros A Aeq dec a b l H. induction l as [| c l IHl]; auto.
destruct (Aeq dec b c) eqn:Heq; simpl in H; rewrite Heq in H.
  - right. auto.
  - destruct H; [rewrite H; intuition | right; auto].
Qed.

Similarly to Forall_remove, if a list was StronglySorted before something was removed then it is also StronglySorted after.

Lemma StronglySorted_remove : \forall \{A\} Aeq_dec r (a:A) l, StronglySorted r l \rightarrow StronglySorted r (remove Aeq_dec a l).
Proof.
intros A Aeq dec a l m H. induction l; auto.
simpl. apply StronglySorted_inv in H. destruct (Aeq_dec a a0).
  - apply IHl. apply H.
  - apply SSorted_cons.
    + apply IHl. apply H.
    + apply Forall_remove. apply H.
Qed.

If the item being removed from a list isn’t in the list, then the list is equal with or without the remove.

Lemma not_In_remove : \forall A Aeq_dec (a:A) l, \neg In a l \rightarrow remove Aeq_dec a l = l.
Proof.
intros A Aeq dec a l H. induction l; auto.
simpl. destruct (Aeq_dec a a0).
  - simpl. rewrite e in H. exfalso. apply H. intuition.
  - rewrite IHl. reflexivity. intro Hin. apply H. intuition.
Qed.

The function remove also distributes over list concatenation.

Lemma remove_distr_app : \forall A Aeq_dec (a:A) l m, remove Aeq_dec a (l ++ m) = remove Aeq_dec a l ++ remove Aeq_dec a m.
Proof.
intros A Aeq_dec a l m. induction l; auto.
simpl. destruct (Aeq_dec a a0); auto.
simpl. f_equal. apply IHl.
Qed.

More interestingly, if two lists were permutations before, they are also permutations after the same element has been removed from both lists.

Lemma remove_Permutation : \forall A Aeq_dec (a:A) l l', Permutation l l' \rightarrow
**Permutation** (remove $Aeq_{dec} a l$) (remove $Aeq_{dec} a l'$).

Proof.
- intros $A$, $Aeq_{dec} a l l' H$. induction $H$.
  - auto.
  - simpl. destruct ($Aeq_{dec} a x$); auto.
  - simpl. destruct ($Aeq_{dec} a y$); destruct ($Aeq_{dec} a x$); auto.
    apply perm_swap.
  - apply Permutation_trans with ($l'$:= (remove $Aeq_{dec} a l'$)); auto.
Qed.

The function $remove$ is also associative with itself.

**Lemma remove_remove** : $\forall \{A\} Aeq_{dec} (a b:A) l$,
remove $Aeq_{dec} a$ (remove $Aeq_{dec} b l$) =
remove $Aeq_{dec} b$ (remove $Aeq_{dec} a l$).

Proof.
- intros $A$, $Aeq_{dec} a b l$. induction $l$ as $\|c\|$; simpl; auto.
  destruct ($Aeq_{dec} a b$); destruct ($Aeq_{dec} b c$); destruct ($Aeq_{dec} a c$).
  - auto.
  - rewrite $\leftarrow e0$ in $n$. rewrite $c$ in $n$. contradiction.
  - rewrite $\leftarrow e0$ in $n$. rewrite $e0$ in $n$. contradiction.
  - simpl. destruct ($Aeq_{dec} a c$); try contradiction.
    destruct ($Aeq_{dec} b c$); try contradiction. rewrite $IHl$. auto.
  - rewrite $c$ in $n$. rewrite $e0$ in $n$. contradiction.
  - simpl. destruct ($Aeq_{dec} b c$); try contradiction. auto.
  - simpl. destruct ($Aeq_{dec} a c$); try contradiction. auto.
  - simpl. destruct ($Aeq_{dec} a c$); try contradiction.
    destruct ($Aeq_{dec} b c$); try contradiction. rewrite $IHl$. auto.
Qed.

Lastly, if an element is being removed from a particular list twice, the inner $remove$ is redundant and can be removed.

**Lemma remove_pointless** : $\forall \{A\} Aeq_{dec} (a:A) l m$,
remove $Aeq_{dec} a$ (remove $Aeq_{dec} a l \ ++ \ m$) =
remove $Aeq_{dec} a$ (l \ ++ \ m).

Proof.
- intros $A$, $Aeq_{dec} a l m$. induction $l$; auto. simpl.
  destruct ($Aeq_{dec} a a0$) eqn:Heq; auto.
  simpl. rewrite Heq. f_equal. apply $IHl$.
Qed.

### 5.3.7 Facts about nodup and NoDup

Next up - the **NoDup** proposition and the closely related **nodup** function. The first lemma states that if there are no duplicates in a list, then the first two elements of that list must
not be equal.

Lemma NoDup_neq : \forall \{A\} l (a b:A),
\hspace{1em} NoDup (a :: b :: l) \rightarrow a \neq b.

Proof.
intros A l a b Hdup. apply NoDup_cons_iff in Hdup as [].
apply NoDup_cons_iff in H0 as []. intro. apply H. simpl. auto.
Qed.

In a similar vein as many of the other remove lemmas, if there were no duplicates in a list before the remove then there are still none after.

Lemma NoDup_remove : \forall A Aeq_dec (a:A) l,
\hspace{1em} NoDup l \rightarrow NoDup (remove Aeq_dec a l).

Proof.
intros A Aeq_dec a l H.
induction l; auto.
apply IHl; apply NoDup_cons_iff in H. intuition.
apply IHl; apply NoDup_cons iff in H as []. intro. apply H.
apply (In_remove Aeq_dec a0 a l H1).
apply IHl. apply NoDup_cons_iff in H; intuition.
Qed.

Another lemma similar to NoDup_neq is NoDup_forall_neq; if every element in a list is not equal to a certain a, and the list has no duplicates as is, then it is safe to add a to the list without creating duplicates.

Lemma NoDup_forall_neq : \forall A (a:A) l,
\hspace{1em} Forall (fun b \Rightarrow a \neq b) l \rightarrow NoDup l \rightarrow NoDup (a :: l).

Proof.
intros A a l Hf Hn. apply NoDup_cons; auto.
intro induction l; auto.
apply Forall_cons_iff in Hf as []. apply IHl; auto.
apply NoDup_cons iff in Hn. apply Hn.
simpl in H. destruct H; auto. rewrite H in H1. contradiction.
Qed.

This lemma is really just a reformalization of NoDup_remove_2, which allows us to easily prove that some a is not in the preceding elements l1 or the following elements l2 when the whole list l has no duplicates.

Lemma NoDup_In_split : \forall \{A\} (a:A) l l1 l2,
\hspace{1em} l = l1 ++ a :: l2 \rightarrow NoDup l \rightarrow
¬ \ln a l1 ∧ ¬ \ln a l2.

Proof.
  intros A a l1 l2 H H0. rewrite H in H0.
  apply NoDup_remove_2 in H0. split; intro; intuition.
Qed.

Now some facts about the function \texttt{nodup}; if the \texttt{NoDup} predicate is already true about a certain list, then calling \texttt{nodup} on it changes nothing.

Lemma \texttt{no_nodup_NoDup}: ∀ A Aeq dec (l:list A),
  NoDup l →
  nodup Aeq dec l = l.
Proof.
  intros A Aeq dec H. induction l; auto.
  simpl. apply NoDup_cons_iff in H as ||. destruct (in_dec Aeq dec a l).
  contradiction. f_equal. auto.
Qed.

If a list is sorted (with a transitive relation) before calling \texttt{nodup} on it, the list is also sorted after.

Lemma \texttt{Sorted_nodup} : ∀ A Aeq dec r (l:list A),
  Relations_1.Transitive r →
  Sorted r l →
  Sorted r (nodup Aeq dec l).
Proof.
  intros A Aeq dec r l Ht H. apply Sorted_StronglySorted in H; auto.
  apply StronglySorted_Sorted. induction l; auto.
  simpl. apply StronglySorted_inv in H as ||. destruct (in_dec Aeq dec a l).
  - apply IHl. apply H.
  - apply SSorted_cons.
    + apply IHl. apply H.
    + rewrite Forall_forall in H0. apply Forall_forall. intros x Hin.
      apply H0. apply nodup_In in Hin. auto.
Qed.

We can also show that in some cases, if there are repeated calls to \texttt{nodup}, they are “pointless” - in other words, we can remove the inner call and only keep the outer one.

Lemma \texttt{nodup_pointless} : ∀ l m,
  nodup Nat.eq_dec (l ++ nodup Nat.eq_dec m) = nodup Nat.eq_dec (l ++ m).
Proof.
  intros l m. induction l.
  - simpl. rewrite no_nodup_NoDup; auto. apply NoDup_nodup.
  - simpl. destruct in_dec; destruct in_dec.
    + auto.
    + exfalso. apply n. apply in_app_iff in i; destruct i. intuition.
apply nodup_fn in H; intuition.
+ exfalso. apply n. apply in_app_iff in i; destruct i; intuition.
   apply in_app_iff. right. apply nodup_fn; auto.
+ f_equal. auto.
Qed.

And lastly, similarly to our other Permutation lemmas this far, if two lists were permutations of each other before \textit{nodup} they are also permutations after.

This lemma was slightly more complex than previous Permutation lemmas, but the proof is still very similar. It is solved by induction on the Permutation hypothesis. The first and last cases are trivial, and the second case (where we must prove \textit{Permutation} \((x :: l) (x :: l')\)) becomes simple with the use of \textit{Permutation} in.

The last case (where we must show \textit{Permutation} \((x :: y :: l) (y :: x :: l)\)) was slightly complicated by the fact that destructing \texttt{in_dec} gives us a hypothesis like \texttt{ln} \(x (y :: l)\), which seems useless in reasoning about the other list at first. However, by also destructing whether or not \(x\) and \(y\) are equal, we can easily prove this case as well.

\textbf{Lemma Permutation\_nodup : \(\forall A\ \texttt{Aeq\_dec} (l m:\texttt{list} A),\textit{ Permutation} l m \rightarrow \textit{Permutation} (\texttt{nodup} A\texttt{eq\_dec} l) (\texttt{nodup} A\texttt{eq\_dec} m)\).}

\textbf{Proof.}
intros. induction \(H\).
- auto.
- simpl. destruct (\texttt{in\_dec} A\texttt{eq\_dec} x l).
  + apply \textit{Permutation\_in} with \((l\:'=l')\) in i; destruct \texttt{in\_dec};
    try \texttt{contradiction}. auto.
  + assert \((\neg \texttt{ln} x l')\). intro. apply \(n\).
    apply \textit{Permutation\_sym}; auto. destruct \texttt{in\_dec}; try \texttt{contradiction}; auto.
- destruct (\texttt{in\_dec} A\texttt{eq\_dec} y (x :: l)). destruct \(i\).
  + rewrite \(H\). simpl. destruct (A\texttt{eq\_dec} y y); try \texttt{contradiction}.
    destruct \texttt{in\_dec}; auto.
  + simpl. destruct (A\texttt{eq\_dec} x y). destruct \texttt{in\_dec}; destruct (A\texttt{eq\_dec} y x);
    try (\texttt{symmetry in} e; \texttt{contradiction}). rewrite \(e\) in i. destruct \texttt{in\_dec};
    try \texttt{contradiction}. auto.
  assert \((\neg \texttt{ln} y l)\). intro; apply \(n\); rewrite \(e\); auto.
  destruct \texttt{in\_dec}; try \texttt{contradiction}. destruct \texttt{in\_dec}; try \texttt{contradiction}.
  destruct \texttt{in\_dec}; destruct (A\texttt{eq\_dec} y x);
  try (\texttt{symmetry in} e; \texttt{contradiction}). rewrite \(e0\). destruct \texttt{in\_dec};
  try \texttt{contradiction}. auto. destruct (A\texttt{eq\_dec} y x);
  try (\texttt{symmetry in} e; \texttt{contradiction}).
  assert \((\neg \texttt{ln} y l)\). intro; apply \(n0\); rewrite \(e\); auto.
  destruct \texttt{in\_dec}; try \texttt{contradiction}. rewrite \(e0\). apply \texttt{perm\_skip}; auto.
assert (¬ ln y l). intro; apply n; intuition.
destruct in_dec; try contradiction. destruct in_dec;
destruct (Aeq_dec y x); try (symmetry in e; contradiction); auto.
apply perm_swap.
- apply Permutation_trans with (l’=(nodup Aeq_dec l’)); auto.
Qed.

5.3.8 Facts about partition

The final function in the standard library we found it useful to prove facts about is partition. First, we show the relation between partition and filter: filtering a list gives you a result that is equal to the first list partition would return. This lemma is proven one way, and then reformalized to be more useful in later proofs.

Lemma partition_filter fst \{A\} p l :
  \text{fst} (\text{partition} p l) = @\text{filter} A p l.
Proof.
  induction l; auto. simpl. rewrite ← IHl.
  destruct (\text{partition} p l); simpl.
  destruct (p a); auto.
Qed.

Lemma partition_filter fst’ : \forall \{A\} p (l t f : list A),
  \text{partition} p l = (t, f) →
  \text{t} = @\text{filter} A p l.
Proof.
  intros A p l t f H.
  rewrite ← partition_filter fst.
  now rewrite H.
Qed.

We would like to be able to state a similar fact about the second list returned by partition, but clearly these are all the elements “thrown out” by filter. Instead, we first create a simple definition for negating a function, and prove two quick facts about the relation between some p and \text{neg} p.

Definition neg \{A:Type\} := \text{fun} (p:A→\text{bool}) ⇒ \text{fun} a ⇒ \text{negb} (p a).

Lemma neg_true_false : \forall \{A\} p (a:A),
  p a = true ↔ \text{neg} p a = false.
Proof.
  intros A p a. unfold neg. split; intro.
  - rewrite H. auto.
  - destruct (p a); intuition.
Qed.

Lemma neg_false_true : \forall \{A\} p (a:A),
\[ p \ a = \text{false} \leftrightarrow \neg p \ a = \text{true} \]

Proof.

\begin{itemize}
  \item intros A p a. unfold neg. split; intro.
  \item rewrite H. auto.
  \item destruct (p a); intuition.
\end{itemize}

Qed.

With the addition of this \texttt{neg} proposition, we can now prove two lemmas relating the second \texttt{partition} list and \texttt{filter} in the same way we proved the lemmas about the first \texttt{partition} list.

Lemma \texttt{partition_filter_snd} \{A\} p l:

\[
\text{snd} (\text{partition} p l) = \text{\texttt{@filter} A (\neg p) l}.
\]

Proof.

\begin{itemize}
  \item induction l; auto. simpl.
  \item rewrite \texttt{← IHl}.
  \item destruct (\texttt{partition} p l); simpl.
  \item destruct (p a) eqn:Hp.
  \item - simpl. apply \texttt{neg_true_false} in Hp. rewrite Hp; auto.
  \item - simpl. apply \texttt{neg_false_true} in Hp. rewrite Hp; auto.
\end{itemize}

Qed.

Lemma \texttt{partition_filter_snd'} : \forall \{A\} p (l t f : \texttt{list} A),

\[
\text{partition} p l = (t, f) \rightarrow f = \text{\texttt{@filter} A (\neg p) l}.
\]

Proof.

\begin{itemize}
  \item intros A p l t f H.
  \item rewrite \texttt{← \texttt{partition_filter_snd}}.
  \item now rewrite H.
\end{itemize}

Qed.

These lemmas about \texttt{partition} and \texttt{filter} are now put to use in two important lemmas about \texttt{partition}. If some list \( l \) is partitioned into two lists \((t, f)\), then every element in \( t \) must return \texttt{true} for the filtering predicate and every element in \( f \) must return \texttt{false}.

Lemma \texttt{part_fst_true} : \forall A p (l t f : \texttt{list} A),

\[
\text{partition} p l = (t, f) \rightarrow (\forall a, \text{In} a t \rightarrow p a = \text{true}).
\]

Proof.

\begin{itemize}
  \item intros A p l t f Hpart a Hin.
  \item assert (Hf: t = \texttt{filter} p l).
  \item - now apply \texttt{partition_filter_fst'} with f.
  \item - assert (Hass := \texttt{filter_in} p a l).
    \item apply Hass.
    \item now rewrite \texttt{← Hf}.
\end{itemize}

Qed.
Lemma part_snd_false : ∀ A p (x t f : list A),
    partition p x = (t, f) →
    (∀ a, ln a f → p a = false).

Proof.
    intros A p l t f Hpart a Hin.
    assert (Hf: f = filter (neg p) l).
    - now apply partition_filter_snd' with t.
    - assert (Hass := filter_in (neg p) a l).
        rewrite ← neg_false_true in Hass.
        apply Hass.
        now rewrite ← Hf.
Qed.

Next is a rather obvious but useful lemma, which states that if a list l was split into (t, f) then appending these lists back together results in a list that is a permutation of the original.

Lemma partition_Permutation : ∀ {A} p (l t f:list A),
    partition p l = (t, f) →
    Permutation l (t ++ f).
Proof.
    intros A p l. induction l; intros.
    - simpl in H. inversion H. auto.
    - simpl in H. destruct (partition p l). destruct (p a); inversion H.
        + simpl. apply perm_skip. apply IHl. f_equal. auto.
        + apply Permutation_trans with (l':=a :: l1 ++ t0). apply perm_skip.
            apply Permutation_trans with (l':=t0 ++ l1). apply IHl. f_equal.
            auto. apply Permutation_app_comm.
            apply Permutation_app_comm with (l:=a :: l1).
Qed.

The last and hardest fact about partition states that if the list being partitioned was already sorted, then the resulting two lists will also be sorted. This seems simple, as partition iterates through the elements in order and maintains the order in its children, but was surprisingly difficult to prove.

After performing induction, the next step was to destruct f a, to see which of the two lists the induction element would end up in. In both cases, the list that doesn’t receive the new element is already clearly sorted by the induction hypothesis, but proving the other one is sorted is slightly harder.

By using Forall_HdRel (defined earlier), we reduced the problem in both cases to only having to show that the new element holds the relation c between all elements of the list it was consed onto. After some manipulation and the use of partition_Permutation and Forall_incl, this follows from the fact that we know the new element holds the relation between all elements of the original list p, and therefore also holds it between the elements of the partitioned list.
Lemma part_Sorted : \( \forall \{ X \} (c:X \to X \to \text{Prop}) f p, \)
Relations_1.Transitive \( c \to \)
Sorted \( c p \to \)
\( \forall l r, \) partition \( f p = (l, r) \to \)
Sorted \( c l \land \) Sorted \( c r. \)

Proof.
intros \( X c f p Htran Hsort. \) induction \( p; \) intros.
- simpl in \( H. \) inversion \( H. \) auto.
- assert \( (H0 := H); \) auto. simpl in \( H. \) destruct (partition \( f p \)) as \([g d].\)
destruct \( f a); \) inversion \( H. \)
+ assert \( (\text{Forall} (c a) g \land \text{Sorted} c g \land \text{Sorted} c r \to \)
Sorted \( c (a :: g) \land \text{Sorted} c r). \)
\times intros \( H4. \) split. apply Sorted_cons. apply \( H4. \) apply Forall_HdRel.
apply \( H4. \) apply \( H4. \)
\times apply \( H1. \) split.
- apply Sorted_StronglySorted in \( Hsort; \) auto.
  apply StronglySorted_inv in \( Hsort \) as \( ||. \)
  apply (Forall_incl _ _ _ H5). apply partition_Permutation in \( H0. \)
  rewrite \( \leftarrow H2 \) in \( H0. \) simpl in \( H0. \) apply Permutation_cons_inv in \( H0. \)
  apply Permutation_incl in \( H0 \) as \( ||. \) unfold incl. unfold incl in \( H6. \)
  intros \( a0 Hin. \) apply \( H6. \) intuition.
- apply \( IHp. \) apply Sorted_inv in \( Hsort; \) apply \( Hsort. f_equal. \) auto.
+ assert \( (\text{Forall} (c a) d \land \text{Sorted} c l \land \text{Sorted} c d \to \)
Sorted \( c l \land \text{Sorted} c (a :: d)). \)
\times intros \( H4. \) split. apply \( H4. \) apply Sorted_cons. apply \( H4. \)
apply Forall_HdRel. apply \( H4. \)
\times apply \( H1. \) split.
- apply Sorted_StronglySorted in \( Hsort; \) auto.
  apply StronglySorted_inv in \( Hsort \) as \( ||. \)
  apply (Forall_incl _ _ _ H5). apply partition_Permutation in \( H0. \)
  rewrite \( \leftarrow H3 \) in \( H0. \) simpl in \( H0. \)
  apply Permutation_trans with \( (l’ := a :: d ++ l) \) in \( H0. \)
  apply Permutation_cons_inv in \( H0. \)
  apply Permutation_trans with \( (l’ := l ++ d) \) in \( H0. \)
  apply Permutation_incl in \( H0 \) as \( ||. \) unfold incl. unfold incl in \( H6. \)
  intros \( a0 Hin. \) apply \( H6. \) intuition. apply Permutation_app_comm.
  apply Permutation_app_comm with \( (l’ := a :: d). \)
- apply \( IHp. \) apply Sorted_inv in \( Hsort; \) apply \( Hsort. f_equal. \) auto.
Qed.
5.4 New Functions over Lists

In order to easily perform the operations we need on lists, we defined three major list functions of our own, each with their own proofs. These generalized list functions all help to make it much easier to deal with our polynomial and monomial lists later in the development.

5.4.1 Distributing two Lists: distribute

The first and most basic of the three is distribute. Similarly to the “FOIL” technique learned in middle school for multiplying two polynomials, this function serves to create every combination of one element from each list. It is done concisely with the use of higher order functions below.

Definition distribute \{A\} (l m : list (list A)) : list (list A) :=
concat (map (fun a ⇒ map (app a) l) m).

The distribute function will play a larger role later, mostly as a part of our polynomial multiplication function. For now, however, there are only two very simple lemmas to be proven, both stating that distributing nil over a list results in nil.

Lemma distribute-nil : ∀ \{A\} (l:list (list A)),
distribute [] l = [].
Proof.
induction l; auto.
Qed.

Lemma distribute-nil_r : ∀ \{A\} (l:list (list A)),
distribute l [] = [].
Proof.
induction l; auto.
Qed.

5.4.2 Cancelling out Repeated Elements: nodup-cancel

The next list function, and possibly the most prolific function in our entire development, is nodup-cancel. Similarly to the standard library nodup function, nodup-cancel takes a list that may or may not have duplicates in it and returns a list without duplicates.

The difference between ours and the standard function is that rather than just removing all duplicates and leaving one of each element, the elements in a nodup-cancel list cancel out in pairs. For example, the list \[1;1;1\] would become \[1\], whereas \[1;1;1;1\] would become \[].

This is implemented with the count_occ function and remove, and is largely the reason for needing so many lemmas about those two functions. If there is an even number of occurrences of an element \(a\) in the original list \(a :: l\), which implies there is an odd number of occurrences of this element in \(l\), then all instances are removed. On the other hand, if there is an odd number of occurrences in the original list, one occurrence is kept, and the rest are removed.
By calling \texttt{nodup\_cancel} recursively on \texttt{xs} before calling \texttt{remove}, Coq is easily able to determine that \texttt{xs} is the decreasing argument, removing the need for a more complicated definition with ‘fuel’.

\begin{verbatim}
Fixpoint nodup\_cancel \{ A \} Aeq \_dec \( l \):list \( A \) :=
    match \( l \) with
    | [] \Rightarrow []
    | x :: xs \Rightarrow
        let count := count\_occ Aeq\_dec \( xs \) \( x \) in
        let \( xs' \) := remove Aeq\_dec \( x \) (nodup\_cancel Aeq\_dec \( xs \)) in
        if even count then x :: \( xs' \) else \( xs' \)
    end.

Now onto lemmas. To begin with, there are a few facts true of \texttt{nodup} that are also true of \texttt{nodup\_cancel}, which are useful in many proofs. \texttt{nodup\_cancel\_in} is the same as the standard library’s \texttt{nodup\_in}, with one important difference: this implication is \textit{not} bidirectional. Because even parity elements are removed completely, not all elements in \( l \) are guaranteed to be in \texttt{nodup\_cancel} \( l \).

\texttt{NoDup\_nodup\_cancel} is much simpler, and effectively exactly the same as \texttt{NoDup\_nodup}.

In these proofs, and most others from this point on, the shape will be very similar to the proof of the corresponding \texttt{nodup} proof. The main difference is that, instead of destructing \texttt{in\_dec} like one would for \texttt{nodup}, we destruct the evenness of \texttt{count\_occ}, as that is what drives the main if statement of the function.

\begin{verbatim}
Lemma nodup\_cancel\_in : \forall A Aeq \_dec a \( l \):list \( A \),
    \texttt{In} \( a \) (nodup\_cancel Aeq\_dec \( l \)) \rightarrow \texttt{In} \( a \) \( l \).
Proof.
    intros A Aeq\_dec a l H. induction \( l \) as [[\( b \) \( l \) IHl]]; auto.
    simpl in \( H \). destruct (Aeq\_dec \( a \) \( b \)).
    - rewrite e. intuition.
    - right. apply IHl. destruct (even (count\_occ Aeq\_dec \( l \) \( b \))).
        + simpl in \( H \). destruct \( H \) rewrite \( H \) in \( n \). contradiction.
            apply \texttt{ln\_remove} in \( H \). auto.
        + apply \texttt{ln\_remove} in \( H \). auto.
Qed.

Lemma NoDup\_nodup\_cancel : \forall A Aeq \_dec \( l \):list \( A \),
\texttt{NoDup} (nodup\_cancel Aeq\_dec \( l \)).
Proof.
    induction \( l \) as [[\( a \) \( l' \) Hrec]]; simpl.
    - constructor.
    - destruct (even (count\_occ Aeq\_dec \( l' \) \( a \))). simpl.
        + apply NoDup\_cons; [apply remove\_ln | apply NoDup\_remove; auto].
        + apply NoDup\_remove; auto.
Qed.
\end{verbatim}
Although not standard library lemmas, the `no_nodup_NoDup` and `Sorted_nodup` facts we proved earlier in this file are also both true of `nodup_cancel`, and proven in almost the same way.

**Lemma no_nodup_cancel_NoDup**: ∀ A Aeq_dec (l:list A),
\[ \text{NoDup } l \rightarrow \text{nodup_cancel } Aeq_dec \ l = l. \]

**Proof**.
intros A Aeq_dec l H. induction l; auto.
simpl. apply NoDup_cons_iff in H as \[. \]
assert (count_occ Aeq_dec l a = 0).
- apply count_occ_not_in. auto.
- rewrite H1. simpl. f_equal. rewrite not_in_remove. auto. intro.
  apply nodup_cancel_in in H2. apply H. auto.
Qed.

**Lemma Sorted_nodup_cancel**: ∀ A Aeq_dec Rel (l:list A),
\[ \text{Relations.1.Transitive } Rel \rightarrow \text{Sorted } Rel \ l \rightarrow \text{Sorted } Rel \ (\text{nodup_cancel } Aeq_dec \ l). \]

**Proof**.
intros A Aeq_dec Rel l Ht H. apply Sorted_StronglySorted in H; auto.
apply StronglySorted_Sorted. induction l; auto.
simpl. apply StronglySorted_inv in H as \[. \]
destruct (even (count_occ Aeq_dec l a)).
- apply SSSorted_cons.
  + apply StronglySorted_remove. apply IHl. apply H.
  + apply Forall_remove. apply Forall_all. rewrite Forall_all in H0.
    intros x Hn. apply H0. apply nodup_cancel_in in Hn. auto.
- apply StronglySorted_remove. apply IHl. apply H.
Qed.

An interesting side effect of the “cancelling” behavior of this function is that while the number of occurrences of an item may change after calling `nodup_cancel`, the evenness of the count never will. If an element was odd before there will be one occurrence, and if it was even before there will be none.

**Lemma count_occ_nodup_cancel**: ∀ \{ A Aeq \} p (a:A),
\[ \text{even } (\text{count_occ } Aeq_dec \ (\text{nodup_cancel } Aeq_dec \ p) \ a) = \text{even } (\text{count_occ } Aeq_dec \ p \ a). \]

**Proof**.
intros A Aeq_dec p a. induction p as \[. \]
destruct (even (count_occ Aeq_dec p b)) eqn:Hb.
- simpl. destruct (Aeq_dec b a).
  + rewrite e. rewrite count_occ_remove. rewrite e in Hb.
    repeat rewrite even_suc. rewrite \[. \] negb_odd in Hb.

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rewrite \text{Bool.negb}\_true\_iff in \text{Hb}. rewrite \text{Hb}. auto.
+ rewrite \text{count}\_occ\_neq_remove; auto.
- simpl. destruct \((\text{Aeq}\_\text{dec} b a)\).
+ rewrite \text{e}. rewrite \text{count}\_occ\_remove. rewrite \text{e} in \text{Hb}.
repeat rewrite \text{even}\_\text{succ}. rewrite \leftarrow \text{negb}\_\text{odd} in \text{Hb}.
rewrite \text{Bool.negb}\_\text{false}\_iff in \text{Hb}. rewrite \text{Hb}. auto.
+ rewrite \text{count}\_occ\_neq_remove; auto.
Qed.

The \text{Permutation}\_\text{nodup} lemma was challenging to prove before, and this version for \text{nodup}\_\text{cancel} faces the same problems. The first and fourth cases are easy, and the second isn’t too bad after using \text{count}\_occ\_\text{Permutation}. The third case faces the same problems as before, but requires some extra work when transitioning from reasoning about \text{count}\_occ \((x :: l) y\) to \text{count}\_occ \((y :: l) x\).

This is accomplished by using \text{even}\_\text{succ}, \text{negb}\_\text{odd}, and \text{negb}\_\text{true}\_\text{iff}. In this way, we can convert something saying \text{even} \((S n) = \text{true} to \text{even} n = \text{false}.

\text{Lemma nodup}\_\text{cancel}\_\text{Permutation} : \forall A \text{Aeq} \_\text{dec} \((l l’)\_\text{list} A),
\text{Permutation} l l’ \rightarrow
\text{Permutation} \text{(nodup}\_\text{cancel} A \text{Aeq} \_\text{dec} l) \text{(nodup}\_\text{cancel} A \text{Aeq} \_\text{dec} l’).
\text{Proof}.
intros A \text{Aeq} \_\text{dec} l l’ H. induction H.
- auto.
- simpl. destruct even \text{eqn:Hevn}.
  + rewrite \text{(count}\_occ\_Permutation \_ \_ \_ \_ \_ H) in \text{Hevn}. rewrite \text{Hevn}.
    apply \text{perm}\_\text{skip}. apply \text{remove}\_\text{Permutation}. apply \text{IHPermutation}.
  + rewrite \text{(count}\_occ\_\text{Permutation} \_ \_ \_ \_ \_ H) in \text{Hevn}. rewrite \text{Hevn}.
    apply \text{remove}\_\text{Permutation}. apply \text{IHPermutation}.
- simpl. destruct \((\text{even} \text{(count}\_occ \text{Aeq}\_\text{dec} l x)) \text{eqn:Hevx}\);
  destruct \((\text{even} \text{(count}\_occ \text{Aeq}\_\text{dec} l y)) \text{eqn:Hevy}; destruct \((\text{Aeq}\_\text{dec} x y)\).
  + rewrite \text{even}\_\text{succ}. rewrite \leftarrow \text{negb}\_\text{odd} in \text{Hevy}.
    rewrite \text{Bool.negb}\_\text{true}\_\text{iff} in \text{Hevy}. rewrite \text{Hevy}. destruct \((\text{Aeq}\_\text{dec} y x)\);
    \text{try (rewrite e in n; contradiction). rewrite even}\_\text{succ}.
    rewrite \leftarrow \text{negb}\_\text{odd} in \text{Hevx}. rewrite \text{Bool.negb}\_\text{true}\_\text{iff} in \text{Hevx}.
    rewrite \text{Hevx}. \text{simpl. destruct} \((\text{Aeq}\_\text{dec} y x)\); \text{try contradiction}.
    \text{destruct} \((\text{Aeq}\_\text{dec} x y)\); \text{try contradiction}. \text{rewrite remove}\_\text{remove}. auto.
  + rewrite \text{Hevy}. \text{simpl. destruct} \((\text{Aeq}\_\text{dec} y x)\);
  \text{try (symmetry in e; contradiction). destruct} \((\text{Aeq}\_\text{dec} x y)\);
  \text{try contradiction}. \text{rewrite Hevx}. \text{rewrite remove}\_\text{remove}. apply \text{perm}\_\text{swap}.
  + rewrite \leftarrow e in \text{Hevy}. rewrite \text{Hevy} in \text{Hevx}. \text{inversion Hevx}.
  + rewrite \text{Hevy}. \text{simpl. destruct} \((\text{Aeq}\_\text{dec} y x)\);
  \text{try (symmetry in e; contradiction). rewrite Hevx. apply perm}\_\text{skip}.
  \text{rewrite remove}\_\text{remove}. auto.
  + rewrite e in \text{Hevx}. rewrite \text{Hevx} in \text{Hevy}. \text{inversion Hevy}.

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As mentioned earlier, in the original definition of the function, it was helpful to reverse the order of \texttt{remove} and the recursive call to \texttt{nodup-cancel}. This is possible because these operations are associative, which is proven below.

\textbf{Lemma \texttt{nodup-cancel-rem-assoc}}: \( \forall \{A\} \ A_{eq\_dec} (a:A) \ p, \)
\[
\texttt{remove} \ A_{eq\_dec} a \ (\texttt{nodup-cancel} \ A_{eq\_dec} p) = \texttt{nodup-cancel} \ A_{eq\_dec} (\texttt{remove} \ A_{eq\_dec} a \ p).
\]

\textbf{Proof.}

\texttt{intros} \(A\ \texttt{A_{eq\_dec}}\ a\ \texttt{p}\); \texttt{induction} \(p\); \texttt{auto}.
\texttt{simpl. destruct} \(\texttt{even eqn:Hevn}\).
- \texttt{simpl. destruct} \((\texttt{A_{eq\_dec}} a\ a0)\).
  + \texttt{rewrite} \(\leftrightarrow\ \texttt{e. rewrite not\_ln\_remove}\); \texttt{auto. apply remove\_ln}.
  + \texttt{simpl. rewrite count\_occ\_neq\_remove}; \texttt{auto. rewrite Hevn}.
    \texttt{f\_equal. rewrite} \(\leftrightarrow\ \texttt{IHp. rewrite remove\_remove}\); \texttt{auto}.
- \texttt{destruct} \((\texttt{A_{eq\_dec}} a\ a0)\).
  + \texttt{rewrite} \(\leftrightarrow\ \texttt{e. rewrite not\_ln\_remove}\); \texttt{auto. apply remove\_ln}.
  + \texttt{simpl. rewrite count\_occ\_neq\_remove}; \texttt{auto. rewrite Hevn}.
    \texttt{rewrite remove\_remove}. \texttt{rewrite} \(\leftrightarrow\ \texttt{IHp. auto}\).

\textbf{Qed.}

The entire point of defining \texttt{nodup-cancel} was so that repeated elements in a list cancel out; clearly then, if an entire list appears twice it will cancel itself out. This proof would be much easier if the order of \texttt{remove} and \texttt{nodup-cancel} was swapped, but the above proof of the two being associative makes it easier to manage.

\textbf{Lemma \texttt{nodup-cancel-self}}: \( \forall \{A\} \ A_{eq\_dec} (l:\texttt{list} A), \)
\[
\texttt{nodup-cancel} \ A_{eq\_dec} (l ++ l) = [\].
\]

\textbf{Proof.}

\texttt{intros} \(A\ \texttt{A_{eq\_dec}}\ p\); \texttt{induction} \(p\); \texttt{auto}.
\texttt{simpl. destruct} \(\texttt{even eqn:Hevn}\).
- \texttt{rewrite count\_occ\_app in Hevn. destruct} \((\texttt{count\_occ} \ A_{eq\_dec} p\ a)\ \texttt{eqn:Hx.} \)
  + \texttt{simpl in Hevn. destruct} \((\texttt{A_{eq\_dec}} a\ a)\); \texttt{try contradiction}.  

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rewrite $Hx$ in $Hevn$. inversion $Hevn$.
+ simpl in $Hevn$. destruct ($Aeq\_dec a a$); try contradiction.
rewrite $Hx$ in $Hevn$. rewrite add\_comm in $Hevn$.
simpl in $Hevn$. destruct (plus $n$ $n$) eqn:Help. inversion $Hevn$.
replace (plus $n$ $n$) with (plus $0$ ($2 \times n$)) in Help.
pose (even\_add\_mul $2$ $n$). pose (even\_succ $n0$). rewrite $\leftarrow$ Help in $e1$.
rewrite $e0$ in $e1$. simpl in $e1$. apply even\_spec in $Hevn$. symmetry in $e1$.
apply odd\_spec in $e1$. apply (Even\_Odd\_False _ $Hevn$) in $e1$. inversion $e1$.
simpl. auto.
- clear $Hevn$. rewrite nodup\_cancel\_remove\_assoc. rewrite remove\_distr\_app.
simpl. destruct ($Aeq\_dec a a$); try contradiction.
rewrite $\leftarrow$ remove\_distr\_app. rewrite $\leftarrow$ nodup\_cancel\_remove\_assoc.
rewrite IH$p$. auto.
Qed.

Next up is a useful fact about $\ln$ that results from nodup\_cancel. Because when there's an even number of an element they all get removed, we can say that there will not be any in the resulting list.

Lemma not\_in\_nodup\_cancel : $\forall \{A Aeq\_dec\} (m:A) p$,
even (count\_occ $Aeq\_dec p m$) $=$ $true$ $\rightarrow$
$\neg \ln m$ (nodup\_cancel $Aeq\_dec p$).
Proof.
intros $A Aeq\_dec m$ $p$ $H$. induction $p$; auto.
intro. simpl in $H$. destruct ($Aeq\_dec a m$).
- simpl in $H0$. rewrite even\_succ in $H$. rewrite $\leftarrow$ negb\_even in $H$.
  rewrite Bool.negb\_true\_iff in $H$. rewrite $\leftarrow e$ in $H$. rewrite $H$ in $H0$.
  rewrite $e$ in $H0$. apply remove\_ln in $H0$. inversion $H0$.
- apply IH$p$; auto. simpl in $H0$. destruct (even (count\_occ $Aeq\_dec p a$)).
  + destruct $H0$; try contradiction. apply ln\_remove in $H0$. auto.
  + apply ln\_remove in $H0$. auto.
Qed.

Similarly to the above lemma, because $a$ will already be removed from $p$ by nodup\_cancel, whether or not a remove is added doesn't make a difference.

Lemma nodup\_extra\_remove : $\forall \{A Aeq\_dec\} (a:A) p$,
even (count\_occ $Aeq\_dec p a$) $=$ $true$ $\rightarrow$
nodup\_cancel $Aeq\_dec p$ $=$
nodup\_cancel $Aeq\_dec$ (remove $Aeq\_dec a p$).
Proof.
intros $A Aeq\_dec a$ $p$ $H$. induction $p$ as $[|b|]$; auto. simpl.
destruct ($Aeq\_dec a b$).
- rewrite $e$ in $H$. simpl in $H$. destruct ($Aeq\_dec b b$); try contradiction.
  rewrite even\_succ in $H$. rewrite $\leftarrow$ negb\_even in $H$.
Lastly, one of the toughest \texttt{nodup-cancel} lemmas. Similarly to \texttt{nodup-pointless}, if \texttt{nodup-cancel} is going to be applied later, there is no need for it to be applied twice. This lemma proves to be very useful when proving that two different polynomials are equal, because, as we will see later, there are often repeated calls to \texttt{nodup-cancel} inside one another. This lemma makes it significantly easier to deal with, as we can remove the redundant \texttt{nodup-cancels}.

This proof proved to be challenging, mostly because it is hard to reason about the parity of the same element in two different lists. In the proof, we begin with induction over \( p \), and then move to destructing the count of \( a \) in each list. The first case follows easily from the two even hypotheses, \texttt{count_occ_app}, and a couple other lemmas. The second case is almost exactly the same, except \( a \) is removed by \texttt{nodup-cancel} and never makes it out front, so the call to \texttt{perm_skip} is removed.

The third case, where \( a \) appears an odd number of times in \( p \) and an even number of times in \( q \), is slightly different, but still solved relatively easily with the use of \texttt{nodup-extra_remove}. The fourth case is by far the hardest. We begin by asserting that, since the count of \( a \) in \( q \) is odd, there must be at least one, and therefore we can rewrite with \texttt{ln_split} to get \( q \) into the form of \( l1 ++ a ++ l2 \). We then assert that, since the count of \( a \) in \( q \) is odd, the count in \( l1 ++ l2 \), or \( q \) with one \( a \) removed, must surely be even. These facts, combined with \texttt{remove_distr_app}, \texttt{count_occ_app}, and \texttt{nodup-cancel_remove_assoc}, allow us to slowly but surely work \( a \) out to the front and eliminate it with \texttt{perm_skip}. All that is left to do at that point is to perform similar steps in the induction hypothesis, so that both \( IH_p \) and our goal are in terms of \( l1 \) and \( l2 \). \( IH_p \) is then used to finish the proof.

\begin{lemma}
\texttt{nodup-cancel_pointless} : \( \forall \{ A \ Aeq\_dec \} \ (p \ q:\text{list} \ A), \)
\begin{align*}
\textbf{Permutation} & \quad (\texttt{nodup-cancel Aeq\_dec} \ (\texttt{nodup-cancel Aeq\_dec} \ p \ ++ \ q)) \\
& \quad (\texttt{nodup-cancel Aeq\_dec} \ (p \ ++ \ q)).
\end{align*}
\end{lemma}

\textbf{Proof}.
\begin{verbatim}
intros A Aeq_dec p q. induction p; auto.
destruct (even (count_occ Aeq_dec p a)) eqn:Hevp;
destruct (even (count_occ Aeq_dec q a)) eqn:Hevq.
- simpl. rewrite Hevp. simpl. rewrite count_occ_app, count_occ_remove. simpl.
  rewrite count_occ_app, even_add, Hevp, Hevq. simpl. apply perm_skip.
  rewrite nodup_cancel_remove_assoc. rewrite remove_pointless.
  rewrite ← nodup_cancel_remove_assoc. apply remove_Permutation. apply IHp.
\end{verbatim}
- simpl. rewrite $\text{Hevp}$. simpl. rewrite $\text{count_occ_app}$, $\text{count_occ_remove}$. simpl. rewrite $\text{count_occ_app}$, $\text{even_add}$, $\text{Hevp}$, $\text{Hevq}$. simpl. rewrite $\leftarrow \text{nodup_cancel_remove_assoc}$. rewrite $\text{remove_pointless}$. rewrite $\leftarrow \text{nodup_cancel_remove_assoc}$. apply $\text{Permutation}$. apply $\text{IH} p$.
- simpl. rewrite $\text{Hevp}$. rewrite $\text{count_occ_app}$, $\text{even_add}$, $\text{Hevp}$, $\text{Hevq}$. simpl. rewrite $\leftarrow \text{(nodup_extra_remove a)}$.
  + rewrite $\text{remove_pointless}$. rewrite $\leftarrow \text{nodup_cancel_remove_assoc}$.
  + rewrite $\text{count_occ_app}$.
  + rewrite $\text{even_add}$. rewrite $\text{count_occ_remove}$.
  rewrite $\text{Hevp}$. auto.
- assert ($\text{count_occ Aeq}_\text{dec} q a > 0$). destruct ($\text{count_occ ~ q ~}$).
  inversion $\text{Hevq}$. apply $\text{gt_Sn_O}$. apply $\text{count_occ_Ln}$ in $H$.
  apply $\text{in_split}$ in $H$ as $\langle l1 \mid l2 \mid H \rangle$. rewrite $H$. simpl $\text{nodup_cancel}$ at 2.
  rewrite $\text{Hevp}$. simpl $\text{app}$. rewrite $H$ in $\text{IH} p$. simpl $\text{nodup_cancel}$ at 3.
  rewrite $\text{count_occ_app}$. rewrite $\text{even_add}$. rewrite $\text{Hevp}$. rewrite $\leftarrow H$ at 2.
  rewrite $\text{Hevp}$. simpl. apply $\text{Permutation}$. rewrite $\leftarrow \text{trans}$ with ($l' := \text{nodup_cancel Aeq}_\text{dec} (a :: \text{remove Aeq}\_\text{dec} a (\text{nodup_cancel Aeq}_\text{dec} p) ++ l1 ++ l2)$).
  + apply $\text{nodup_cancel}\_\text{Permutation}$. rewrite $\text{app_assoc}$. apply $\text{Permutation}\_\text{sym}$.
  rewrite $\text{app_assoc}$. apply $\text{Permutation}\_\text{middle}$ with ($l2 := l2$) ($l1 := \text{remove Aeq}_\text{dec} a (\text{nodup_cancel Aeq}_\text{dec} p) ++ l1$).
  + assert ($\text{even (count_occ Aeq}_\text{dec} (l1 ++ l2) a) = \text{true}$).
    rewrite $H$ in $\text{Hevp}$. rewrite $\text{count_occ_app}$ in $\text{Hevp}$. simpl in $\text{Hevp}$.
    destruct ($\text{Aeq}_\text{dec} a a$); try contradiction. rewrite $\text{plus_comm}$ in $\text{Hevp}$.
    rewrite $\text{plus_Sn_m}$ in $\text{Hevp}$. rewrite $\text{even_succ}$ in $\text{Hevp}$.
    rewrite $\leftarrow \text{negb_even}$ in $\text{Hevp}$. rewrite $\text{Bool}$.$\text{negb_false_iiff}$ in $\text{Hevp}$.
    rewrite $\text{count_occ_app}$. symmetry. rewrite $\text{plus_comm}$. auto.
  simpl. rewrite $\text{count_occ_app}$. rewrite $\text{count_occ_remove}$. simpl.
  replace ($\text{even ~}$) with $\text{true}$. apply $\text{perm}$.$\text{skip}$.
  rewrite $\leftarrow \text{nodup_cancel_remove_assoc ~}$ ($p ++ l1 ++ a :: l2$).
  repeat rewrite $\text{remove_distr_app}$. simpl; destruct ($\text{Aeq}_\text{dec} a a$); try contradiction. rewrite $\text{nodup_cancel_remove_assoc}$.
  rewrite $\text{remove_pointless}$. repeat rewrite $\leftarrow \text{remove_distr_app}$.
  repeat rewrite $\leftarrow \text{nodup_cancel_remove_assoc}$. apply $\text{Permutation}$ with
  ($l' := \text{nodup_cancel Aeq}_\text{dec} (a :: p ++ l1 ++ l2)$) in $\text{IH} p$.
  apply $\text{Permutation}\_\text{sym}$ in $\text{IH} p$.
  apply $\text{Permutation}$ with ($l' := \text{nodup_cancel Aeq}_\text{dec} (a :: \text{nodup_cancel Aeq}_\text{dec} p ++ l1 ++ l2)$) in $\text{IH} p$.
  simpl in $\text{IH} p$. rewrite $\text{count_occ_app}$, $\text{even_add}$, $\text{Hevp}$ in $\text{IH} p$.
  rewrite $\text{IH} 0$ in $\text{IH} p$. simpl in $\text{IH} p$.
  rewrite $\text{count_occ_app}$, $\text{even_add}$, $\text{count_occ_nodup_cancel}$, $\text{Hevp}$, $\text{IH} 0$ in $\text{IH} p$.
  simpl in $\text{IH} p$. apply $\text{Permutation}$; apply $\text{IH} p$.
  × apply $\text{nodup_cancel}\_\text{Permutation}$. rewrite $\text{app_assoc}$.
    apply $\text{Permutation}\_\text{sym}$. rewrite $\text{app_assoc}$. apply $\text{Permutation}\_\text{middle}$ with
(l1 := nodup_cancel Aeq_dec p ++ l1).
\times \text{apply nodup_cancel_Permutation. rewrite app_assoc.}
\text{apply Permutation_sym. rewrite app_assoc. apply Permutation_middle with}
(l1 := p ++ l1).

Qed.

This lemma is simply a reformalization of the above for convenience, which follows simply because of Permutation_app_comm.

Lemma nodup_cancel_pointless_r : \forall \{A\ Aeq\ dec\} (p : list A),
\text{Permutation}
(nodup_cancel Aeq_dec (p ++ nodup_cancel Aeq_dec q))
(nodup_cancel Aeq_dec (p ++ q)).

Proof.
intros A Aeq_dec p q. apply Permutation_trans with (l' := nodup_cancel Aeq_dec
(nodup_cancel Aeq_dec q ++ p)). apply nodup_cancel_Permutation.
\text{apply Permutation_app_comm.}
\text{apply Permutation_sym. apply Permutation_trans with (l' := nodup_cancel}
Aeq_dec (q ++ p)). apply nodup_cancel_Permutation.
\text{apply Permutation_app_comm. apply Permutation_sym.}
\text{apply nodup_cancel_pointless.}

Qed.

An interesting side effect of nodup_cancel_pointless is that now we can show that nodup_cancel almost “distributes” over app. More formally, to prove that the nodup_cancel of two lists appended together is a permutation of nodup_cancel applied to two other lists appended, it is sufficient to show that the first of each and the second of each are permutations after applying nodup_cancel to them individually.

Lemma nodup_cancel_app_Permutation : \forall \{A\ Aeq\ dec\} (a b c d : list A),
\text{Permutation (nodup_cancel Aeq_dec a) (nodup_cancel Aeq_dec b) \rightarrow}
\text{Permutation (nodup_cancel Aeq_dec c) (nodup_cancel Aeq_dec d) \rightarrow}
\text{Permutation (nodup_cancel Aeq_dec (a ++ c)) (nodup_cancel Aeq_dec (b ++ d)).}

Proof.
intros A Aeq_dec a b c d H H0. rewrite \leftarrow (nodup_cancel_pointless a),
\leftarrow (nodup_cancel_pointless b), \leftarrow (nodup_cancel_pointless_r c),
\leftarrow (nodup_cancel_pointless_r d).
apply nodup_cancel_Permutation. apply Permutation_app; auto.

Qed.

5.4.3 Comparing Parity of Lists: parity_match

The final major definition over lists we wrote is parity_match. parity_match is closely related to nodup_cancel, and allows us to make statements about lists being equal after applying nodup_cancel to them. Clearly, if an element appears an even number of times in both lists,
then it won’t appear at all after \texttt{nodup\_cancel}, and if an element appears an odd number of times in both lists, then it will appear once after \texttt{nodup\_cancel}. The ultimate goal of creating this definition is to prove a lemma that if the parity of two lists matches, they are permutations of each other after applying \texttt{nodup\_cancel}.

The definition simply states that for all elements, the parity of the number of occurrences in each list is equal.

\begin{definition}
\textbf{parity\_match} \{A\} \ Aeq\_dec \ (l \ m:\text{list} \ A) : \text{Prop} := \\
\forall x, \text{even} (\text{count\_occ} \ Aeq\_dec \ l \ x) = \text{even} (\text{count\_occ} \ Aeq\_dec \ m \ x).
\end{definition}

A useful lemma in working towards this proof is that if the count of every variable in a list is even, then there will be no variables in the resulting list. This is relatively easy to prove, as we have already proven \texttt{not\_in\_nodup\_cancel} and can contradict away the other cases.

\begin{lemma}
\textbf{even\_nodup\_cancel} : \forall \{A \ Aeq\_dec\} \ (p:\text{list} \ A), \\
(\forall x, \text{even} (\text{count\_occ} \ Aeq\_dec \ p \ x) = \text{true}) \rightarrow \\
(\forall x, \neg \text{ln} \ x \ (\text{nodup\_cancel} \ Aeq\_dec \ p)).
\end{lemma}

\begin{proof}
\text{intros} \ A \ Aeq\_dec \ p \ H \ m. \ \text{intro. induction} \ p.
- \text{inversion} \ H0.
- \text{simpl in *}. \ \text{pose} \ (H \ m) \ as \ H1. \ \text{symmetry in} \ H1. \ \text{destruct} \ (Aeq\_dec \ a \ m).
  + \text{symmetry in} \ H1. \ \text{rewrite} \ e \ in \ H1. \ \text{rewrite} \ even\_succ \ in \ H1.
  \text{rewrite} \ negb\_even \ in \ H1. \ \text{rewrite} \ Boolean.negb\_true\_iff \ in \ H1.
  \text{rewrite} \ H1 in \ H0. \ \text{rewrite} \ e \ in \ H0. \ \text{apply} \ \text{remove\_ln} \ in \ H0. \ \text{inversion} \ H0.
  + \text{destruct} \ (even (\text{count\_occ} \ Aeq\_dec \ p \ a)).
  \times \ \text{destruct} \ H0; \ \text{try contradiction}. \ \text{apply} \ \text{ln\_remove} \ in \ H0. \ \text{symmetry in} \ H1.
  \text{apply} \ \text{not\_in\_nodup\_cancel} \ in \ H1. \ \text{contradiction}.
  \times \ \text{apply} \ \text{ln\_remove} \ in \ H0. \ \text{symmetry in} \ H1. \ \text{apply} \ \text{not\_in\_nodup\_cancel} \ in \ H1.
  \text{contradiction}.
\end{proof}

\text{Qed.}

The above lemma can then be used in combination with \texttt{nothing\_in\_empty} to easily prove \texttt{parity\_match\_empty}, which will be useful in two cases of our goal lemma.

\begin{lemma}
\textbf{parity\_match\_empty} : \forall \{A \ Aeq\_dec\} \ (q:\text{list} \ A), \\
\text{parity\_match} \ Aeq\_dec \ [] \ q \rightarrow \\
\textbf{Permutation} \ [] \ (\text{nodup\_cancel} \ Aeq\_dec \ q).
\end{lemma}

\begin{proof}
\text{intros} \ A \ Aeq\_dec \ q \ H. \ \text{unfold} \ \text{parity\_match} \ in \ H. \ \text{simpl in} \ H.
\text{symmetry in} \ H. \ \text{pose} \ (even\_nodup\_cancel \ q \ H). \ \text{apply} \ \text{nothing\_in\_empty} \ in \ n.
\text{rewrite} \ n. \ \text{auto}.
\end{proof}

\text{Qed.}

The \texttt{parity\_match} definition is also reflexive, symmetric, and transitive, and knowing this will make future proofs easier.

\begin{lemma}
\textbf{parity\_match\_refl} : \forall \{A \ Aeq\_dec\} \ (l:\text{list} \ A),
\end{lemma}
Lemma `parity_match`: \[ \forall \{A \text{ Aeq } \text{ dec}\} \ (l m:\text{list } A), \]
\[ \text{parity_match \ Aeq } \text{ dec } l m \leftrightarrow \text{parity_match \ Aeq } \text{ dec } m l. \]
Proof.
\[ \text{intros } l m. \text{ unfold parity_match. split; intros } H x; \text{ auto.} \]
Qed.

There are also a few interesting facts that can be proved about elements being consed onto lists in a `parity_match`. First is that if the parity of two lists is equal, then the parities will also be equal after adding another element to the front, and vice versa.

Lemma `parity_match_cons`: \[ \forall \{A \text{ Aeq } \text{ dec}\} \ (a:A) \ l1 l2, \]
\[ \text{parity_match \ Aeq } \text{ dec } (a :: l1) (a :: l2) \leftrightarrow \text{parity_match \ Aeq } \text{ dec } l1 l2. \]
Proof.
\[ \text{intros } A \text{ Aeq } \text{ dec } a l1 l2. \text{ unfold parity_match. split; intros } H x. \]
\[- \text{ pose } (H x). \text{ symmetry in } e. \text{ simpl in } e. \text{ destruct } (\text{Aeq } \text{ dec } a x); \text{ auto.} \]
\[- \text{ repeat rewrite even succ in } e. \text{ repeat rewrite } \leftrightarrow \neg \text{negb even in } e. \]
\[- \text{ apply Bool.negb_sym in } e. \text{ rewrite Bool.negb_involutive in } e. \text{ auto.} \]
\[- \text{ simpl. destruct } (\text{Aeq } \text{ dec } a x); \text{ auto.} \]
\[- \text{ repeat rewrite even succ. repeat rewrite } \leftrightarrow \neg \text{negb even.} \]
\[- \text{ apply Bool.negb_sym. rewrite Bool.negb_involutive. auto.} \]
Qed.

Similarly, adding the same element twice to a list does not change the parities of any elements in the list.

Lemma `parity_match_double`: \[ \forall \{A \text{ Aeq } \text{ dec}\} \ (a:A) \ l, \]
\[ \text{parity_match \ Aeq } \text{ dec } (a :: a :: l) l. \]
Proof.
\[ \text{intros } A \text{ Aeq } \text{ dec } a l. \text{ unfold parity_match. intros } x. \text{ simpl.} \]
\[ \text{destruct } (\text{Aeq } \text{ dec } a x); \text{ auto.} \]
Qed.
The last `cons parity_match` lemma states that if you remove an element from one list and add it to the other, the parity will not be affected. This follows because if they both had an even number of \( a \) before they will both have an odd number after, and if it was odd before it will be even after.

**Lemma parity_match_cons_swap :** \( \forall \{ A \ \text{Aeq}_{\text{dec}} \} (a:A) \ l1 \ l2, \)

- `parity_match Aeq_{\text{dec}} (a :: l1) l2 \rightarrow`
- `parity_match Aeq_{\text{dec}} l1 (a :: l2)`.

**Proof.**
- `intros A Aeq_{\text{dec}} a l1 l2 H`. `apply (parity_match_cons a) in H.`
- `apply parity_match_sym in H`. `apply parity_match_trans with (r:=l1) in H.`
- `apply parity_match_sym in H`. `auto. apply parity_match_double.`

**Qed.**

This next lemma states that if we know that some element \( a \) appears in the rest of the list an even number of times, than clearly it appears in \( l2 \) an odd number of times and must be in the list.

**Lemma parity_match_In :** \( \forall \{ A \ \text{Aeq}_{\text{dec}} \} (a:A) \ l1 \ l2, \)

- `even (count_occ Aeq_{\text{dec}} l1 a) = true \rightarrow`
- `parity_match Aeq_{\text{dec}} (a :: l1) l2 \rightarrow`
- `\( l1 a \ l2 \)`.

**Proof.**
- `intros A Aeq_{\text{dec}} a l1 l2 H H0`. `apply parity_match_cons_swap in H0.`
- `rewrite H0 in H`. `simpl in H`. `destruct (Aeq_{\text{dec}} a a); try contradiction.`
- `rewrite even_succ in H`. `rewrite \( \leftarrow \:\text{negb}_{\text{even}} \) in H.`
- `rewrite \( \text{Bool}_{\text{negb}}_{\text{true}}_{\text{iff}} \) in H.`
- `assert (count_occ Aeq_{\text{dec}} l2 a > 0). destruct count_occ. inversion H.`
- `apply gt_Sn_O. apply count_occ_In in H1. auto.`

**Qed.**

The last fact to prove before attempting the big lemma is that if two lists are permutations of each other, then their parities must match because they contain the same elements the same number of times.

**Lemma Permutation parity_match :** \( \forall \{ A \ \text{Aeq}_{\text{dec}} \} (p q:\text{list} \ A), \)

- \( \text{Permutation} \ p \ q \rightarrow \text{parity_match Aeq}_{\text{dec}} p \ q \).

**Proof.**
- `intros A Aeq_{\text{dec}} p q H`. `induction H.`
  - `auto.`
  - `apply parity_match_cons. auto.`
  - `repeat apply parity_match_cons_swap. unfold parity_match. intros x0.`
    - `simpl. destruct Aeq_{\text{dec}}; destruct Aeq_{\text{dec}};`
    - `repeat (rewrite even_succ; rewrite odd_succ); auto.`
  - `apply parity_match_trans with (r:=l1); auto.`

**Qed.**
Finally, the big one. The first three cases are straightforward, especially now that we have already proven `parity_match_empty`. The third case is more complicated. We begin by destructing if \(a\) and \(a0\) are equal. In the case that they are, the proof is relatively straightforward; `parity_match_cons`, `perm_skip`, and `remove_Permutation` take care of it.

In the case that they are not equal, we next destruct if the number of occurrences is even or not. If it is odd, we can use `parity_match_In` and `ln_split` to rewrite \(l2\) in terms of \(a\). From there, we use permutation facts to rearrange \(a\) to be at the front, and the rest of the proof is similar to the proof when \(a\) and \(a0\) are equal.

The final case is when they are not equal and the number of occurrences is even. After using `parity_match_swap`, we can get to a point where we know that \(a\) appears in \(q++a0\) an even number of times. This means that \(a\) will not be in \(q++a0\) after applying `nodup_cancel`, so we can rewrite with `not_In_remove` in the reverse direction to get the two sides of the permutation goal to be more similar. Then, because it is wrapped in `remove a`, we can clearly add an \(a\) on the inside without it having any effect. Then all that is left is to apply `remove_Permutation`, and we end up with a goal matching the induction hypothesis.

This lemma is very powerful, especially when dealing with `nodup_cancel` with functions applied to the elements of a list. This will come into play later in this file.

Lemma `parity_nodup_cancel_Permutation`:
\[
\forall \{A \ Aeq \ dec\} \ (p \ q : \text{list} \ A),
\text{parity_match} \ Aeq \ dec \ p \ q \rightarrow
\text{Permutation} \ (\text{nodup_cancel} \ Aeq \ dec \ p) \ (\text{nodup_cancel} \ Aeq \ dec \ q).
\]

Proof.
- `intros A Aeq dec p q H. generalize dependent q`.
- `induction p; induction q; intros.`
- `auto.`
- `simpl nodup_cancel at 1. apply parity_match_empty. auto.`
- `simpl nodup_cancel at 2. apply Permutation_sym. apply parity_match_empty. apply parity_match_sym. auto.`
- `clear IHq. destruct (Aeq_dec a a0).`
- `+ rewrite e. simpl. rewrite e in H. apply parity_match_cons in H. destruct even eqn:Hev; rewrite H in Hev; rewrite Hev.`
  - `apply perm_skip. apply remove_Permutation. auto.`
  - `apply remove_Permutation. auto.`
- `+ simp nodup_cancel at 1. destruct even eqn:Hev.`
  - `assert (Hev':=Hev).`
  - `apply parity_match_In with (l2:=a0 :: q) in Hev; auto.`
  - `deconstruct Hev. symmetry in H0. contradiction.`
  - `apply ln_split in H0 as [l1 [l2 H0]]. rewrite H0. apply Permutation_sym. apply Permutation_trans with
    (l':=nodup_cancel Aeq_dec (a :: l2 ++ a0 :: l1)).`
  - `apply nodup_cancel_Permutation. rewrite app_comm_cons. apply (Permutation_app_comm). simpl. rewrite H0 in H.`
  - `apply parity_match_trans with (r:=a :: l2 ++ a0 :: l1) in H.`
apply parity_match_cons in H. rewrite H in Hev'. rewrite Hev'.
apply perm_skip. apply remove_Permutation. apply Permutation_sym.
apply IHp. auto. rewrite app_comm_cons. apply Permutation_parity_match.
apply Permutation_app_comm.
× apply parity_match_cons_swap in H. rewrite H in Hev. assert (Hev2:=Hev).
rewrite count_occ_Permutation with (l':=a :: q ++ [a0]) in Hev.
simpl in Hev. destruct (Aeq_dec a a); try contradiction.
rewrite even_succ in Hev. rewrite ← negb_even in Hev.
rewrite Bool.negb_false_iff in Hev.
rewrite ← (not_in_remove _ Aeq_dec a).
assert (∀ l, remove Aeq_dec a (nodup_cancel Aeq_dec l) =
  remove Aeq_dec a (nodup_cancel Aeq_dec (a :: l))).
intros l. simpl. destruct (even (count_occ _ l a)).
simpl. destruct (Aeq_dec a a); try contradiction.
rewrite (not_in_remove _ _ _ (remove _ _)). auto. apply remove_in.
rewrite (not_in_remove _ _ _ (remove _ _)). auto. apply remove_in.
rewrite (H0 (a0 :: q)). apply remove_Permutation. apply IHp. auto.
apply not_in_nodup_cancel.
rewrite count_occ_Permutation with (l':=a0 :: q) in Hev.
auto. replace (a0 :: q) with ([a0] ++ q); auto.
apply Permutation_app_comm. apply perm_skip.
replace (a0 :: q) with ([a0] ++ q); auto. apply Permutation_app_comm.
Qed.

5.5 Combining nodup_cancel and Other Functions

5.5.1 Using nodup_cancel over map

Our next goal is to prove things about the relation between nodup_cancel and map over lists. In particular, we want to prove a lemma similar to nodup_cancel_pointless, that allows us to remove redundant nodup_cancels.

The challenging part of proving this lemma is that it is often hard to reason about how, for example, the number of times a appears in p relates to the number of times f a appears in map f p. Many of the functions we map across lists in practice are not one-to-one, meaning that there could be some b such that f a = f b. However, at the end of the day, these repeated elements will cancel out with each other and the parities will match, hence why parity_nodup_cancel_Permutation is extremely useful.

To begin, we need to prove a couple facts comparing the number of occurrences of elements in a list. The first lemma states that the number of times some a appears in p is less than or equal to the number of times f a appears in map f p.

Lemma count_occ_map_lt : ∀ {A Aeq_dec} p (a:A) f,
  count_occ Aeq_dec p a ≤ count_occ Aeq_dec (map f p) (f a).
Proof.
  intros A Aeq_dec p a f. induction p. auto. simpl. destruct Aeq_dec.
  - rewrite e. destruct Aeq_dec; try contradiction. simpl. apply le_n_S. auto.
  - destruct Aeq_dec; auto.
Qed.

Building off this idea, the next lemma states that the number of times \( f \, a \) appears in \( \text{map} \, f \, p \) with \( a \) removed is equal to the count of \( f \, a \) in \( \text{map} \, f \, p \) minus the count of \( a \) in \( p \).

**Lemma count_occ_map_sub** : \( \forall \{ A \, \text{Aeq} \} \, f \, (a:A) \, p, \)
\[
\text{count}_\text{occ} \, \text{Aeq} \, \text{dec} \, (\text{map} \, f \, (\text{remove} \, \text{Aeq} \, \text{dec} \, a \, p)) \, (f \, a) = \\
\text{count}_\text{occ} \, \text{Aeq} \, \text{dec} \, (\text{map} \, f \, p) \, (f \, a) \, - \, \text{count}_\text{occ} \, \text{Aeq} \, \text{dec} \, p \, a.
\]

**Proof.**
  intros A Aeq_dec f a p. induction p; auto. simpl. destruct Aeq_dec.
  - rewrite e. destruct Aeq_dec; try contradiction. destruct Aeq_dec; try contradiction. simpl. rewrite ← e. auto.
  - simpl. destruct Aeq_dec.
    + destruct Aeq_dec. symmetry in e; contradiction. rewrite IHp.
      rewrite sub_succ_l. auto. apply count_occ_map_lt.
    + destruct Aeq_dec. symmetry in e; contradiction. auto.
Qed.

It is also true that if there is some \( x \) that is not equal to \( f \, a \), then the count of that \( x \) in \( \text{map} \, f \, p \) is the same as the count of \( x \) in \( \text{map} \, f \, p \) with \( a \) removed.

**Lemma count_occ_map_neq_remove** : \( \forall \{ A \, \text{Aeq} \} \, f \, (a:A) \, p \, x, \)
\[
x \neq f \, a \rightarrow \\
\text{count}_\text{occ} \, \text{Aeq} \, \text{dec} \, (\text{map} \, f \, (\text{remove} \, \text{Aeq} \, \text{dec} \, a \, p)) \, x = \\
\text{count}_\text{occ} \, \text{Aeq} \, \text{dec} \, (\text{map} \, f \, p) \, x.
\]

**Proof.**
  intros. induction p as ||b|; auto. simpl. destruct (Aeq_dec a b).
  - destruct Aeq_dec. rewrite ← e in e0. symmetry in e0. contradiction. auto.
  - simpl. destruct Aeq_dec; auto.
Qed.

The next lemma is similar to count_occ_map_lt, except it involves some \( b \) where \( a \) is not equal to \( b \), but \( f \, a = f \, b \). Then clearly, the sum of \( a \) in \( p \) and \( b \) in \( p \) is less than the count of \( f \, a \) in \( \text{map} \, f \, p \).

**Lemma f_equal_sum_lt** : \( \forall \{ A \, \text{Aeq} \} \, f \, (a:A) \, b \, p, \)
\[
b \neq a \rightarrow (f \, a) = (f \, b) \rightarrow \\
\text{count}_\text{occ} \, \text{Aeq} \, \text{dec} \, p \, b + \\
\text{count}_\text{occ} \, \text{Aeq} \, \text{dec} \, p \, a \leq \\
\text{count}_\text{occ} \, \text{Aeq} \, \text{dec} \, (\text{map} \, f \, p) \, (f \, a).
\]

**Proof.**
  intros A Aeq_dec f a b p Hne Hfe. induction p as ||c|; auto. simpl.
destruct Aeq_dec.
- rewrite e. destruct Aeq_dec; try contradiction. rewrite Hfe.
   destruct Aeq_dec; try contradiction. simpl. apply le_n_S.
   rewrite ← Hfe. auto.
- destruct Aeq_dec.
  + rewrite e. destruct Aeq_dec; try contradiction. rewrite plus_comm.
    simpl. rewrite plus_comm. apply le_n_S. auto.
  + destruct Aeq_dec.
    × apply le_S. auto.
    × auto.
Qed.

For the next lemma, we once again try to compare the count of a to the count of f a, but also involve nodup_cancel. Clearly, there is no way for there to be more a’s in p than f a’s in map f p even with the addition of nodup_cancel.

Lemma count_occ_nodup_map_lt : \forall \{A Aeq dec\} p f (a:A),
  count_occ Aeq_dec (nodup_cancel Aeq_dec p) a ≤
  count_occ Aeq_dec (map f (nodup_cancel Aeq_dec p)) (f a).
Proof.
intros A Aeq_dec p f a. induction p as [||b]; auto. simpl.
destruct even eqn:Hev.
- simpl. destruct Aeq_dec.
  + rewrite e. destruct Aeq_dec; try contradiction. apply le_n_S. auto.
  rewrite count_occ_remove. apply le_0_l.
  + rewrite count_occ_neq_remove; auto. rewrite not_IN_remove.
    destruct Aeq_dec; firstorder. apply not_in_nodup_cancel; auto.
- destruct (Aeq_dec b a) eqn:Hba.
  + rewrite e. rewrite count_occ_remove. apply le_0_l.
  + rewrite count_occ_neq_remove; auto.
    destruct (Aeq_dec (f b) (f a)) eqn:Hfb.
    × rewrite ← e. rewrite count_occ_map_sub. rewrite e.
      apply le_add_le_sub_l. apply f_equal_sum_lt; auto.
    × rewrite count_occ_map_neq_remove; auto.
Qed.

All of these lemmas now come together for the core one, a variation of nodup_cancel_pointless but involving map f. We begin by applying parity_nodup_cancel_Permutation, and destructing if a appears in p an even number of times or not.

The even case is relatively easy to prove, and only involves using the usual combination of even_succ, not_IN_remove, and not_in_nodup_cancel.

The odd case is trickier, and where we involve all of the newly proved lemmas. If x and f a are not equal, the proof follows just from count_occ_map_neq_remove and the induction hypothesis.

If they are equal, we begin by rewriting with count_occ_map_sub and even_sub. After a
few more rewrites, it becomes the case that we need to prove that the boolean equivalence of the parities of $f \ a$ in $\text{map} \ f \ p$ and $a$ in $p$ is equal to the negated parity of $f \ a$ in $\text{map} \ f \ p$. Because we know that $a$ appears in $p$ an odd number of times from destructing even earlier, this follows immediately.

\textbf{Lemma nodup\textunderscore cancel\_map} : $\forall \ \{A \ \text{Aeq} \ \text{dec}\} (p:\text{list} \ A) \ f,$
\begin{align*}
\text{Permutation} & \\
(\text{nodup\textunderscore cancel} \ \text{Aeq} \ \text{dec} \ (\text{map} \ f \ (\text{nodup\textunderscore cancel} \ \text{Aeq} \ \text{dec} \ p))) & \\
(\text{nodup\textunderscore cancel} \ \text{Aeq} \ \text{dec} \ (\text{map} \ f \ p)).
\end{align*}
\textbf{Proof}.
\begin{align*}
\text{intros} \ A \ \text{Aeq} \ \text{dec} \ p \ f. \ & \text{apply parity\textunderscore nodup\textunderscore cancel\_Permutation}. \\
\text{unfold parity\textunderscore match. intros} \ x. \ & \text{induction} \ p; \text{auto. simpl.} \\
\text{destruct} \ (\text{even} \ (\text{count\textunderscore occ} \ _ \ p \ a)) \ & \text{eqn:Hev.} \\
- \ & \text{simpl. destruct} \ \text{Aeq} \ \text{dec}. \\
+ \ & \text{repeat rewrite} \ \text{even\_succ}. \ & \text{repeat rewrite} \leftrightarrow \ \text{negb\_even}. \\
\ & \text{rewrite} \ \text{not\_ln\_remove. rewrite} \ \text{IHp. auto. apply not\_in\_nodup\textunderscore cancel. auto.} \\
+ \ & \text{rewrite not\_ln\_remove. apply} \ \text{IHp. apply not\_in\_nodup\textunderscore cancel. auto.} \\
- \ & \text{simpl. destruct} \ \text{Aeq} \ \text{dec}. \\
+ \ & \text{rewrite} \leftrightarrow \ e. \ & \text{rewrite} \ \text{count\textunderscore occ\_map\_sub. rewrite} \ \text{even\_sub}. \\
\ & \text{rewrite} \leftrightarrow \ e \ \text{in} \ \text{IHp. rewrite} \ \text{IHp. rewrite} \ \text{count\textunderscore occ\_nodup\textunderscore cancel}. \\
\ & \text{rewrite} \ \text{Hev. rewrite} \ \text{even\_succ. rewrite} \leftrightarrow \ \text{negb\_even}. \\
\ & \text{destruct} \ (\text{even} \ (\text{count\textunderscore occ} \ _ \ (\text{map} \ f \ p) \ _)) \ \text{eqn:Hev.} \\
\ & \text{apply} \ \text{count\textunderscore occ\_nodup\_map\_lt.} \\
+ \ & \text{rewrite} \ \text{count\textunderscore occ\_map\_neq\_remove}; \text{auto.}
\end{align*}
\textbf{Qed.}

\textbf{5.5.2 Using nodup\textunderscore cancel over concat map}

Similarly to map, the same property of not needing repeated nodup\textunderscore cancels applies when the lists are being flattened and mapped over. This final section of the file seeks to, in very much the same way as earlier, prove this.

We begin with a simple lemma about math that will come into play soon - if a number is less than or equal to 1, then it is either 0 or 1. This is immediately solved with firstorder logic.

\textbf{Lemma n\textunderscore le\_1} : $\forall \ n,$
\begin{align*}
 n \leq 1 & \rightarrow \ n = 0 \lor \ n = 1.
\end{align*}
\textbf{Proof}.
\begin{align*}
\text{intros} \ n \ \text{H}. \ & \text{induction} \ n; \text{firstorder.}
\end{align*}
\textbf{Qed.}

The main difference between this section and the section about map is that all of the functions being mapped will clearly be returning lists as their output, and then being concatenated with the rest of the result. This makes things slightly harder, as we can’t reason
about the number of times, for example, some \( f \) \( a \) appears in a list. Instead, we have to reason about the number of times that some \( x \) appears in a list, where \( x \) is one of the elements of the list \( f \) \( a \).

In practice, these lemmas are only going to be applied in situations where every \( f \) \( a \) has no duplicates in it. In other words, as the lemma above states, there will be either 0 or 1 of each \( x \) in a list. The next two lemmas prove some consequences of this.

First is that if the count of \( x \) in \( f \) \( a \) is 0, then clearly removing \( a \) from some list \( p \) will not affect the count of \( x \) in the concatenated version of the list.

**Lemma count_occ_map_sub_not_in:** \( \forall \{ A \ Aeq_dec \} f \ (a:A) \ p, \)
\[ \forall \ x, \ \text{count}_\text{occ} \ Aeq\_dec \ (f \ a) \ x = 0 \rightarrow \]
\[ \text{count}_\text{occ} \ Aeq\_dec \ (\text{concat} \ (\text{map} \ f \ (\text{remove} \ Aeq\_dec \ a \ p))) \ x = \]
\[ \text{count}_\text{occ} \ Aeq\_dec \ (\text{concat} \ (\text{map} \ f \ p)) \ x. \]

Proof.
- intros \( A \ Aeq_dec \ f \ a \ p \ x \ H \). induction \( p \) as \([|b|] \); auto. simpl.
- rewrite \( \text{count}_\text{occ}_\text{app} \). destruct \( Aeq\_dec \).
  - rewrite \( e \) in \( H \). rewrite \( H \). firstorder.
  - simpl. rewrite \( \text{count}_\text{occ}_\text{app} \). auto.
Qed.

On the other hand, if the count of some \( x \) in \( f \) \( a \) is 1, then the count of \( a \) in the original list must be less than or equal to the count of \( x \) in the final list, depending on if some \( b \) exists such that \( f \) \( a \) also contains \( x \). More useful is the fact that if \( x \) appears once in \( f \) \( x \), the count of \( x \) in the final list with \( a \) removed is equal to the count of \( x \) in the final list minus the count of \( a \) in the list. Both of these proofs are relatively straightforward, and mostly follow from firstorder logic.

**Lemma count_occ_concat_map_lt:** \( \forall \{ A \ Aeq_dec \} p \ (a:A) \ f \ x, \)
\[ \text{count}_\text{occ} \ Aeq\_dec \ p \ a \leq \text{count}_\text{occ} \ Aeq\_dec \ (\text{concat} \ (\text{map} \ f \ p)) \ x. \]

Proof.
- intros \( A \ Aeq_dec \ p \ a \ f \ x \ H \). induction \( p \). auto. simpl. destruct \( Aeq\_dec \).
  - rewrite \( e \). rewrite \( \text{count}_\text{occ}_\text{app} \). rewrite \( H \). simpl. firstorder.
  - rewrite \( \text{count}_\text{occ}_\text{app} \). induction \( (\text{count}_\text{occ} \ Aeq\_dec \ (f \ a0) \ x) \); firstorder.
Qed.

**Lemma count_occ_map_sub_in:** \( \forall \{ A \ Aeq_dec \} f \ (a:A) \ p, \)
\[ \forall \ x, \ \text{count}_\text{occ} \ Aeq\_dec \ (f \ a) \ x = 1 \rightarrow \]
\[ \text{count}_\text{occ} \ Aeq\_dec \ (\text{concat} \ (\text{map} \ f \ (\text{remove} \ Aeq\_dec \ a \ p))) \ x = \]
\[ \text{count}_\text{occ} \ Aeq\_dec \ (\text{concat} \ (\text{map} \ f \ p)) \ x - \text{count}_\text{occ} \ Aeq\_dec \ p \ a. \]

Proof.
- intros \( A \ Aeq_dec \ f \ a \ p \ x \ H \). induction \( p \) as \([|b|] \); auto. simpl.
  - destruct \( Aeq\_dec \).
  - rewrite \( e \). destruct \( Aeq\_dec \); try contradiction. rewrite \( \text{count}_\text{occ}_\text{app} \).
    - rewrite \( e \) in \( H \). rewrite \( H \). simpl. rewrite \( e \) auto.
- simpl. destruct \textit{Aeq\_dec}. symmetry in \textit{e}. contradiction.
  repeat rewrite count\_occ\_app. rewrite \textit{IHp}. rewrite add\_sub\_assoc. auto.
  apply count\_occ\_concat\_map\_lt; auto.
Qed.

Continuing the pattern of proving similar facts as we did during the \textit{map} proof, we now prove a version of \textit{f\_equal\_sum\_lt} involving \textit{concat}. This lemma states that, if we know there will be no duplicates in \textit{f x} for all \textit{x}, and that there are some \textit{a} and \textit{b} such that they are not equal but \textit{x} in \textit{f a} and \textit{f b}, then clearly the sum of the count of \textit{a} and the count of \textit{b} is less than or equal to the count of \textit{x} in the list after applying the function and flattening.

**Lemma \textit{f\_equal\_concat\_sum\_lt}**: \( \forall \{A \textit{Aeq\_dec}\} f (a:A) b p x, \)
\[
b \neq a \rightarrow \]
\[
(\forall x, \textit{NoDup} (f x)) \rightarrow \]
\[
\text{count\_occ} \textit{Aeq\_dec} (f a) x = 1 \rightarrow \]
\[
\text{count\_occ} \textit{Aeq\_dec} (f b) x = 1 \rightarrow \]
\[
\text{count\_occ} \textit{Aeq\_dec} p b + \]
\[
\text{count\_occ} \textit{Aeq\_dec} p a \leq \]
\[
\text{count\_occ} \textit{Aeq\_dec} (\text{concat} (\text{map} f p)) x. \]

**Proof**.
intros A \textit{Aeq\_dec} f a b p x Hne Hnd Hfa Hfb. induction \textit{p} as \[\mid c\]; auto. simpl.
destruct \textit{Aeq\_dec}.
- rewrite \textit{e}. destruct \textit{Aeq\_dec}; try contradiction. rewrite count\_occ\_app.
  firstorder.
- destruct \textit{Aeq\_dec}.
  + rewrite \textit{e}. rewrite count\_occ\_app. firstorder.
  + rewrite count\_occ\_app. pose (Hnd \textit{c}).
    rewrite (NoDup\_count\_occ \textit{Aeq\_dec}) in n1. pose (n1 x).
    apply n\_le\_1 in l. clear n1. destruct \textit{l}; firstorder.
Qed.

The last step before we are able to prove \textit{nodup\_cancel\_concat\_map} is to actually involve \textit{nodup\_cancel} rather than just \textit{remove}. This lemma states that given \textit{f x} has no duplicates and \textit{a} appears once in \textit{f a}, the count of \textit{a} in \textit{p} after applying \textit{nodup\_cancel} is less than or equal to the count of \textit{x} after applying \textit{concat map} and \textit{nodup\_cancel}.

The first cases, when the count is even, are relatively straightforward. The second cases, when the count is odd, are slightly more complicated. We destruct if \textit{a} and \textit{b} (where \textit{b} is our induction element) are equal. If they are, then the proof is solved by firstorder logic. On the other hand, if they are not, we make use of our \textit{n\_le\_1} fact proved before to find out how many times \textit{x} appears in \textit{f b}. If it is zero, then we rewrite with the 0 fact proved earlier and are done. In the final case, we rewrite with the 1 subtraction fact we proved earlier, and it follows from \textit{f\_equal\_concat\_sum\_lt}.

**Lemma \textit{count\_occ\_nodup\_concat\_map\_lt}**: \( \forall \{A \textit{Aeq\_dec}\} p f (a:A) x, \)
\[
(\forall x, \textit{NoDup} (f x)) \rightarrow \]
count_occ $Aeq_{dec} (f \ a) x = 1 \to$
count_occ $Aeq_{dec} (\text{nodup-cancel} Aeq_{dec} p) a \leq$
count_occ $Aeq_{dec} (\text{concat} (\text{map} f (\text{nodup-cancel} Aeq_{dec} p))) x.$

Proof.

intros $A Aeq_{dec} p f a x Hn H$. induction $p$ as $||b||$; auto. simpl.
destruct even eqn:$Hev$.
- simpl. destruct $Aeq_{dec}$.
  + rewrite $e$. rewrite count_occ_remove, count_occ_app. rewrite $H$. firstorder.
  + rewrite count_occ_neq_remove; auto. rewrite not_in_remove.
    rewrite count_occ_app. firstorder. apply not_in_cancel. auto.
- destruct $(Aeq_{dec} b a)$ eqn:$Hba$.
  + rewrite $e$. rewrite count_occ_remove. firstorder.
  + rewrite count_occ_neq_remove; auto. assert $(Hn1::=(Hn b))$.
    rewrite (NoDup_count_occ $Aeq_{dec}$) in $Hn1$. assert $(Hn2::=(Hn1 x))$.
    clear $Hn1$. apply n_le_1 in $Hn2$. destruct $Hn2$.
      rewrite count_occ_map_sub_not_in; auto.
      rewrite count_occ_map_sub_in _ _ (nodup_cancel $Aeq_{dec} p$) in $H0$ as $H1$.
      rewrite $H1$. apply le_add_le_sub_l. apply f_equal_concat_sum_lt; auto.
Qed.

Finally, the proof we’ve been building up to. Once again, we begin the proof by converting
to a parity_match problem and then perform induction on the list. The case where $a$
appears an even number of times in the list is easy, and follows from the same combination of
count_occ_app and even_add that we have used before.

The case where $a$ appears an odd number of times is slightly more complex. Once again,
we apply n_le_1 to determine how many times our $x$ appears in $f \ a$. If it is zero times, we
use count_occ_map_sub_not_in like above, and then the induction hypothesis solves it. If $x$
appears once in $f \ a$, we instead use count_occ_map_sub_in combined with even_sub. Then,
after rewriting with the induction hypothesis, we can easily solve the lemma with the use of
count_occ_nodup_cancel.

Lemma nodup_cancel_concat_map : $\forall \{A Aeq_{dec}\} (p:list A) f,$
$(\forall x, \text{NoDup} (f x)) \to$

Permutation
  $(\text{nodup-cancel} Aeq_{dec} (\text{concat} (\text{map} f (\text{nodup-cancel} Aeq_{dec} p))))$
  $(\text{nodup-cancel} Aeq_{dec} (\text{concat} (\text{map} f p)))$.

Proof.

intros $A Aeq_{dec} p f H$. apply parity_nodup_cancel_Permutation.
unfold parity_match. intros $x$. induction $p$; auto. simpl.
destruct (even (count_occ _ p a)) eqn:$Hev$.
- simpl. repeat rewrite count_occ_app. repeat rewrite even_add.
  rewrite not_in_remove. rewrite $IHp$. auto. apply not_in_nodup-cancel auto.
  - assert $(H0::=(H a))$. rewrite (NoDup_count_occ $Aeq_{dec}$) in $H0$.
    assert $(H1::=(H0 x))$. clear $H0$. apply n_le_1 in $H1$. rewrite count_occ_app.
rewrite even_add. destruct H1.
+ apply (count_occ_map_sub_not_in _ _ (nodup_cancel Aeq_dec p)) in H0 as H1.
  rewrite H0, H1, IHp. simpl.
  destruct (even (count_occ _ (concat (map f p)) x)); auto.
+ apply (count_occ_map_sub_in _ _ (nodup_cancel Aeq_dec p)) in H0 as H1.
  rewrite H0, H1, even_sub, IHp. simpl. rewrite count_occ_nodup_cancel.
  rewrite Hev. destruct (even (count_occ _ (concat (map f p)) x)); auto.
  apply count_occ_nodup_concat_map_lt; auto.
Qed.
Chapter 6

Library B_Unification.poly

Require Import Arith.
Require Import List.
Import ListNotations.
Require Import FunctionalExtensionality.
Require Import Sorting.
Require Import Permutation.
Import Nat.
Require Export list_util.

6.1 Monomials and Polynomials

6.1.1 Data Type Definitions

Now that we have defined those functions over lists and proven all of those facts about them, we can begin to apply all of them to our specific project of unification. The first step is to define the data structures we plan on using.

As mentioned earlier, because of the ten axioms that hold true during $B$-unification, we can represent all possible terms with lists of lists of numbers. The numbers represent variables, and a list of variables is a monomial, where each variable is multiplied together. A polynomial, then, is a list of monomials where each monomial is added together.

In this representation, the term 0 is represented as the empty polynomial, and the term 1 is represented as the polynomial containing only the empty monomial.

In addition to the definitions of var, mono, and poly, we also have definitions for var_eq_dec and mono_eq_dec; these are a proofs of decidability of varailes and monomials respectively. They make use of a special Coq data structure that allows them to be used as a comparison function - for example, we can destruct (mono_eq_dec a b) to compare the two cases where $a = b$ and $a \neq b$. In addition to being useful in some proofs, this is also needed by some functions, such as remove and count_occ, since they compare variables and monomials.

Definition var := nat.
Definition var_eq_dec := Nat.eq_dec.
Definition mono := list var.
Definition mono_eq_dec := (list_eq_dec Nat.eq_dec).
Definition poly := list mono.

6.1.2 Comparisons of monomials and polynomials

In order to easily compare monomials, we make use of the \texttt{lex} function we defined at the beginning of the \texttt{list_util} file. For convenience, we also define \texttt{mono_lt}, which is a proposition that states that some monomial is less than another.

Definition mono_cmp := lex compare.
Definition mono_lt m n := mono_cmp m n = Lt.

A simple but useful definition is \texttt{vars}, which allows us to take any polynomial and get a list of all the variables in it. This is simply done by concatenating all of the monomials into one large list of variables and removing any repeated variables.

Clearly then, there will never be any duplicates in the \texttt{vars} of some polynomial.

Definition vars \( p : \text{poly} \) : list var := nodup var_eq_dec (concat \( p \)).

Hint Unfold vars.

Lemma NoDup_vars : \( \forall (p : \text{poly}), \)
\textbf{NoDup} (vars \( p \)).

Proof.
intros \( p \). unfold vars. apply NoDup_nodup.
Qed.

This next lemma allows us to convert from a statement about \texttt{vars} to a statement about the monomials themselves. If some variable \( x \) is not in the variables of a polynomial \( p \), then every monomial in \( p \) must not contain \( x \).

Lemma in_mono_in_vars : \( \forall \ x \ p, \)
\( (\forall m : \text{mono}, \ln m \ p \rightarrow \neg \ln x m) \leftrightarrow \neg \ln x (\text{vars} \ p) \).

Proof.
intros \( x \ p \). split.
- intros \( H \). induction \( p \).
  + simpl. auto.
  + unfold not in * \( \). intro. apply IHp.
    \times intros \( m \) Hin \( \). apply \( H \). intuition.
    \times unfold vars in *. apply nodup \( \ln \) in \( H0 \). apply nodup \( \ln \). simpl in \( H0 \).
    apply in \( \text{app or in} \) \( H0 \). destruct \( H0 \).
    \- exfalso. apply \( (H a) \). intuition. auto.
    \- auto.
- intros \( H \) \( m \) Hin Hin'. apply \( H \). clear \( H \). induction \( p \).
inversion Hin.
+ unfold vars in *. rewrite nodup_ln. rewrite nodup_ln in IHp. simpl.
  apply in_or_app. destruct Hin.
  × left. rewrite H. auto.
  × auto.
Qed.

6.1.3 Stronger Definitions

Because, as far as Coq is concerned, any list of natural numbers is a monomial, it is necessary to define a few more predicates about monomials and polynomials to ensure our desired properties hold. Using these in proofs will prevent any random list from being used as a monomial or polynomial.

Monomials are simply lists of natural numbers that, for ease of comparison, are sorted least to greatest. A small sublety is that we are insisting they are sorted with lt, meaning less than, rather than le, or less than or equal to. This way, the Sorted predicate will insist that each number is less than the one following it, thereby preventing any values from being equal to each other. In this way, we simultaneously enforce the sorting and lack of duplicated values in a monomial.

Definition is_monomial (m : mono) : Prop := Sorted lt m.

Polynomials are sorted lists of lists, where all of the lists in the polynomial are monomials. Similarly to the last example, we use mono lt to simultaneously enforce sorting and no duplicates.

Definition is_polynomial (p : poly) : Prop :=
  Sorted mono lt p ∧ ∀ m, In m p → is_monomial m.

Hint Unfold is_monomial is_polynomial.

Hint Resolve NoDup_cons NoDup_nil Sorted_cons.

There are a few useful things we can prove about these definitions too. First, because of the sorting, every element in a monomial is guaranteed to be less than the element after it.

Lemma mono_order : ∀ x y m,
  is_monomial (x :: y :: m) →
  x < y.

Proof.
  unfold is_monomial.
  intros x y m H.
  apply Sorted_inv in H as [].
  apply HdRel_inv in H0.
  apply H0.
Qed.

Similarly, if x :: m is a monomial, then m is also a monomial.
Lemma mono_cons : \forall x m,  
is_mono (x :: m) \rightarrow  
is_mono m.  
Proof.  
unfold is_mono.  
intros x m H. apply Sorted_inv in H as []. apply H.  
Qed.  
The same properties hold for is_poly as well; any list in a polynomial is guaranteed to be less than the lists after it, and if m :: p is a polynomial, we know both that p is a polynomial and that m is a monomial.  
Lemma poly_order : \forall m n p,  
is_poly (m :: n :: p) \rightarrow  
mono_lt m n.  
Proof.  
unfold is_poly.  
intros.  
destruct H.  
apply Sorted_inv in H as [].  
apply HdRel_inv in H1.  
apply H1.  
Qed.  
Lemma poly_cons : \forall m p,  
is_poly (m :: p) \rightarrow  
is_poly p \land is_mono m.  
Proof.  
unfold is_poly.  
intros.  
destruct H.  
apply Sorted_inv in H as[].  
split.  
- split; auto.  
  intros. apply H0, in_cons, H2.  
- apply H0, in_eq.  
Qed.  
Lastly, for completeness, nil is both a polynomial and monomial, the polynomial representation for one as we described before is a polynomial, and a singleton variable is a polynomial.  
Lemma nil_is_mono :  
is_mono [].  
Proof.  
unfold is_mono. auto.
Qed.

Lemma nil_is_poly :
  is_poly [].
Proof.
  unfold is_poly. split; auto.
  intro; contradiction.
Qed.

Lemma one_is_poly :
  is_poly [[]].
Proof.
  unfold is_poly. split; auto.
  intro. intro. simpl in H. destruct H.
  - rewrite ← H. apply nil_is_mono.
  - inversion H.
Qed.

Lemma var_is_poly : ∀ x,
  is_poly [[x]].
Proof.
  intros x. unfold is_poly. split.
  - apply Sorted_cons; auto.
  - intros m H. simpl in H; destruct H; inversion H.
    unfold is_mono. auto.
Qed.

In unification, a common concept is a ground term, or a term that contains no variables. If some polynomial is a ground term, then it must either be equal to 0 or 1.

Lemma no_vars_is_ground : ∀ p,
  is_poly p →
  vars p = [] →
  p = [] ∨ p = [[]].
Proof.
  intros p H H0. induction p; auto.
  induction a.
  - destruct IHp.
    + apply poly_cons in H. apply H.
    + unfold vars in H0. simpl in H0. apply H0.
    + rewrite H1. auto.
    + rewrite H1 in H unfold is_poly in H. destruct H. inversion H.
      inversion H6. inversion H8.
    - unfold vars in H0. simpl in H0. destruct in_dec in H0.
      + rewrite ← nodup_in in i rewrite H0 in i. inversion i.
      + inversion H0.

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Qed.

Hint Resolve mono_order mono_cons poly_order poly_cons nil_is_mono nil_is_poly
            var_is_poly one_is_poly.

6.2 Sorted Lists and Sorting

Clearly, because we want to maintain that our monomials and polynomials are sorted at all
times, we will be dealing with Coq's Sorted proposition a lot. In addition, not every list
we want to operate on will already be perfectly sorted, so it is often necessary to sort lists
ourselves. This next section serves to give us all of the tools necessary to operate on sorted
lists.

6.2.1 Sorting Lists

In order to sort our lists, we will make use of the Sorting module in the standard library,
which implements a version of merge sort.

For sorting variables in a monomial, we can simply reuse the already provided NatSort
module.

Module Import VARSORT := NATSORT.

Sorting the monomials in a polynomial is slightly more complicated, but still straightforward
thanks to the Sorting module. First, we need to define a MONOORDER, which must
be a total less-than-or-equal-to comparator.

This is accomplished by using our mono_cmp defined earlier, and simply returning true
for either less than or equal to.

We also prove a relatively simple lemma about this new MONOORDER, which states that
if \( x \leq y \) and \( y \leq x \), then \( x \) must be equal to \( y \).

Require Import Orders.

Module MONOORDER <: TOTALLeBOOL.

Definition t := mono.

Definition leb m n :=
    match mono_cmp m n with
    | Lt => true
    | Eq => true
    | Gt => false
    end.

Infix "\le m" := leb (at level 35).

Lemma leb_total :
    \forall m n, (m \le m n = true) \lor (n \le m m = true).

Proof.
    intros n m. unfold "\le m". destruct (mono_cmp n m) eqn:Hcomp; auto.

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Lemma leb_both_eq : \forall x y, 
  is_true (MonoOrder.leb x y) \rightarrow 
  is_true (MonoOrder.leb y x) \rightarrow 
  x = y.

Proof.
  intros x y H H0. unfold is_true, MonoOrder.leb in *.
  destruct (mono_cmp y x) eqn:Hyx; destruct (mono_cmp x y) eqn:Hxy;
  unfold mono_cmp in *;
  try (apply lex_rev_lt_gt in Hxy; rewrite Hxy in Hyx; inversion Hyx);
  try (apply lex_rev_lt_gt in Hyx; rewrite Hxy in Hyx; inversion Hxy);
  try inversion H; try inversion H0.
  apply lex_eq in Hxy; auto.
Qed.

After this order has been defined and its totality has been proven, we simply define a new MONOSORT module to be a sort based on this MONOORDER.

Now, we have a simple sort function for both monomials and polynomials, as well as a few useful lemmas about the sort functions’ correctness.

Module Import MONOSORT := Sort MONOORDER.

One technique that helps us deal with the difficulty of sorted lists is proving that each of our four comparators - lt, VarOrder, mono_lt, and MONOORDER - are all transitive. This allows us to seamlessly pass between the standard library’s Sorted and StronglySorted propositions, making many proofs significantly easier.

All four of these are proved relatively easily, mostly by induction and destructing the comparison of the individual values.

Lemma lt_Transitive :
  Relations_1.Transitive lt.

Proof.
  unfold Relations_1.Transitive. intros. apply lt_trans with (m:=y); auto.
Qed.

Lemma VarOrder_Transitive :
  Relations_1.Transitive (fun x y \Rightarrow is_true (NatOrder.leb x y)).

Proof.
  unfold Relations_1.Transitive, is_true.
  induction x, y, z; intros; try reflexivity; simpl in *.
  - inversion H.
  - inversion H.
  - inversion H0.
  - apply IHx with (y:=y); auto.
Qed.

Lemma mono_lt_Transitive : Relations_1_TRANSITIVE mono_lt.
Proof.

unfold Relations_1_TRANSITIVE, is_true, mono_lt, mono_cmp.
induction x, y, z; intros; try reflexivity; simpl in *.
- inversion H.
- inversion H0.
- inversion H.
- inversion H0.
- destruct \( \text{(a \neq n0)} \) \( eqn:Han0 \).
  + apply compare_eq_iff in \( Han0 \). rewrite \( Han0 \) in H.
  destruct \( (n \neq n0) \) \( eqn:Hn0 \).
  × rewrite compare_antisym in \( Hn0 \). unfold CompOpp in \( Hn0 \).
  destruct \( (n0?\neq n) \); try inversion \( Hn0 \). apply \( (IHx \_ \_ H H0) \).
  × rewrite compare_antisym in \( Hn0 \). unfold CompOpp in \( Hn0 \).
  destruct \( (n0?\neq n0) \); try inversion \( Hn0 \). inversion H.
  × inversion H0.
  + auto.
  + destruct \( (n \neq n0) \) \( eqn:Hnn0 \).
  × apply compare_eq_iff in \( Hnn0 \). rewrite \( Hnn0 \) in H. rewrite \( Han0 \) in H.
  inversion H.
  × apply compare_lt_iff in \( Hnn0 \). apply compare_gt_iff in \( Han0 \).
  apply lt_trans with \( (n= n0) \) in \( Han0 \); auto. apply compare_lt_iff in \( Han0 \).
  rewrite compare_antisym in \( Han0 \). unfold CompOpp in \( Han0 \).
  destruct \( (a?= n0) \); try inversion \( Han0 \). inversion H.
  × inversion H0.
Qed.

Lemma MonoOrder_Transitive :
Relations_1_TRANSITIVE (fun x y \Rightarrow is_true (MonoOrder.leb x y)).
Proof.

unfold Relations_1_TRANSITIVE, is_true, MonoOrder.leb, mono_cmp.
induction x, y, z; intros; try reflexivity; simpl in *.
- inversion H.
- inversion H.
- inversion H0.
- destruct \( \text{(a \neq n)} \) \( eqn:Han \).
  + apply compare_eq_iff in \( Han \). rewrite \( Han \). destruct \( (n \neq n0) \) \( eqn:Hn0 \).
  × apply \( (IHx \_ \_ H H0) \).
  × reflexivity.
  × inversion H0.
  + destruct \( (n \neq n0) \) \( eqn:Hn0 \).
× apply compare_eq_iff in Hn0. rewrite ← Hn0. rewrite Han. reflexivity.
× apply compare_lt_iff in Han. apply compare_lt_iff in Hn0.
  apply (lt_trans a n n0 Han) in Hn0. apply compare_lt_iff in Hn0.
  rewrite Hn0. reflexivity.
× inversion H0.
+ inversion H.
Qed.

6.2.2 Sorting and Permutations

The entire purpose of ensuring our monomials and polynomials remain sorted at all times is so that two polynomials containing the same elements are treated as equal. This definition obviously lends itself very well to the use of the Permutation predicate from the standard library, which explains why we proved so many lemmas about permutations during list_util.

When comparing equality of polynomials or monomials, this sort function is often extremely tricky to deal with. Induction over a list being passed to sort is nearly impossible, because the induction element a is not guaranteed to be the least value, so will not easily make it outside of the sort function. As a result, the induction hypothesis is almost always useless.

To combat this, we will prove a series of lemmas relating sort to Permutation, since clearly sorting has no effect when we are comparing the lists in an unordered fashion. The simplest of these lemmas is that if either term of a Permutation is wrapped in a sort function, we can easily get rid of it without changing the provability of these statements.

Lemma Permutation_VarSort_l : ∀ m n,
  Permutation m n ↔ Permutation (VarSort.sort m) n.
Proof.
  intros m n. split; intro.
  - apply Permutation_trans with (l':=m). apply Permutation_sym.
    apply VarSort.Permuted_sort. apply H.
  - apply Permutation_trans with (l':=(VarSort.sort m)).
    apply VarSort.Permuted_sort. apply H.
Qed.

Lemma Permutation_VarSort_r : ∀ m n,
  Permutation m n ↔ Permutation m (VarSort.sort n).
Proof.
  intros m n. split; intro.
  - apply Permutation_sym. rewrite ← Permutation_VarSort_l.
    apply Permutation_sym; auto.
  - apply Permutation_sym. rewrite → Permutation_VarSort_l.
    apply Permutation_sym; auto.
Qed.

Lemma Permutation_MonoSort_r : ∀ p q,
**Permutation** \( p q \leftrightarrow \text{Permutation} \ (\text{sort} \ q) \).

Proof.

intros \( p q \). split; intro \( H \).
- apply Permutation\_trans with \((l' := q)\). apply \( H \). apply Permut\_sort.
- apply Permutation\_trans with \((l' := \text{sort} \ q))\). apply \( H \). apply Permutation\_sym.
  apply Permut\_sort.

Qed.

Lemma Permutation\_MonoSort\_l : \( \forall \ p \ q, \)
**Permutation** \( p q \leftrightarrow \text{Permutation} \ (\text{sort} \ p) \ q \).

Proof.

intros \( p q \). split; intro \( H \).
- apply Permutation\_sym. rewrite \( \leftarrow \) Permutation\_MonoSort\_r.
  apply Permutation\_sym. auto.
- apply Permutation\_sym. rewrite Permutation\_MonoSort\_r.
  apply Permutation\_sym. auto.

Qed.

More powerful is the idea that, if we know we are dealing with sorted lists, there is no difference between proving lists are equal and proving they are Permutations. While this seems intuitive, it is actually fairly complicated to prove in Coq.

For monomials, the proof begins by performing induction on both lists. The first three cases are very straightforward, and the only challenge comes from the third case. We approach the third case by first comparing the two induction elements, \( a \) and \( a' \).

This forms three goals for us - one where \( a = a' \), one where \( a < a' \), and one where \( a > a' \). The first goal is extremely straightforward, and follows from the induction hypothesis almost immediately after using a few compare lemmas.

This leaves us with the next two goals, which seem to be more challenging at first. However, some further thought leads us to the conclusion that both goals should both be contradictions. If the lists are both sorted, and they contain all the same elements, then they should have the same element, at the head of the list, which is the least element of the set. This element is clearly \( a \) for the first list, and \( a' \) for the second. However, our destruct of compare has left us with a hypothesis stating that they are not equal! This is the source of the contradiction.

To get Coq to see our contradiction, we first make use of the Transitive lemmas we proved earlier to convert to StronglySorted. This allows us to get a hypothesis in the second goal that states that \( a' \) must be less than everything in the second list. Because \( a \) is not equal to \( a' \), this implied that \( a \) is somewhere else in the second list, and therefore \( a' \) is less than \( a \). This clearly contradicts the fact that \( a < a' \). The third goal looks the same, but in reverse.

Lemma Permutation\_Sorted\_mono\_eq : \( \forall \ (m n : \text{mono}), \)

\[ \text{Permutation} \ m n \rightarrow \]
\[ \text{Sorted} \ (\text{fun} \ n m \Rightarrow \text{is\_true} \ (\text{leb} \ n m)) \ m \rightarrow \]
\[ \text{Sorted} \ (\text{fun} \ n m \Rightarrow \text{is\_true} \ (\text{leb} \ n m)) \ n \rightarrow \]
\[ m = n. \]
Proof.

intros m n Hp Hsl Hsm. generalize dependent n.
induction m; induction n; intros.
- reflexivity.
- apply Permutation_nil in Hp. auto.
- apply Permutation_sym, Permutation_nil in Hp. auto.
- clear IHn. apply Permutation_incl in Hp as Hp'. destruct Hp'.
  destruct (a ?= a0) eqn:Hcomp.
  + apply compare_eq_iff in Hcomp. rewrite Hcomp in *.
    apply Permutation_cons_inv in Hp. f_equal; auto.
  apply IHm.
  × apply Sorted_inv in Hsl. apply Hsl.
  × apply Hp.
  × apply Sorted_inv in Hsm. apply Hsm.
  + apply compare_lt_iff in Hcomp as Hneq. apply incl_cons_inv in H.
    destruct H. apply Sorted_StronglySorted in Hsm.
    apply StronglySorted_inv in Hsm as [].
    × simpl in H. destruct H; try (rewrite H in Hneq; apply lt_irrefl in Hneq;
      contradiction). pose (Forall_In _ _ _ H H3). simpl in i.
    unfold is_true in i. apply leb_le in i. apply lt_not_le in Hneq.
    contradiction.
  × apply VarOrder_Transitive.
  + apply compare_gt_iff in Hcomp as Hneq. apply incl_cons_inv in H0.
    destruct H0.
    apply Sorted_StronglySorted in Hsl. apply StronglySorted_inv in Hsl as [].
    × simpl in H0. destruct H0; try (rewrite H0 in Hneq;
      apply gt_irrefl in Hneq; contradiction). pose (Forall_In _ _ _ H0 H3).
      simpl in i. unfold is_true in i. apply leb_le in i.
      apply lt_not_le in Hneq. contradiction.
  × apply VarOrder_Transitive.
Qed.

We also wish to prove the same thing for polynomials. This proof is identical in spirit,
as we do the same double induction, destructing of compare, and find the same two contra-
dictions. The only difference is the use of lemmas about lex instead of compare, since now
we are dealing with lists of lists.

Lemma Permutation_Sorted_eq : ∀ (l m : list mono),
  Permutation l m →
  Sorted (fun x y ⇒ is_true (MonoOrder.leb x y)) l →
  Sorted (fun x y ⇒ is_true (MonoOrder.leb x y)) m →
  l = m.
Proof.
  intros l m Hp Hsl Hsm. generalize dependent m.
induction \( l \); induction \( m \); intros.
- reflexivity.
- apply Permutation_nil in \( Hp \). auto.
- apply Permutation_sym, Permutation_nil in \( Hp \). auto.
- clear IHm. apply Permutation_incl in \( Hp \) as \( Hp' \). destruct \( Hp' \).
  destruct (mono_cmp a a0) eqn:Hcomp.
  + apply lex_eq in Hcomp. rewrite Hcomp in *.
    apply Permutation_cons_inv in \( Hp \). f_equal; auto.
    apply IHl.
    × apply Sorted_inv in \( Hsl \). apply \( Hsl \).
    × apply \( Hp \).
    × apply Sorted_inv in \( Hsm \). apply \( Hsm \).
  + apply lex_neq' in Hcomp as Hneq. apply incl_cons_inv in \( H \). destruct \( H \).
    apply Sorted_StronglySorted in \( Hsm \). apply StronglySorted_inv in \( Hsm \) as \[].
    × simpl in \( H \). destruct \( H \); try (rewrite \( H \) in Hneq; contradiction).
      pose (Forall_In _ _ _ _ H H3). simpl in \( i \). unfold is_true,
      MonoOrder.leb. mono_cmp in \( i \). apply lex_rev_lt_gt in Hcomp.
      rewrite Hcomp in \( i \). inversion \( i \).
    × apply MonoOrder_Transitive.
  + apply lex_neq' in Hcomp as Hneq. apply incl_cons_inv in \( H0 \). destruct \( H0 \).
    apply Sorted_StronglySorted in \( Hsl \). apply StronglySorted_inv in \( Hsl \) as \[].
    × simpl in \( H0 \). destruct \( H0 \); try (rewrite \( H0 \) in Hneq; contradiction).
      pose (Forall_In _ _ _ _ H0 H3). simpl in \( i \). unfold is_true in \( i \).
      unfold MonoOrder.leb in \( i \). rewrite Hcomp in \( i \). inversion \( i \).
    × apply MonoOrder_Transitive.
Qed.

Another useful form of these two lemmas is that if at any point we are attempting to prove that \( \text{sort} \) of one list equals \( \text{sort} \) of another, we can ditch the \( \text{sort} \) and instead prove that the two lists are permutations. These lemmas will come up a lot in future proofs, and has made some of our work much easier.

Lemma Permutation_sort_mono_eq : \( \forall l m, \)
\[\text{Permutation } l \leftrightarrow \text{VarSort.sort } l = \text{VarSort.sort } m.\]

Proof.
  intros \( l \) \( m \); split; intros \( H \).
  - assert \( (H0 : \text{Permutation } (\text{VarSort.sort } l) (\text{VarSort.sort } m)) \).
    + apply Permutation_trans with \( (l':=(\text{VarSort.sort } l)) (l':=m) \)
      \( (l':=\text{VarSort.sort } m) \).
      × apply Permutation_sym. apply Permutation_sym in \( H \).
      apply (Permutation_trans H (VarSort.Permuted_sort l)).
      × apply VarSort.Permuted_sort.
    + apply \( (\text{Permutation_SORTED_mono_eq } _ _ H0 \ (\text{VarSort.LocallySorted_sort } l) \ (\text{VarSort.LocallySorted_sort } m)) \).
- assert (Permutation (VarSort.sort l) (VarSort.sort m)).
  + rewrite H. apply Permutation_refl.
    apply (Permutation_trans p) in H0. apply Permutation_sym in p0.
    apply (Permutation_trans H0) in p0. apply p0.
Qed.

Lemma Permutation_sort_eq : ∀ l m,  
Permutation l m ↔ sort l = sort m.
Proof.
  intros l m. split; intros H.
  - assert (H0 : Permutation (sort l) (sort m)).
    + apply Permutation_trans with (l:=sort l) (l':=m) (l'':=sort m).
      × apply Permutation_sym. apply Permutation_sym in H.
      apply (Permutation_trans H (Permuted_sort l)).
      × apply Permuted_sort.
    + apply (Permutation_Sorted_eq _ _ H0 (LocallySorted_sort l)
      (LocallySorted_sort m)).
  - assert (Permutation (sort l) (sort m)).
    + rewrite H. apply Permutation_refl.
    + pose (Permuted_sort l). pose (Permuted_sort m).
      apply (Permutation_trans p) in H0. apply Permutation_sym in p0.
      apply (Permutation_trans H0) in p0. apply p0.
Qed.

6.3 Repairing Invalid Monomials & Polynomials

Clearly, there is a very strict set of rules we would like to be true about all of the polynomials and monomials we workd with. These rules are, however, relatively tricky to maintain when it comes to writing functions that operate over monomials and polynomials. Rather than rely on our ability to define every function to perfectly maintain this set of rules, we decided to define two functions to “repair” any invalid monomials or polynomials. These functions, given a list of variables or a list of list of variables, will apply a few functions to them such that at the end, we are left with a properly formatted monomial or polynomial.

6.3.1 Converting Between lt and le

A small problem with the sort function provided by the standard library is that it requires us to use a le comparator, as opposed to lt like we use in our is_mono and is_poly definitions. However, as we said before, because our lists have no duplicates le and lt are equivalent. Obviously, though, saying this isn’t enough - we must prove it for it to be useful to us in proofs.
The first step to proving this is proving that this is true when dealing with the **HdRel** definition that **Sorted** is built on top of. These lemmas state that, if \( a \) holds the \( le \) relation with a list, and there are also no duplicates in \( a :: l \), that \( a \) also holds the \( lt \) relation with the list. These proofs are both relatively straightforward, especially with the use of the NoDup_neq lemma proven earlier.

**Lemma** \( \text{HdRel} \ _\text{le} \ _\text{lt} \):
\[
\forall \ a \ m, \quad \text{HdRel} \ (\text{fun} \ n \ m \Rightarrow \text{is} \ _\text{true} \ (\text{leb} \ n \ m)) \ a \ m \land \text{NoDup} \ (a :: m) \rightarrow \\
\text{HdRel} \ \text{lt} \ a \ m.
\]

**Proof.**

- **intros** \( a \ m \ []).\ **remember** \( (\text{fun} \ n \ m \Rightarrow \text{is} \ _\text{true} \ (\text{leb} \ n \ m)) \) as \( le \).
- **destruct** \( m \).
  - **apply** HdRel_nil.
  - **apply** HdRel_cons. **apply** HdRel_inv in \( H \).
    - **apply** (NoDup_neq _ a n) in \( H0 \); **intuition**. **rewrite** Heqle in \( H \).
    - **unfold** is_true in \( H \). **apply** leb_le in \( H \). **destruct** \( (a \ ?= n) \) eqn:Hcomp.
      - **apply** compare_eq_iff in \( Hcomp \). **contradiction**.
      - **apply** compare_lt_iff in \( Hcomp \). **apply** Hcomp.
      - **apply** compare_gt_iff in \( Hcomp \). **apply** leb_correct_conv in \( Hcomp \).
        - **apply** leb_correct in \( H \). **rewrite** \( H \) in \( Hcomp \). **inversion** \( Hcomp \).

**Qed.**

**Lemma** \( \text{HdRel} \ _\text{mono} \ _\text{le} \ _\text{lt} \):
\[
\forall \ a \ p, \quad \text{HdRel} \ (\text{fun} \ n \ m \Rightarrow \text{is} \ _\text{true} \ (\text{MonoOrder.leb} \ n \ m)) \ a \ p \land \text{NoDup} \ (a :: p) \rightarrow \\
\text{HdRel} \ \text{mono} \ _\text{lt} \ a \ p.
\]

**Proof.**

- **intros** \( a \ p \ []).\ **remember** \( (\text{fun} \ n \ m \Rightarrow \text{is} \ _\text{true} \ (\text{MonoOrder.leb} \ n \ m)) \) as \( le \).
- **destruct** \( p \).
  - **apply** HdRel_nil.
  - **apply** HdRel_cons. **apply** HdRel_inv in \( H \).
    - **apply** (NoDup_neq _ a l) in \( H0 \); **intuition**. **rewrite** Heqle in \( H \).
    - **unfold** is_true in \( H \). **unfold** MonoOrder.leb in \( H \). **unfold** monoltk.
      - **destruct** \( (\text{mono} \ _\text{cmp} \ a \ l) \) eqn:Hcomp.
        - **apply** lex_eq in \( Hcomp \). **contradiction**.
        - **reflexivity**.
        - **inversion** \( H \).

**Qed.**

Now, to apply these lemmas - we prove that if a list is **Sorted** with a \( le \) operator and has no duplicates, that it is also **Sorted** with the corresponding \( lt \) operator.

**Lemma** \( \text{VarSort} \ _\text{Sorted} \):
\[
\forall \ m, \quad \text{Sorted} \ (\text{fun} \ n \ m \Rightarrow \text{is} \ _\text{true} \ (\text{leb} \ n \ m)) \ m \land \text{NoDup} \ m \rightarrow \\
\text{Sorted} \ \text{lt} \ m.
\]

**Proof.**
intros m []. remember (fun n m ⇒ is_true (le n m)) as le.

induction m.
- apply Sorted_nil.
- apply Sorted_inv in H. apply Sorted_cons.
  + apply IHm.
    × apply H.
    × apply NoDup_cons_iff in H0. apply H0.
  + apply HdRel_le_lt. split.
    × rewrite ← Heqle. apply H.
    × apply H0.

Qed.

Lemma MonoSort_Sorted : ∀ p, Sorted (fun n m ⇒ is_true (MonoOrder.le n m)) p ∧ NoDup p → Sorted mono_lt p.

Proof.
  intros p []. remember (fun n m ⇒ is_true (MonoOrder.le n m)) as le.
  induction p.
- apply Sorted_nil.
- apply Sorted_inv in H. apply Sorted_cons.
  + apply IHp.
    × apply H.
    × apply NoDup_cons_iff in H0. apply H0.
  + apply HdRel_mono_le_lt. split.
    × rewrite ← Heqle. apply H.
    × apply H0.

Qed.

For convenience, we also include the inverse - if a list is Sorted with an lt operator, it is also Sorted with the matching le operator.

Lemma Sorted_VarSorted : ∀ (m : mono), Sorted lt m → Sorted (fun n m ⇒ is_true (le n m)) m.

Proof.
  intros m H. induction H.
  - apply Sorted_nil.
  - apply Sorted_cons.
    + apply IHSorted.
    + destruct l.
      × apply HdRel_nil.
      × apply HdRel_cons. apply HdRel_inv in H0. apply lt_le_incl in H0.
        apply leb_le in H0. apply H0.
  
Qed.
Lemma Sorted_MonoSorted : \(\forall (p : \text{poly}),\)
\(\text{Sorted \ mono\_lt} \ p \rightarrow \text{Sorted} (\text{fun} \ n \ m \Rightarrow \text{is\_true} (\text{MonoOrder\_leb} \ n \ m)) \ p.\)

Proof.
\(\text{intros } p \ H. \ \text{induction } H.\)
- apply Sorted_nil.
- apply Sorted_cons.
  + apply IHSorted.
  + destruct \(l.\)
    \(\times\) apply HdRel_nil.
    \(\times\) apply HdRel_cons. apply HdRel_inv in \(H0.\) unfold MonoOrder.leb.
    rewrite \(H0.\) auto.

Qed.

Another obvious side effect of what we have just proven is that if a list is Sorted with an lt operator, clearly there are no duplicates, as no elements are equal to each other.

Lemma NoDup_VarSorted : \(\forall m,\)
\(\text{Sorted \ lt} \ m \rightarrow \text{NoDup} \ m.\)

Proof.
\(\text{intros } m \ H. \ \text{apply Sorted\_StronglySorted in } H.\)
- induction \(m;\) auto.
  apply StronglySorted_inv in \(H\) as \([].\) apply NoDup\_forall\_neq.
  + apply Forall\_forall. intros \(x \ H1n.\) rewrite Forall\_forall in \(H0.\)
    apply \(lt\_neq.\) apply \(H0.\) apply \(H1n.\)
    + apply IHm. apply \(H.\)
  - apply \(lt\_Transitive.\)

Qed.

Lemma NoDup_MonoSorted : \(\forall p,\)
\(\text{Sorted \ mono\_lt} \ p \rightarrow \text{NoDup} \ p.\)

Proof.
\(\text{intros } p \ H. \ \text{apply Sorted\_StronglySorted in } H.\)
- induction \(p;\) auto.
  apply StronglySorted_inv in \(H\) as \([].\) apply NoDup\_forall\_neq.
  + apply Forall\_forall. intros \(x \ H1n.\) rewrite Forall\_forall in \(H0.\)
    pose (\text{lex\_neq'} \ a \ x). destruct \(a0.\) apply \(H1\) in \(H0;\) auto.
    + apply IHp. apply \(H.\)
  - apply \(mono\_lt\_Transitive.\)

Qed.

There are a few more useful lemmas we would like to prove about our sort functions before we can define and prove the correctness of our repair functions. Mostly, we want to know that sorting a list has no effect on some properties of it.

Specifically, if an element was in a list before it was sorted, it is also in it after, and vice
versa. Similarly, if a list has no duplicates before being sorted, it also has no duplicates after.

**Lemma** \(\text{In \_ \_sorted : } \forall a \, l,\)

\[\text{In} \, a \, l \leftrightarrow \text{In} \, a \, (\text{sort} \, l).\]

**Proof.**
- intros \(a \, l\). pose \((\text{MonoSort.Permuted_sort} \, l)\). split; intros \(Hin\).
- apply \((\text{Permutation_in} \, p \, Hin)\).
- apply \((\text{Permutation_in'} \, (\text{Logic.eq_refl} \, a) \, p)\). auto.

Qed.

**Lemma** \(\text{NoDup.VarSort : } \forall (m : \text{mono}),\)

\[\text{NoDup} \, m \rightarrow \text{NoDup} \, (\text{VarSort.sort} \, m).\]

**Proof.**
- intros \(m \, Hdup\). pose \((\text{VarSort.Permuted_sort} \, m)\).
- apply \((\text{Permutation_NoDup} \, p \, Hdup)\).

Qed.

**Lemma** \(\text{NoDup.MonoSort : } \forall (p : \text{poly}),\)

\[\text{NoDup} \, p \rightarrow \text{NoDup} \, (\text{MonoSort.sort} \, p).\]

**Proof.**
- intros \(p \, Hdup\). pose \((\text{MonoSort.Permuted_sort} \, p)\).
- apply \((\text{Permutation_NoDup} \, p \, Hdup)\).

Qed.

### 6.3.2 Defining the Repair Functions

Now time for our definitions. To convert a list of variables into a monomial, we first apply \(\text{nodup}\), which removes all duplicates. We use \(\text{nodup}\) rather than \(\text{nodup_cancel}\) because \(x+x \approx_B x\), so we want one copy to remain. After applying \(\text{nodup}\), we use our \text{VARSORT} module to sort the list from least to greatest.

**Definition** \(\text{make_mono} \,(l: \text{list} \, \text{nat}) : \text{mono} :=\)

\[\text{VarSort.sort} \, (\text{nodup} \, \text{var_eq_dec} \, l).\]

The process of converting a list of list of variables into a polynomial is very similar. First we map across the list applying \(\text{make_mono}\), so that each sublist is properly formatted. Then we apply \(\text{nodup_cancel}\) to remove duplicates. In this case, we use \(\text{nodup_cancel}\) instead of \(\text{nodup}\) because \(x+x = 0\), so we want pairs to cancel out. Lastly, we use our \text{MONOSORT} module to sort the list.

**Definition** \(\text{make_poly} \,(l: \text{list} \, \text{mono}) : \text{poly} :=\)

\[\text{MonoSort.sort} \, (\text{nodup_cancel} \, \text{mono_eq_dec} \, (\text{map} \, \text{make_mono} \, l)).\]

**Lemma** \(\text{make_poly_refold : } \forall p,\)

\[\text{sort} \, (\text{nodup_cancel} \, \text{mono_eq_dec} \, (\text{map} \, \text{make_mono} \, p)) = \text{make_poly} \, p.\]

**Proof.** auto. Qed.
Now to prove the correctness of these lists - if you apply \texttt{make\_mono} to something, it is then guaranteed to satisfy the \texttt{is\_mono} proposition. This proof is relatively straightforward, as we have already done most of the work with \texttt{VarSort\_Sorted}; all that is left to do is show that \texttt{make\_mono} \( m \) is \texttt{Sorted} and has no duplicates, which is obvious considering that is exactly what \texttt{make\_mono} does!

**Lemma make\_mono\_is\_mono:** \( \forall m, \)
\[ \texttt{is\_mono} (\texttt{make\_mono} \ m). \]

**Proof.**
- \( \text{intros } \ m. \ \text{unfold } \texttt{is\_mono}, \texttt{make\_mono}. \ \text{apply } \texttt{VarSort\_Sorted}. \ \text{split}. \)
  - \( \text{apply } \texttt{VarSort.LocallySorted} \ m. \)
  - \( \text{apply } \texttt{NoDup\_VarSort} \ m. \ \text{apply } \texttt{NoDup\_nodup}. \)

Qed.

The proof for \texttt{make\_poly\_is\_poly} is almost identical, with the addition of one part. The \texttt{is\_poly} predicate still asks us to prove that the list is \texttt{Sorted}, which follows from \texttt{MonoSort\_Sorted} like above. The only difference is that \texttt{is\_poly} also asks us to show that each element in the list is \texttt{mono}, which follows from the use of a few \texttt{ln} lemmas and the \texttt{make\_mono\_is\_mono} we just proved thanks to the \texttt{map} in \texttt{make\_poly}.

**Lemma make\_poly\_is\_poly:** \( \forall p, \)
\[ \texttt{is\_poly} (\texttt{make\_poly} \ p). \]

**Proof.**
- \( \text{intros } \ p. \ \text{unfold } \texttt{is\_poly}, \texttt{make\_poly}. \ \text{split}. \)
  - \( \text{apply } \texttt{MonoSort\_Sorted}. \ \text{split}. \)
    - \( \text{apply } \texttt{MonoSort.LocallySorted} \ m. \)
      - \( \text{apply } \texttt{NoDup\_MonoSort} \ m. \ \text{apply } \texttt{NoDup\_nodup\_cancel}. \)
  - \( \text{intros } \ m \ Hm. \ \text{apply } \texttt{ln\_sorted} \ m. \ \text{apply } \texttt{nodup\_cancel}\_\text{in} \ m. \)
    - \( \text{apply } \texttt{in\_map}\_\text{iff} \ m. \ \text{destruct } \ Hm. \ \text{destruct } \ H. \ \text{rewrite} \leftarrow \ H. \)
    - \( \text{apply } \texttt{make\_mono\_is\_mono}. \)

Qed.

**Hint** \texttt{Resolve make\_poly\_is\_poly make\_mono\_is\_mono}.  

### 6.3.3 Facts about \texttt{make\_mono}

Before we dive into more complicated proofs involving these repair functions, there are a few simple lemmas we can prove about them.

First is that if some variable \( x \) was in a list before \texttt{make\_mono} was applied, it must also be in it after, and vice-versa.

**Lemma make\_mono\_ln:** \( \forall x \ m, \)
\[ \texttt{ln} \ x (\texttt{make\_mono} \ m) \leftrightarrow \texttt{ln} \ x \ m. \]

**Proof.**
- \( \text{intros } \ x \ m. \ \text{split}; \ \text{intro } \ H. \)
  - \( \text{unfold } \texttt{make\_mono} \ m. \ \text{pose} (\texttt{VarSort}\_\texttt{Permuted}\_\texttt{sort} (\texttt{nodup} \ \texttt{var}\_\texttt{eq}\_\texttt{dec} \ m)). \)
apply Permutation_sym in p. apply (Permutation_in _ p) in H.
apply nodup_Ln in H. auto.
- unfold make_mono. pose (VarSort.Permutated_sort (nodup var_eq_dec m)).
  apply Permutation_in with (l:=(nodup var_eq_dec m)); auto. apply nodup_Ln.
  auto.
Qed.

In addition, if some list \( m \) is already a monomial, removing anything from it will not
change that.

Lemma remove_is_mono : \( \forall \; x \; m, \)
  is_mono m \( \rightarrow \)
  is_mono (remove var_eq_dec x m).
Proof.
  intros x m H. unfold is_mono in *. apply StronglySorted_Sorted.
  apply StronglySorted_remove. apply Sorted_StronglySorted in H. auto.
  apply lt_Transitive.
Qed.

If we know that some \((l1 ++ x :: l2)\) is a mono, then clearly it is still a monomial if we
remove the x from the middle, as this will not affect the sorting at all.

Lemma mono_middle : \( \forall \; x \; l1 \; l2, \)
  is_mono \((l1 ++ x :: l2)\) \( \rightarrow \)
  is_mono \((l1 ++ l2)\).
Proof.
  intros x l1 l2 H. unfold is_mono in *. apply Sorted_StronglySorted in H.
  apply StronglySorted_Sorted. induction l1.
  - rewrite app_nil_l in *. apply StronglySorted_inv in H as []; auto.
  - simpl in *. apply StronglySorted_inv in H as |]. apply SSorted_cons; auto.
    apply Forall_forall. rewrite Forall_forall in H0. intros x0 Hin.
    apply H0. apply in_app_iff in Hin as |]; intuition.
  - apply lt_Transitive.
Qed.

Due to the nature of sorting, make_mono is commutative across list concatenation.

Lemma make_mono_app_comm : \( \forall \; m \; n, \)
  make_mono \((m ++ n)\) = make_mono \((n ++ m)\).
Proof.
  intros m n. apply Permutation_sort_mono_eq. apply Permutation_nodup.
  apply Permutation_app_comm.
Qed.

Finally, if a list \( m \) is a member of the list resulting from map make_mono, then clearly it
is a monomial.

Lemma mono_in_map_make_mono : \( \forall \; p \; m, \)
In \( m \) (map make\_mono \( p \)) \( \rightarrow \) is\_mono \( m \).

Proof.
intros. apply in\_map\_iff in \( H \) as \([x \mathbb{I}]\). rewrite \( \leftarrow H \). auto.
Qed.

**6.3.4 Facts about make\_poly**

If two lists are permutations of each other, then they will be equivalent after applying make\_poly to both.

Lemma make\_poly\_Permutation : \( \forall p q \),
\[\text{Permutation } p q \rightarrow \text{make\_poly } p = \text{make\_poly } q.\]
Proof.
intros. unfold make\_poly.
apply Permutation\_sort\_eq, nodup\_cancel\_Permutation, Permutation\_map.
auto.
Qed.

Because we have shown that sort and Permutation are equivalent, we can easily show that make\_poly is commutative accross list concatenation.

Lemma make\_poly\_app\_comm : \( \forall p q \),
\[\text{make\_poly } (p ++ q) = \text{make\_poly } (q ++ p).\]
Proof.
intros \( p q \). apply Permutation\_sort\_eq.
apply nodup\_cancel\_Permutation. apply Permutation\_map.
apply Permutation\_app\_comm.
Qed.

During make\_poly, we both sort and call nodup\_cancel. A lemma that is useful in some cases shows that it doesn’t matter what order we do these in, as nodup\_cancel will maintain the order of a list.

Lemma sort\_nodup\_cancel\_assoc : \( \forall l \),
\[\text{sort } (\text{nodup\_cancel mono\_eq\_dec } l) = \text{nodup\_cancel mono\_eq\_dec } (\text{sort } l).\]
Proof.
intros \( l \). apply Permutation\_Sorted\_eq.
- pose (Permuted\_sort (nodup\_cancel mono\_eq\_dec \( l \))).
  apply Permutation\_sym in \( p \). apply (Permutation\_trans \( p \)). clear \( p \).
  apply NoDup\_Permutation.
+ apply NoDup\_nodup\_cancel.
+ apply NoDup\_nodup\_cancel.
+ intros \( x \). split.
  \( \times \) intros \( H \). apply Permutation\_in with (\( l := \text{nodup\_cancel mono\_eq\_dec } l \)).
    apply nodup\_cancel\_Permutation. apply Permuted\_sort. auto.
  \( \times \) intros \( H \).
apply Permutation_in with (l:=nodup_cancel mono_eq_dec (sort l)).
apply nodup_cancel_Permutation. apply Permutation_sym.
apply Permuted_sort. auto.
- apply LocallySorted_sort.
- apply Sorted_nodup_cancel.
  + apply MonoOrder_Transitive.
  + apply LocallySorted_sort.
Qed.

Another obvious but useful lemma is that if a monomial \( m \) is in a list resulting from applying make_poly, is is clearly a monomial.

Lemma mono_in_make_poly : \( \forall p \ m, \)
\( \mathsf{In} \ m \ (\text{make}_\text{poly} \ p) \rightarrow \mathsf{is\_mono} \ m. \)
Proof.
intros. unfold make_poly in H. apply \( \mathsf{In\_sorted} \) in H.
apply nodup_cancel_in in H. apply (mono_in_map_make_mono _ _ H).
Qed.

### 6.4 Proving Functions “Pointless”

In the list_util file, we have two lemmas revolving around the idea that, in some cases, calling \texttt{nodup\_cancel} is “pointless”. The idea here is that, when comparing very complicated terms, it is sometimes beneficial to either add or remove an extra function call that has no effect on the final term. Until this point, we have only proven this about \texttt{nodup\_cancel} and \texttt{remove}, but there are many other cases where this is true, which will make our more complex proofs much easier. This section serves to prove this true of most of our functions.

#### 6.4.1 Working with sort Functions

The next two lemmas very simply prove that, if a list is already Sorted, then calling either \texttt{VarSort} or \texttt{MonoSort} on it will have no effect. This is relatively obvious, and is extremely easy to prove with our \texttt{Permutation / Sorted} lemmas from earlier.

Lemma no_sort_VarSorted : \( \forall m, \)
\( \mathsf{Sorted} \ l \ m \rightarrow \)
\( \mathsf{VarSort}_\text{sort} \ m = m. \)
Proof.
intros m H. apply Permutation_Sorted_mono_eq.
- apply Permutation_sym. apply VarSort.Permuted_sort.
- apply VarSort.LocallySorted_sort.
- apply Sorted_VarSorted. auto.
Qed.

Lemma no_sort_MonoSorted : \( \forall p, \)
\textbf{Sorted} \ mono\_lt \ p \ \rightarrow \ 
\text{MonoSort.sort} \ p = p.

\textbf{Proof.}
\begin{itemize}
    \item intros \ p \ \text{H. unfold make.poly. apply Permutation\_Sorted.eq.}
    \item - apply Permutation\_sym. apply Permutted\_sort.
    \item - apply LocallySorted\_sort.
    \item - apply Sorted\_MonoSorted. auto.
\end{itemize}
Qed.

The following lemma more closely aligns with the format of the \texttt{nodup\_cancel\_pointless}
lemma from \texttt{listutil}. It states that if the result of appending two lists is already going to be
sorted, there is no need to sort the intermediate lists.

This also applies if the sort is wrapped around the right argument, thanks to the \texttt{Permu-
tation} lemmas we proved earlier.

\textbf{Lemma sort\_pointless} : \ \forall \ p \ q,
\begin{itemize}
    \item sort (sort \ p ++ \ q) =
    \item sort (p ++ q).
\end{itemize}
\textbf{Proof.}
\begin{itemize}
    \item intros \ p \ q. apply Permutation\_sort.eq.
    \item apply Permutation\_app\_tail. apply Permutation\_sym.
    \item apply Permutted\_sort.
\end{itemize}
Qed.

\subsection{6.4.2 Working with make\_mono}

There are a couple forms that the proof of \texttt{make\_mono} being pointless can take. Firstly,
because we already know that \texttt{make\_mono} simply applies functions to get the list into a
form that satisfies \texttt{is\_mono}, it makes sense to prove that if some list is already a mono
that \texttt{make\_mono} will have no effect. This is proved with the help of \texttt{no\_sort\_VarSorted} and
\texttt{no\_nodup\_NoDup}.

\textbf{Lemma no\_make\_mono} : \ \forall \ m,
\begin{itemize}
    \item is\_mono \ m \ \rightarrow
    \item make\_mono \ m = m.
\end{itemize}
\textbf{Proof.}
\begin{itemize}
    \item unfold make\_mono, is\_mono. intros \ m \ \text{H. rewrite no\_sort\_VarSorted.}
    \item - apply no\_nodup\_NoDup. apply NoDup\_VarSorted in \ H. auto.
    \item - apply Sorted\_nodup; auto. apply lt\_Transitive.
\end{itemize}
Qed.

We can also prove the more standard form of \texttt{make\_mono\_pointless}, which states that if
there are nested calls to \texttt{make\_mono}, we can remove all except the outermost layer.

\textbf{Lemma make\_mono\_pointless} : \ \forall \ m \ a,
\begin{itemize}
    \item make\_mono \ (m ++ \texttt{make\_mono} \ a) = \texttt{make\_mono} \ (m ++ a).
\end{itemize}
Proof.
  intros m a. apply Permutation_sort_mono_eq. rewrite ← (nodup_pointless _ a).
  apply Permutation_nodup. apply Permutation_app_head. unfold make_mono.
  rewrite ← Permutation_VarSort_l. auto.
Qed.

Similarly, if we already know that all of the elements in a list are monomials, then
mapping make Mono across the list will have no effect on the entire list.

Lemma no_map_makeMono : ∀ p, 
  (∀ m, In m p → is_mono m) → 
  map makeMono p = p.
Proof.
  intros p H. induction p; auto.
  simpl. rewrite no_makeMono.
  - f_equal. apply IHp. intros m Hin. apply H. intuition.
  - apply H. intuition.
Qed.

Lastly, the pointless proof that more closely aligns with what we have done so far - if
makePoly is already being applied to a list, there is no need to have a call to map makeMono
on the inside.

Lemma map_makeMono_pointless : ∀ p q, 
  makePoly (map makeMono p ++ q) = 
  makePoly (p ++ q).
Proof.
  intros p q. destruct p; auto.
  simpl. unfold makePoly, simpl map.
  rewrite (no_makeMono (makeMono l)); auto. rewrite map_app. rewrite map_app.
  rewrite (no_map_makeMono (map _ _)). auto. intros m Hin.
  apply in_map_iff in Hin. destruct Hin as [x[]]. rewrite ← H. auto.
Qed.

6.4.3 Working with makePoly

Finally, we work to prove some lemmas about makePoly as a whole being pointless. These
proofs are built upon the previous few lemmas, which prove that we can remove the compo-
nents of makePoly one by one.

First up, we have a lemma that shows that if p already has no duplicates and everything
in the list is a mono, then nodup_cancel and map makeMono will both have no effect. This
lemma turns out to be very useful after something like Permutation_sort_eq has been applied,
as it can strip away the other two functions of makePoly.

Lemma unsorted_poly : ∀ p,
  NoDup p →
∀ m, In m p → is_mono m →

nodup_cancel mono_eq_dec (map make_mono p) = p.

Proof.
intros p Hdup Hin. rewrite no_map_make_mono; auto.
apply no_nodup_cancel_NoDup; auto.
Qed.

Similarly to no_make_mono, it is very straightforward to prove that if some list p is already a polynomial, then make_poly has no effect.

Lemma no_make_poly : ∀ p,
is_poly p →
make_poly p = p.

Proof.
unfold make_poly, is_poly. intros m ||. rewrite no_sort_MonoSorted.
- rewrite no_nodup_cancel_NoDup.
  + apply no_map_make_mono. intros m0 Hin. apply H0. auto.
  + apply NoDup_MonoSorted in H. rewrite no_map_make_mono; auto.
- apply Sorted_nodup_cancel.
  + apply mono_lt_Transitive.
  + rewrite no_map_make_mono; auto.

Qed.

Now onto the most important lemma. In many of the later proofs, there will be times where there are calls to make_poly nested inside of each other, or long lists of arguments appended together inside of a make_poly. In either case, the ability to add and remove extra calls to make_poly as we please proves to be very powerful.

To prove make_poly_pointless, we begin by proving a weaker version that insists that all of the arguments of p and q are all monomials. This addition makes the proof significantly easier. As one might expect, the proof is completed by using Permutation_sort_eq to remove the sort calls, nodup_cancel_pointless to remove the nodup_cancel calls, and no_map_make_mono to get rid of the map make_mono calls. After this is done, the two sides are identical.

Lemma make_poly_pointless_weak : ∀ p q,
(∀ m, In m p → is_mono m) →
(∀ m, In m q → is_mono m) →
make_poly (make_poly p ++ q) =
make_poly (p ++ q).

Proof.
intros p q Hmp Hmq. unfold make_poly.
repeat rewrite no_map_make_mono; intuition.
apply Permutation_sort_eq. rewrite sort_nodup_cancel_assoc.
rewrite nodup_cancel_pointless. apply nodup_cancel_Permutation.
apply Permutation_sym. apply Permutation_app_tail. apply Permutee_sort.
- simpl in H. rewrite in_app_iff in H. destruct H; intuition.
Now, to make the stronger and easier to use version, we simply rewrite in the opposite direction with `map_make_mono_pointless` to add extra calls of `map make_mono` in! Ironically, this proof of `make_poly_pointless` is a great example of why these “pointless” lemmas are so useful. While we can clearly tell that adding the extra call to `map make_mono` makes no difference, it makes proving things in a way that Coq understands dramatically easier at times.

After rewriting with `map_make_mono_pointless`, clearly both arguments contain all monomials, and we can use `make_poly_pointless_weak` to prove the stronger version.

Lemma `make_poly_pointless`:
\[ \forall p \ q, \quad \text{make\_poly} (\text{make\_poly} p + + q) = \text{make\_poly} (p + + q). \]

Proof.
1. `intros p q. rewrite make_poly_app_comm.`
2. `rewrite \leftarrow \text{map\_make\_mono\_pointless}. rewrite make_poly_app_comm.`
3. `rewrite \leftarrow (\text{map\_make\_mono\_pointless} p). rewrite (make_poly_app_comm \_ q).`  
4. `rewrite \leftarrow (\text{map\_make\_mono\_pointless} q).`  
5. `rewrite (make_poly_app_comm \_ (\text{map make\_mono} p)).`  
6. `rewrite (\text{make\_poly\_pointless\_weak} (\text{map make\_mono} p)). unfold make\_poly.`  
7. `rewrite (\text{no\_map\_make\_mono} (\text{map make\_mono} p)). auto.`  
8. `apply mono\_in\_map\_make\_mono. apply mono\_in\_map\_make\_mono.`  
9. `apply mono\_in\_map\_make\_mono.`

Qed.

For convenience, we also prove that it applies on the right side by using `make_poly_app_comm` twice.

Lemma `make_poly_pointless_r`:
\[ \forall p \ q, \quad \text{make\_poly} (p + + \text{make\_poly} q) = \text{make\_poly} (p + + q). \]

Proof.
1. `intros p q. rewrite make_poly_app_comm. rewrite make_poly_pointless.`
2. `apply make_poly_app_comm.`

Qed.

### 6.5 Polynomial Arithmetic

Now, the foundation for operations on polynomials has been put in place, and we can begin to get into the real meat - our arithmetic operators. First up is addition. Because we have so cleverly defined our `make_poly` function, addition over our data structures is as simple as appending the two polynomials and repairing the result back into a proper polynomial.
We also include a simple refold lemma for convenience, and a quick proof that the result of addPP is always a polynomial.

Definition addPP (p q : poly) : poly :=
    make_poly (p ++ q).

Lemma addPP_refold : ∀ p q,
    make_poly (p ++ q) = addPP p q.
Proof.
    auto.
Qed.

Lemma addPP_is_poly : ∀ p q,
    is_poly (addPP p q).
Proof.
    intros p q. apply make_poly_is_poly.
Qed.

Similarly, the definition for multiplication becomes much easier with the creation of make_poly. All we need to do is use our distribute function defined earlier to form all combinations of one monomial from each list, and call make_poly on the result.

Definition mulPP (p q : poly) : poly :=
    make_poly (distribute p q).

Lemma mulPP_is_poly : ∀ p q,
    is_poly (mulPP p q).
Proof.
    intros p q. apply make_poly_is_poly.
Qed.

Hint Resolve addPP_is_poly mulPP_is_poly.

While this definition is elegant, sometimes it is hard to work with. This has led us to also create a few more definitions of multiplication. Each is just slightly different from the last, which allows us to choose the level of completeness we need for any given multiplication proof while knowing that at the end of the day, they are all equivalent.

Each of these new definitions breaks down multiplication into two steps - multiplying a monomial times a polynomial, and multiplying a polynomial times a polynomial. Multiplying a monomial times a polynomial is simply appending the monomial to each monomial in the polynomial, and multiplying two polynomials is just multiplying each monomial in one polynomial times the other polynomial.

The difference in each of the following definitions comes from the intermediate step. Because we know that mulPP will call make_poly, there is no need to call make_poly on the result of mulMP, as shown in the first definition. However, some proofs are made easier if the result of mulMP is wrapped in map make_mono, and some are made easier if the result is wrapped in a full make_poly. As a result, we have created each of these definitions, and choose between them to help make our proofs easier.
We also include a refolding method for each, for convenience, and a proof that each new version is equivalent to the last.

**Definition mulMP** \( (p : \text{poly}) \ (m : \text{mono}) : \text{poly} := \)
\[
\text{map} \ (\text{app} \ m) \ p.
\]

**Definition mulPP'** \( (p q : \text{poly}) : \text{poly} := \)
\[
\text{make}_\text{poly} \ (\text{concat} \ (\text{map} \ (\text{mulMP} \ p) \ q)).
\]

**Lemma mulPP'_refold** : \( \forall \ p q, \)
\[
\text{make}_\text{poly} \ (\text{concat} \ (\text{map} \ (\text{mulMP} \ p) \ q)) = \text{mulPP'} \ p q.
\]
**Proof.** auto. Qed.

**Lemma mulPP_mulPP'** : \( \forall \ (p q : \text{poly}), \)
\[
\text{mulPP} \ p q = \text{mulPP'} \ p q.
\]
**Proof.**
\[
\text{intros} \ p q. \ \text{unfold} \ \text{mulPP}, \ \text{mulPP'}. \ \text{induction} \ q; \ \text{auto}.
\]
Qed.

Next, the version including a map make mono:

**Definition mulMP'** \( (p : \text{poly}) \ (m : \text{mono}) : \text{poly} := \)
\[
\text{map make}_\text{mono} \ (\text{map} \ (\text{app} \ m) \ p).
\]

**Definition mulPP''** \( (p q : \text{poly}) : \text{poly} := \)
\[
\text{make}_\text{poly} \ (\text{concat} \ (\text{map} \ (\text{mulMP'} \ p) \ q)).
\]

**Lemma mulPP''_refold** : \( \forall \ p q, \)
\[
\text{make}_\text{poly} \ (\text{concat} \ (\text{map} \ (\text{mulMP'} \ p) \ q)) = \text{mulPP''} \ p q.
\]
**Proof.** auto. Qed.

**Lemma mulPP'_mulPP''** : \( \forall \ p q, \)
\[
\text{mulPP'} \ p q = \text{mulPP''} \ p q.
\]
**Proof.**
\[
\text{intros} \ p q. \ \text{unfold} \ \text{mulPP'}, \ \text{mulPP''}, \ \text{mulMP}, \ \text{mulMP'}, \ \text{make}_\text{poly}.
\]
\[
\text{rewrite} \ \text{concat}_\text{map}_\text{map}.
\]
\[
\text{rewrite} \ \text{(no_map}_\text{make}_\text{mono} \ (\text{map} \ _. )); \ \text{auto}.
\]
\[
\text{intros. apply in_map_iff in} \ H \ \text{as} \ [n \ ||].
\]
\[
\text{rewrite} \ \text{←} \ H.
\]
\[
\text{auto}.
\]
Qed.

And finally, the version including a full make poly:

**Definition mulMP''** \( (p : \text{poly}) \ (m : \text{mono}) : \text{poly} := \)
\[
\text{make}_\text{poly} \ (\text{map} \ (\text{app} \ m) \ p).
\]

**Definition mulPP'''** \( (p q : \text{poly}) : \text{poly} := \)
\[
\text{make}_\text{poly} \ (\text{concat} \ (\text{map} \ (\text{mulMP''} \ p) \ q)).
\]
Lemma mulPP''_refold : ∀ p q,
    make_poly (concat (map (mulMP'' p) q)) =
    mulPP'' p q.
Proof. auto. Qed.

In order to make the proof of going from mulPP'' to mulPP''' easier, we begin by proving
that we can go from their corresponding mulMPs if they are wrapped in a make_poly.

Lemma mulMP'_mulMP'' : ∀ m p q,
    make_poly (mulMP' p m ++ q) = make_poly (mulMP'' p m ++ q).
Proof.
    intros m p q. unfold mulMP', mulMP''. rewrite make_poly_app_comm.
    rewrite ← map_make_mono_pointless. rewrite make_poly_app_comm.
    rewrite ← make_poly_pointless. unfold make_poly at 2.
    rewrite (no_map_make_pointless. unfold make_poly at 3.
    rewrite (make_poly_app_comm _ q). rewrite ← (map_make_mono_pointless q).
    rewrite make_poly_app_comm. auto. apply mono_in_map_make_mono.
Qed.

Lemma mulPP''_mulPP''' : ∀ p q,
    mulPP'' p q = mulPP''' p q.
Proof.
    intros p q. induction q. auto. unfold mulPP'', mulPP'''. simpl.
    rewrite mulMP'_mulMP''.
    repeat rewrite ← (make_poly_pointless_r _ (concat _)).
    f_equal. f_equal. apply IHq.
Qed.

Again, for convenience, we add lemmas to skip from mulPP to any of the other varieties.

Lemma mulPP_mulPP'' : ∀ p q,
    mulPP p q = mulPP'' p q.
Proof.
    intros. rewrite mulPP_mulPP', mulPP'_mulPP''. auto.
Qed.

Lemma mulPP_mulPP''' : ∀ p q,
    mulPP p q = mulPP''' p q.
Proof.
    intros. rewrite mulPP_mulPP'', mulPP''_mulPP'''. auto.
Qed.

Hint Unfold addPP mulPP mulPP' mulPP'' mulPP''' mulMP mulMP' mulMP''.

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6.6 Proving the 10 $B$-unification Axioms

Now that we have defined our operations so carefully, we want to prove that the 10 standard $B$-unification axioms all apply. This is extremely important, as they will both be needed in the higher-level proofs of our unification algorithm, and they show that our list-of-list setup is actually correct and equivalent to any other representation of a term.

6.6.1 Axiom 1: Additive Inverse

We begin with the inverse and identity for each addition and multiplication. First is the additive inverse, which states that for all terms $x$, $(x + x) \downarrow_p 0$.

Thanks to the definition of `nodup_cancel` and the previously proven `nodup_cancel_self`, this proof is extremely simple.

```
Lemma addPP_p_p : \forall p, 
    addPP p p = [].
Proof. 
    intros p. unfold addPP. unfold make_poly. rewrite map_app. 
    rewrite nodup_cancel_self. auto. 
Qed.
```

6.6.2 Axiom 2: Additive Identity

Next, we prove the additive identity: for all terms $x$, $(0 + x) \downarrow_p = x \downarrow_p$. This also applies in the right direction, and is extremely easy to prove since we already know that appending `nil` to a list results in that list.

Something to note is that, unlike some of the other of the ten axioms, this one is only true if $p$ is already a polynomial. Clearly, if it wasn’t, `addPP` would not return the same $p$, but rather `make_poly p`, since `addPP` will only return proper polynomials.

```
Lemma addPP_0 : \forall p, 
    is_poly p \rightarrow 
    addPP [] p = p.
Proof. 
    intros p Hpoly. unfold addPP. simpl. apply no_make_poly. auto. 
Qed.

Lemma addPP_0r : \forall p, 
    is_poly p \rightarrow 
    addPP p [] = p.
Proof. 
    intros p Hpoly. unfold addPP. rewrite app_nil_r. apply no_make_poly. auto. 
Qed.
```
6.6.3 Axiom 3: Multiplicative Identity - 1

Now onto multiplication. In $B$-unification, there are two multiplicative identities. We begin with the easier to prove of the two, which is 1. In other words, for any term $x$, $(x \cdot 1) \downarrow_P = x \downarrow_P$.

This proof is also very simply proved because of how appending nil works.

Lemma $\text{mulPP}_1r : \forall p, \is_poly p \rightarrow \text{mulPP} p \[bracket = p.$
Proof.
intros $p H$. unfold mulPP. distribute. simpl. rewrite app-nil-r.
rewrite map_id. apply no_make_poly. auto.
Qed.

6.6.4 Axiom 4: Multiplicative Inverse

Next is the multiplicative inverse, which states that for any term $x$, $(0 \cdot x) \downarrow_P = 0$.

This is proven immediately by the distribute_nil lemmas we proved in list UTIL.

Lemma $\text{mulPP}_0 : \forall p, \text{mulPP} \[\bbracket = \bbracket.$
Proof.
intros $p q$. unfold mulPP. rewrite (@distribute_nil var). auto.
Qed.

Lemma $\text{mulPP}_0r : \forall p, \text{mulPP} p \[\bbracket = \bbracket.$
Proof.
intros $p q$. unfold mulPP. rewrite (@distribute_nil_r var). auto.
Qed.

6.6.5 Axiom 5: Commutativity of Addition

The next of the ten axioms states that, for all terms $x$ and $y$, $(x + y) \downarrow_P = (y + x) \downarrow_P$.

This axiom is also rather easy, and follows entirely from the make_poly_app_comm lemma we proved earlier due to our clever addition definition.

Lemma $\text{addPP_comm} : \forall p q, \text{addPP} p q = \text{addPP} q p.$
Proof.
intros $p q$. unfold addPP. apply make_poly_app_comm.
Qed.

6.6.6 Axiom 6: Associativity of Addition

The next axiom states that, for all terms $x$, $y$, and $z$, $(x + (y + z)) \downarrow_P = ((x + y) + z) \downarrow_P$. 

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Thanks to \texttt{addPP\_comm} and all of the “pointless” lemmas we proved earlier, this proof is much easier than it might have been otherwise. These lemmas allow us to easily manipulate the operations until we end by proving that \texttt{p ++ q ++ r} is a permutation of \texttt{q ++ r ++ p}.

Lemma \texttt{addPP\_assoc} : \( \forall p q r, \)
\[
\text{addPP}\ (\text{addPP} \ p \ q) \ r = \text{addPP} \ p \ (\text{addPP} \ q \ r).
\]
Proof.
\begin{align*}
\text{intros } p q r. &\ \text{rewrite } (\text{addPP\_comm} \_ \ (\text{addPP} \_ \_)). \ \text{unfold } \text{addPP}. \\
\text{repeat rewrite } \text{make\_poly\_pointless}. &\ \text{repeat rewrite } \text{← } \text{app\_assoc}. \\
\text{apply } \text{Permutation\_sort\_eq}. &\ \text{apply nodup\_cancel\_Permutation}. \\
\text{apply } \text{Permutation\_map}. &\ \text{rewrite } (\text{app\_assoc} \ q). \\
\text{apply } \text{Permutation\_app\_comm} \text{ with } (l' := q ++ r).
\end{align*}
Qed.

\section{Axiom 7: Commutativity of Multiplication}

Now onto the harder half of the axioms. This next one states that for all terms \( x \) and \( y \), \( (x * y) \downarrow_\mathcal{P} = (y * x) \downarrow_\mathcal{P} \). In order to prove this, we have opted to use the second version of \texttt{mulPP}, which wraps the monomial multiplication in a \texttt{map make\_mono}.

The proof begins with double induction, and the first three cases are rather simple. The fourth case is slightly more complicated, but the \texttt{make\_poly\_pointless} lemma we proved earlier plays a huge role in making it simpler. We begin by simplifying, so that the \( m \) created by induction on \( q \) is distributed across the list on the left side, and the \( a \) created by induction on \( p \) is distributed across the list on the right side. Then, we use \texttt{make\_poly\_pointless} to surround the rightmost term - which now has \( a \) but not \( m \) on the left and \( m \) but not \( a \) on the right - with \texttt{make\_poly}. This additional \texttt{make\_poly} allows us to refold the mess of \texttt{maps} and \texttt{concat}s into \texttt{mulPP}, like they used to be. From there, we use the two induction hypotheses to apply commutativity, remove the redundant \texttt{make\_polys} we added, and simplify again.

In this way, we are able to cause both \( a \) and \( m \) to be distributed across the whole list on both the left and right sides of the equation. At this point, it simply requires some rearranging of \texttt{app} with the help of \texttt{Permutation}, and our left and right sides are equal.

Without the help of \texttt{make\_poly\_pointless}, we would not have been able to use the induction hypotheses until much later in the proof, and the proof would have been dramatically longer. This also makes it more readable as you step through the proof, as we can seamlessly move between the original form including \texttt{mulPP} and the more functional form consisting of \texttt{map} and \texttt{concat}.

Lemma \texttt{mulPP\_comm} : \( \forall p q, \)
\[
\text{mulPP} \ p \ q = \text{mulPP} \ q \ p.
\]
Proof.
\begin{align*}
\text{intros } p q. &\ \text{repeat rewrite mulPP\_mulPP"}. \\
\text{generalize dependent } q. &\ \text{induction } p; \ \text{induction } as \ ||m|. \\
\text{auto.}
\end{align*}
- unfold mulPP", mulMP'. simpl. rewrite (@concat_map Nil mono). auto.
- unfold mulPP", mulMP'. simpl. rewrite (@concat_map Nil mono). auto.
- unfold mulPP". simpl. rewrite (app_comm_cons _ _ (make_mono (a++m))).
  rewrite ← make_poly_pointless_r rewrite mulPP"\_refold rewrite ← IHp.
  unfold mulPP". rewrite make_poly_pointless_r. simpl. unfold mulMP' at 2.
  rewrite app_comm_cons rewrite ← make_poly_pointless_r.
  rewrite mulPP"\_refold rewrite IHq. unfold mulPP".
  rewrite make_poly_pointless_r. simpl. unfold mulMP' at 1.
  rewrite app_comm_cons rewrite app_assoc rewrite ← make_poly_pointless_r.
  rewrite mulPP"\_refold rewrite ← IHp. unfold mulPP".
  rewrite make_poly_pointless_r. simpl. rewrite (app_assoc (map _ (map _ q))).
  apply Permutation\_sort_eq apply nodup\_cancel\_Permutation.
  apply Permutation\_map rewrite make_mono\_app\_comm apply perm\_skip.
  apply Permutation\_app\_tail apply Permutation\_app\_comm.
Qed.

6.6.8 Axiom 8: Associativity of Multiplication

The eighth axiom states that, for all terms x, y, and z, (x \ast (y \ast z)) \downarrow_P = ((x \ast y) \ast z) \downarrow_P.

This one is also fairly complicated, so we will start small and build up to it. First, we prove a convenient side effect of make\_poly\_pointless, which allows us to simplify mulPP into a mulMP and a mulPP. Unlike commutativity, for this proof we opt to use the version of mulPP that includes a make\_poly in its mulMP, in addition to the map make\_mono version used previously.

**Lemma mulPP"\_cons : \forall q a p,**

\[ \text{make\_poly (mulMP' q a ++ mulPP" q p)} = \text{mulPP" q (a::p)}. \]

**Proof.**

\[ \text{intros q a p. unfold mulPP" rewrite make\_poly\_pointless\_r. auto.} \]

Qed.

Next is a deceptively easy lemma map\_app\_make\_poly, which is the primary application of nodup\_cancel\_map, proven in list\_util. It states that if we are applying make\_poly twice, we can remove the second application, even if there is a map app in between them. Clearly, here, the map app is in reference to mulMP.

**Lemma map\_app\_make\_poly : \forall m p,**

\[ \forall a, \text{In a p} \rightarrow \text{is\_mono a)} \rightarrow \text{make\_poly (map (app m) (make\_poly p)) = make\_poly (map (app m) p).} \]

**Proof.**

\[ \text{intros m p Hm. apply Permutation\_sort\_eq.} \]

apply Permutation\_trans with (l':=(nodup\_cancel Mono\_eq\_dec (map make\_mono (map (app m) (nodup\_cancel Mono\_eq\_dec (map make\_mono p)))))).

apply nodup\_cancel\_Permutation. repeat apply Permutation\_map.

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unfold make_poly. rewrite ← Permutation_MonoSort_. auto.
rewrite (no_map_make_mono p); auto. repeat rewrite map_map.
apply nodup_cancel_map.
Qed.

The **map_app_make_poly** lemma is then immediately applied here, to state that since **mulMP''** already applies **make_poly** to its result, we can remove any **make_poly** calls inside.

**Lemma mulMP''_make_poly**: \( \forall p \ m, \)
\[
(\forall a, \text{In} a p \rightarrow \text{is mono} a) \rightarrow \\
\text{mulMP''}(\text{make_poly} p) m = \\
\text{mulMP''} p m.
\]
**Proof.**
intrs p m. unfold mulMP''. apply map_app_make_poly.
Qed.

This very simple lemma states that since **mulMP** is effectively just a **map**, it distributes over **app**.

**Lemma mulMP'_app**: \( \forall p q m, \)
\[
\text{mulMP'}(p ++ q) m = \\
\text{mulMP'} p m ++ \text{mulMP'} q m.
\]
**Proof.**
intrs p q m. unfold mulMP'. repeat rewrite map_app. auto.
Qed.

Now into the meat of the associativity proof. We begin by proving that **mulMP'** is associative. This proof is straightforward, and is proven by induction with the use of **make_mono_pointless** and **Permutation_sort_mono_eq**.

**Lemma mulMP'_assoc**: \( \forall q a m, \)
\[
\text{mulMP'}(\text{mulMP'} q a) m = \\
\text{mulMP'}(\text{mulMP'} q m) a.
\]
**Proof.**
intrs q a m. unfold mulMP'. induction q; auto.
simpl. repeat rewrite make_mono_pointless. f_equal.
- apply Permutation_sort_mono_eq. apply Permutation_nodup.
  - apply Permutation_app_assoc. apply Permutation_app_tail.
  - apply Permutation_app_comm.
  - apply IHq.
Qed.

For the final associativity proof, we begin by using the commutativity lemma to make it so that \( q \) is on the leftmost side of the multiplications. This means that it will never be the polynomial being mapped across, and allows us to do induction on just \( p \) and \( r \) instead of all three. Thus \( p \) becomes \( a :: p \), and \( r \) becomes \( m :: r \).

The first three cases are easily solved with some rewrites and a call to auto, so we move on to the fourth. Similarly to the commutativity proof, the main struggle here is forcing
mulPP to map across the same term on both sides of the equation. This is accomplished in a very similar way - by simplifying, using make_poly_pointless to get mulPP back in the goal, and then applying the two induction hypotheses to reorder the terms.

The crucial point is when we rewrite with mulMP'_mulMP'', allowing us to wrap our mulMPs in make_poly and make use of the lemmas we proved earlier in this section. This technique enables us to reorder the multiplications in a way that is convenient for us; ((q*[a :: p]) * m) ↓P becomes ((q * a) * m) ↓P + (q * p) * m) ↓P. At the end of all of this rewriting, we are left with the original (p * q * r) ↓P as the last term of both sides, and (q * p * m) ↓P and (q * r * a) ↓P as the middle terms of both. These three terms are easily eliminated with the standard Permutation lemmas, because they are on both sides.

The only remaining challenge comes from the first term on each side; on the left, we have ((q * a) * m) ↓P, and on the right we have ((q * m) * a) ↓P. This is where the above mulMP'_assoc lemma comes into play, solving the last piece of the associativity lemma.

Lemma mulPP_assoc : ∀ p q r.
  mulPP (mulPP p q) r = mulPP p (mulPP q r).

Proof.
  intros p q r. rewrite (mulPP_comm _ (mulPP q _)). rewrite (mulPP_comm p _).
  generalize dependent r. induction p; induction r as [|m];
  repeat rewrite mulPP_0; repeat rewrite mulPP_0r; auto.
  repeat rewrite mulPP_mulPP'' in *. unfold mulPP''. simpl.
  repeat rewrite ← (make_poly_pointless_r _ (concat _)).
  repeat rewrite mulPP''_refold. repeat rewrite (mulPP''_cons q).
  pose (IHp (m::r)). repeat rewrite mulPP_mulPP'' in e. rewrite ← e.
  rewrite IHr. unfold mulPP'' at 2, mulPP'' at 4. simpl.
  repeat rewrite make_poly_pointless_r. repeat rewrite app_assoc.
  repeat rewrite ← (make_poly_pointless_r _ (concat _)).
  repeat rewrite mulPP''_refold. pose (IHp r).
  repeat rewrite mulPP_mulPP'' in e0. rewrite ← e0.
  repeat rewrite ← app_assoc. repeat rewrite mulMP'_mulMP''.
  repeat rewrite ← mulPP''_cons. repeat rewrite mulMP''_make_poly.
  repeat rewrite ← mulMP'_mulMP''. repeat rewrite app_assoc.
  apply Permutation_sort_eq. apply nodup_cancel_Permutation.
  apply Permutation_map. apply Permutation_app_tail. repeat rewrite mulMP'_app.
  rewrite mulMP'_assoc. repeat rewrite ← app_assoc. apply Permutation_app_head.
  apply Permutation_app_comm. intros a0 Hin. apply in_app_iff in Hin as ||.
  unfold mulMP' in H. apply in_map_iff in H as |x||. rewrite ← H; auto.
  apply (make_poly_is_poly (map (mulMP' q) r)). auto.
  intros a0 Hin. apply in_app_iff in Hin as ||. unfold mulMP' in H.
  apply in_map_iff in H as |x||. rewrite ← H; auto.
  apply (make_poly_is_poly (map (mulMP' q) p)). auto.
Qed.

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6.6.9 Axiom 9: Multiplicative Identity - Self

Next comes the other multiplicative identity mentioned earlier. This axiom states that for all terms \( x \), \( (x \times x) \downarrow_P = x \downarrow_P \).

To begin, we prove that this holds for monomials; \( (m \times m) \downarrow_P = m \downarrow_P \). This proof uses a combination of Permutation_Sorted_mono_eq and induction. We then use the standard Permutation lemmas to move the induction variable \( a \) out to the front, and show that nodup removes one of the two as. After that, perm_skip and the induction hypothesis solve the lemma.

Lemma make_mono_self : \( \forall m, is Mono m \rightarrow make_mono (m \times m) = m \).

Proof.
intros m H. apply Permutation_Sorted_mono_eq.
- induction m; auto. unfold make_mono. rewrite \(-\) Permutation_VarSort\_l.
  simpl. assert \((\ln a (m \times a : : m))\).
  intuition. destruct in_dec; try contradiction.
  apply Permutation_trans with \((l' := \text{nodup var_eq_dec (a : : m \times m)})\).
  apply Permutation_nodup. apply Permutation_app_comm.
  simpl. assert \((\ln a (m \times m))\).
  apply NoDup_VarSorted in H as H1. apply NoDup_cons_iff in H1.
  intro. apply H1. apply in_app_iff in H2; intuition.
  destruct in_dec; try contradiction. apply perm_skip.
  apply Permutation_VarSort\_l in IHm. auto. apply (mono_cons \_ \_ H).
- apply VarSort.LocallySorted_sort.
- apply Sorted_VarSorted. apply H.
Qed.

The full proof of the self multiplicative identity is much longer, but in a way very similar to the proof of commutativity. We begin by doing induction and simplifying, which distributes one of the induction variables across the list on the left side. This leaves us with \( a \times a \) as the leftmost term, which is easily replaced with \( a \) with the above lemma and then removed from both sides with perm_skip.

At this point we are left with a goal of the form \( (a \times [a :: p]) \downarrow_P + + ([a :: p] \times p) \downarrow_P = p \downarrow_P \) which is not particularly easy to deal with. However, by rewriting with mulPP_comm, we can force the second term on the left to simplify further.

This leaves us with something along the lines of \( (a \times [a :: p]) \downarrow_P + + (a \times [a :: p]) \downarrow_P + + (p \times p) \downarrow_P = p \downarrow_P \) which is much more workable! We know that \( (p \times p) \downarrow_P = p \downarrow_P \) from the induction hypothesis, so this is then removed from both sides and all that is left is to prove that the same term added together twice is equal to an empty list. This follows from the nodup_cancel_self lemma used to prove addPP_p_p, and finished the proof of this lemma.

Lemma mulPP_p_p : \( \forall p, is Poly p \rightarrow \)
mulPP \ p \ p = \ p.

Proof.
intros \ p \ H. rewrite mulPP_mulPP'. rewrite mulPP'_mulPP''.
apply Permutation_Sorted_eq.
- induction \ p; auto. unfold mulPP''; make_poly.
  rewrite \ Permutation_MonoSort_l. simpl map at 1.
  apply poly_cons in \ H as \ H1. destruct \ H1. rewrite make_mono_self; auto.
  rewrite no_make_mono; auto. rewrite map_app. apply Permutation_trans with
    (l':=nodup_cancel mono_eq_dec (map make_mono (concat (map (mulMP' (a ::
        p)) (a :: p)) ++ a :: map make_mono (map make_mono (map (app a) p)))).
  apply nodup_cancel_Permutation. rewrite app_comm_cons.
  apply Permutation_app_comm.
rewrite \ nodup_cancel_pointless. apply Permutation_trans with
  (l':=nodup_cancel mono_eq_dec ((nodup_cancel mono_eq_dec (map make_mono
    (concat (map (mulMP' p) (a :: p)) ++ a :: map make_mono (map make_mono
      (map (app a) p))))). apply nodup_cancel_Permutation.
  apply Permutation_app_tail. apply Permutation_sort_eq.
  repeat rewrite make_poly_refold. repeat rewrite mulPP''_refold.
  repeat rewrite \ mulPP'_mulPP''. repeat rewrite \ mulPP_mulPP'.
  apply mulPP_comm.
rewrite nodup_cancel_pointless. apply Permutation_trans with
  (l':=nodup_cancel mono_eq_dec (a :: map make_mono (map make_mono
    (map (app a) p)) ++ (a :: map make_mono (map make_mono
      (map (app a) p)))). apply nodup_cancel_Permutation.
  apply Permutation_app_comm. simpl map. rewrite map_app. unfold mulMP' at 1.
  repeat rewrite (no_map_make_mono (map make_mono _)).
  try apply mono_in_map_make_mono. rewrite (app_assoc (map _ _)).
  apply Permutation_trans with (l':=nodup_cancel mono_eq_dec ((map make_mono
    (map (app a) p) ++ map make_mono (map (app a) p)) ++ a :: map make_mono
    (map (mulMP' p) p)))). apply nodup_cancel_Permutation.
  apply Permutation_middle. rewrite \ nodup_cancel_pointless.
  rewrite nodup_cancel_self. simpl app.
  apply Permutation_trans with (l':=nodup_cancel mono_eq_dec (map make_mono
    (concat (map (mulMP' p) p)) ++ \ [a])). apply nodup_cancel_Permutation.
  replace a :: map make_mono (concat (map (mulMP' p) p)) with \ [a] ++ map
    make_mono (concat (map (mulMP' p) p))); auto. apply Permutation_app_comm.
  rewrite \ nodup_cancel_pointless. apply Permutation_trans with
    (l':=nodup_cancel mono_eq_dec (p ++ \ [a])). apply nodup_cancel_Permutation.
    apply Permutation_app_tail. unfold mulPP'', make_poly in \ Hp.
  rewrite \ Permutation_MonoSort_l in \ Hp. apply \ Hp; auto.
  replace a with \ [a]++p); auto. rewrite no_nodup_cancel_NoDup.
  apply Permutation_app_comm. apply Permutation_NoDup with (l:=a :: p).
6.6.10 Axiom 10: Distribution

Finally, we are left with the most intimidating of the axioms - distribution. This states, as one would expect, that for all terms \( x, y, \) and \( z \), \((x \times (y + z)) \downarrow_p = ((x \times y) + (x \times z)) \downarrow_p \).

In a similar approach to what we have done for some of the other lemmas, we begin by proving this on a smaller scale, working with just \texttt{mulMP} and \texttt{addPP}. This lemma is once again solved easily by the \texttt{map_app_make_poly} we proved while working on multiplication associativity, combined with \texttt{make_poly_pointless}.

**Lemma mulMP''_distr_addPP**: \( \forall m p q, \)
\[
is\text{poly } p \rightarrow is\text{poly } q \rightarrow \text{mulMP''}(\text{addPP } p q) m = \text{addPP}(\text{mulMP'' } p m)(\text{mulMP'' } q m).
\]

**Proof.**

\[
\text{intros } m p q \text{ } Hp \text{ } Hq. \text{ } \text{unfold } \text{mulMP''}, \text{addPP}. \text{rewrite } \text{map_app_make_poly}.
\]
\[
\text{rewrite } \text{make_poly_pointless}. \text{rewrite } \text{make_poly_app_comm}.
\]
\[
\text{rewrite } \text{make_poly_pointless}. \text{rewrite } \text{make_poly_app_comm}.
\]
\[
\text{rewrite } \text{map_app}. \text{auto. intros } a \text{ } Hin. \text{ } \text{apply in}_{app} \iff \text{in } Hin \text{ as } [].
\]
\[
\text{apply } Hp. \text{auto. apply } Hq. \text{auto.}
\]

Qed.

For the distribution proof itself, we begin by performing induction on \( r \), the element outside of the \texttt{addPP} call initially. We begin by simplifying, and using the usual combination of \texttt{make_poly_pointless} and refolding to convert our goal to a form of \((p + q) \times a) \downarrow_p + + \((p + r) \times r) \downarrow_p \).

We then apply similar tactics on the right side, to convert our goal to a form similar to \((p \times a + q \times a + p \times r + q \times r) \downarrow_p \). The two terms containing \( r \) are easy to deal with, since we know they are equal to the \((p + q) \times r) \downarrow_p \) we have on the left side due to the induction hypothesis. Similarly, the first two terms are known to be equal to \((p + q) \times a) \downarrow_p \) from the \texttt{mulMP_distr_addPP} lemma we just proved. This results in us having the same thing on both sides, thus solving the final of the ten \( B \)-unification axioms.

**Lemma mulPP_distr_addPP**: \( \forall p q r, \)
\[
is\text{poly } p \rightarrow is\text{poly } q \rightarrow \text{mulPP}(\text{addPP } p q) r = \text{addPP}(\text{mulPP } p r)(\text{mulPP } q r).
\]

**Proof.**

\[
\text{intros } p q r \text{ } Hp \text{ } Hq. \text{ } \text{induction } r; \text{auto. rewrite mulPP_mulPP''}. \text{unfold mulPP''}.
\]
\[
\text{simpl. rewrite mulPP_mulPP''}, (\text{mulPP_mulPP'' } q), \text{make_poly_app_comm}.
\]
\[
\text{rewrite } \text{← make_poly_pointless}. \text{rewrite } \text{make_poly_app_comm}.
\]
\[
\text{rewrite mulPP''_refold}.
\]
rewrite addPP_refold. repeat unfold mulPP'' at 2. simpl. unfold addPP at 4.
rewrite make_poly_pointless. rewrite addPP_refold.
rewrite (addPP_comm_ (make_poly _)).
unfold addPP at 4. rewrite make_poly_pointless. rewrite ← app_assoc.
rewrite make_poly_app_comm. rewrite ← app_assoc.
rewrite ← make_poly_pointless.
rewrite mulPP''_refold. rewrite ← app_assoc. rewrite app_assoc.
rewrite make_poly_app_comm.
rewrite ← app_assoc. rewrite ← make_poly_pointless. rewrite mulPP''_refold.
replace (make_poly (mulPP'' p r ++ mulMP' q a ++ mulPP'' q r ++ mulMP' p a))
   with (make_poly ((mulPP'' p r ++ mulPP'' q r) ++ mulMP' p a ++ mulMP' q a)).
rewrite ← make_poly_pointless. rewrite (addPP_refold (mulPP'' _ _)).
rewrite make_poly_app_comm. rewrite addPP_refold.
rewrite mulPP_mulPP'', (mulPP_mulPP'' p), (mulPP_mulPP'' q) in IHr.
rewrite ← IHr. unfold addPP at 4.
rewrite ← make_poly_pointless. unfold addPP. repeat rewrite mulMP'_mulMP''.
rewrite (make_poly_app_comm (mulMP'' _ _)) (mulMP' _ _)).
rewrite mulMP_mulMP''.
rewrite (make_poly_app_comm (mulMP'' _ _)) (mulMP'' _ _)).
repeat rewrite addPP_refold. f_equal. apply mulMP'_distr_addPP; auto.
apply make_poly_Permutation. rewrite ← app_assoc.
apply Permutation_app_head. rewrite app_assoc.
apply Permutation_trans with
   (l':=mulMP' q a ++ mulPP'' q r ++ mulMP' p a).
apply Permutation_app_comm.
auto.
Qed.

For convenience, we also prove that distribution can be applied from the right, which
follows from mulPP_comm and the distribution lemma we just proved.

Lemma mulPP_distr_addPPr : ∀ p q r,
is_poly p → is_poly q →
mulPP r (addPP p q) = addPP (mulPP r p) (mulPP r q).
Proof.
intros p q r Hp Hq. rewrite mulPP_comm. rewrite (mulPP_comm r p).
rewrite (mulPP_comm r q). apply mulPP_distr_addPP; auto.
Qed.

6.7 Other Facts About Polynomials

Now that we have proven the core ten axioms proven, there are a few more useful lemmas
that we will prove to assist us in future parts of the development.
6.7.1 More Arithmetic

Occasionally, when dealing with multiplication, we already know that one of the variables being multiplied in is less than the rest, meaning it would end up at the front of the list after sorting. For convenience and to bypass the work of dealing with the calls to sort and nodup-cancel, the below lemma allows us to rewrite with this concept.

Lemma mulPP_mono_cons : \forall x m, is_mono (x :: m) \rightarrow mulPP [[x]] [m] = [x :: m].

Proof.
intros x m H. unfold mulPP, distribute. simpl. apply Permutation_Sorted_eq.
- apply Permutation_trans with
  (l':=nodup_cancel mono_eq_dec (map make_mono [m ++ [x]])).
  apply Permutation_sym. apply Permuted_sort. rewrite no_nodup_cancel_NoDup.
  simpl. assert (make_mono (m ++ [x]) = x :: m).
  rewrite \rightarrow no_make_mono; auto. apply Permutation_sort_mono_eq.
  repeat rewrite no_nodup_NoDup. replace (x :: m) with ([x] ++ m); auto;
  apply Permutation_app_comm. apply NoDup_VarSorted; apply H.
  apply Permutation_NoDup with (l:=x :: m).
  replace (x :: m) with ([x] ++ m); auto; apply Permutation_app_comm.
  apply NoDup_VarSorted; apply H.
  + rewrite H0. auto.
  + apply NoDup_cons; auto.
- apply LocallySorted_sort.
- apply Sorted_cons; auto.
Qed.

Similarly, if we already know some monomial is less than the polynomials it is being added to, then the monomial will clearly end up at the front of the list.

Lemma addPP_poly_cons : \forall m p, is_poly (m :: p) \rightarrow addPP [m] p = m :: p.

Proof.
intros m p H. unfold addPP. simpl. rewrite no_make_poly; auto.
Qed.

An interesting arithmetic fact is that if we multiply the term \((p*q) + r) \downarrow_P by \((1+q) \downarrow_P, we effectively eliminate the \((p * q) \downarrow_P term and are left with \(((1 + q) * r) \downarrow_P. This will come into play later in the development, as we look to begin building unifiers.

Lemma mulPP_addPP_1 : \forall p q r, is_poly p \rightarrow is_poly q \rightarrow is_poly r \rightarrow mulPP (addPP (mulPP p q) r) (addPP [[]] q) = mulPP (addPP [[]] q) r.

Proof.

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6.7.2 Reasoning about Variables

To more easily deal with the \texttt{vars} definition, we have defined a few definitions about it. First, if some \(x\) is in the variables of \texttt{make\_poly} \(p\), then it must have been in the \texttt{vars} of \(p\) originally. Note that this is not true in the other direction, as \texttt{nodup\_cancel} may remove some variables.

\textbf{Lemma make\_poly\_rem\_vars}: \(\forall p\ x,\)
\[\text{In } x (\text{vars (make\_poly} p)) \rightarrow \text{In } x (\text{vars } p).\]
\textbf{Proof}.
\begin{itemize}
  \item intros \(p\ x\ H.\ induction\ p.\)
  \item - inversion \(H.\)
  \item - unfold vars. simpl. apply nodup\_ln. apply in\_app\_iff.
    \begin{itemize}
      \item unfold vars, make\_poly in \(H.\) apply nodup\_ln in \(H.\)
      \item apply ln\_concat\_exists in \(H\) as \([m \,|\,].\)
      \item apply ln\_sorted in \(H.\) apply nodup\_cancel\_in in \(H.\)
      \item apply in\_map\_iff in \(H\) as \([n \,|\,].\) destruct \(H1.\)
      \item + left. apply make\_mono\_ln. rewrite \(H1.\) rewrite \(H.\) auto.
      \item + right. apply ln\_concat\_exists. \(\exists n.\) split; auto. apply make\_mono\_ln.
        rewrite \(H.\) auto.
    \end{itemize}
\end{itemize}
\texttt{Qed.}

An interesting observation about \texttt{addPP} and our \texttt{vars} function is that clearly, the variables of some \((p + q) \downarrow_p\) is a subset of the variables of \(p\) combined with the variables of \(q\). The next lemma is a more convenient formulation of that fact, using a list of variables \(xs\) rather than comparing them directly.

\textbf{Lemma incl\_vars\_addPP}: \(\forall p\ q\ xs,\)
\[\text{incl (vars } p\) \,xs \wedge \text{incl (vars } q\) \,xs \rightarrow \text{incl (vars (addPP } p\ q)) \,xs.\]
\textbf{Proof}.
\begin{itemize}
  \item unfold incl, addPP.
  \item intros \(p\ q\ xs\ [HinP\ HinQ] x\ HinPQ.\)
  \item apply make\_poly\_rem\_vars in \(HinPQ.\)
  \item unfold vars in \(HinPQ.\)
  \item apply nodup\_ln in \(HinPQ.\)
  \item rewrite concat\_app in \(HinPQ.\)
  \item apply in\_app\_or in \(HinPQ\) as \([Hin\ |\ Hin].\)
  \item - apply \(HinP.\) apply nodup\_ln. auto.
\end{itemize}
- apply HinQ. apply nodup_in. auto.
Qed.

We would like to be able to prove a similar fact about \texttt{mulPP}, but before we can do so, we need to know more about the \texttt{distribute} function. This lemma states that if some \(a\) is in the variables of \(\texttt{distribute} \ l \ m\), then it must have been in either \(\texttt{vars} \ l\) or \(\texttt{vars} \ m\) originally.

\textbf{Lemma} \(\texttt{ln_distribute} : \forall (l \ m:\texttt{poly}) \ a,\)
\(\texttt{ln} \ a (\texttt{vars} (\texttt{distribute} \ l \ m)) \to \texttt{ln} \ a (\texttt{vars} \ l) \lor \texttt{ln} \ a (\texttt{vars} \ m).\)
\textbf{Proof}.
intros \(l \ m \ a \ H\). unfold \texttt{distribute}, \texttt{vars} in \(H\).
apply \(\texttt{ln_concat_exists} \in H\). destruct \(H\) as \([\texttt{ll}]\).
apply \(\texttt{ln_concat_exists} \in H\). destruct \(H\) as \([\texttt{ll1}]\).
apply \(\texttt{ln_map_iff} \in H\). destruct \(H\) as \([x][]\). rewrite \(\leftarrow \ H\) in \(H1\).
apply \(\texttt{ln_app_iff} \in H1\). destruct \(H1\) as \([x0][]\). rewrite \(\leftarrow \ H1\) in \(H0\).
- right. apply nodup_in. apply \(\texttt{ln_concat_exists}\). \(\exists x.\) auto.
- left. apply nodup_in. apply \(\texttt{ln_concat_exists}\). \(\exists x0.\) auto.
Qed.

We can then use this fact to prove our desired fact about \texttt{mulPP}; the variables of \(p\star q \downarrow P\) are a subset of the variables of \(p\) and the variables of \(q\). Once again, this is formalized in a way that is more convenient in later proofs, with an extra list \(xs\).

\textbf{Lemma} \(\texttt{incl_vars_mulPP} : \forall p \ q \ xs,\)
\(\texttt{incl} (\texttt{vars} \ p) \ xs \land \texttt{incl} (\texttt{vars} \ q) \ xs \to \texttt{incl} (\texttt{vars} (\texttt{mulPP} \ p \ q)) \ xs.\)
\textbf{Proof}.
unfold \texttt{incl}, \texttt{mulPP}.
intros \(p \ q \ xs \ [\texttt{HinP} \ HinQ] \ x \ HinPQ\).
apply \(\texttt{make_poly_rem_vars} \in \texttt{HinPQ}\).
apply \(\texttt{ln_distribute} \in \texttt{HinPQ}\). destruct \(\texttt{HinPQ}\).
- apply \(\texttt{HinP}\). auto.
- apply \(\texttt{HinQ}\). auto.
Qed.

\section{6.7.3 Partition with Polynomials}

When it comes to actually performing successive variable elimination later in the development, the \texttt{partition} function will play a big role, so we have opted to prove a few useful facts about its relation to polynomials now.

First is that if you separate a polynomial with any function \(f\), you can get the original polynomial back by adding together the two lists returned by \texttt{partition}. This is relatively easy to prove thanks to the lemma \texttt{partition_Permutation} we proved during \texttt{list_util}.
Lemma part_add_eq : ∀ f p l r,
  is_poly p →
  partition f p = (l, r) →
  p = addPP l r.
Proof.
  intros f p l r H H0. apply Permutation_Sorted_eq.
  - generalize dependent l; generalize dependent r. induction p; intros.
    + simpl in H0. inversion H0. auto.
    + assert (H1:=H0); auto. apply partition_Permutation in H1. simpl in H0.
      destruct (partition f p) as [g d]. unfold addPP, make_poly.
      rewrite ← Permutation_MonoSort_r. rewrite unsorted_poly. destruct (f a);
      inversion H0.
      × rewrite ← H3 in H1. apply H1.
      × rewrite ← H4 in H1. apply H1.
      × destruct H. apply NoDup_MonoSorted in H. apply (Permutation_NoDup H1 H).
      × intros m Hin. apply H. apply Permutation_sym in H1.
        apply (Permutation_in _ H1 Hin).
      - apply Sorted_MonoSorted. apply H.
      - apply Sorted_MonoSorted. apply make_poly_is_poly.
  Qed.

In addition, if you partition some polynomial p with any function f, the resulting two lists will both be proper polynomials, since partition does not affect the order.

Lemma part_is_poly : ∀ f p l r,
  is_poly p →
  partition f p = (l, r) →
  is_poly l ∧ is_poly r.
Proof.
  intros f p l r Hpoly Hpart. destruct Hpoly. split; split.
  - apply (part_Sorted _ _ mono_lt_Transitive H _ _ Hpart).
  - intros m Hin. apply H0. apply elements_in_partition with (x:=m) in Hpart.
    apply Hpart; auto.
  - apply (part_Sorted _ _ mono_lt_Transitive H _ _ Hpart).
  - intros m Hin. apply H0. apply elements_in_partition with (x:=m) in Hpart.
    apply Hpart; auto.
Qed.

### 6.7.4 Multiplication and Remove

Lastly are some rather complex lemmas relating remove and multiplication. Similarly to the partition lemmas, these will come to play a large roll in performing successive variable elimination later in the development.

First is an interesting fact about removing from monomials. If there are two monomials
which are equal after removing some \( x \), and either both contain \( x \) or both do not contain \( x \), then they must have been equal originally. This proof begins by performing double induction, and quickly solving the first three cases.

The fourth case is rather long, and begins by comparing if the \( a \) and \( a_0 \) at the head of each list are equal. The case where they are equal is relatively straightforward; we must also destruct if \( x = a = a_0 \), but regardless of whether they are equal or not, we can easily prove this with the use of the induction hypothesis.

The case where \( a \neq a_0 \) should be a contradiction, as that element is at the head of both lists, and we know the lists are equal after removing \( x \). We begin by comparing if the \( x \) is in the two lists. In the case where it is not in either, we can quickly solve this, as we know the call to remove will do nothing, which immediately gives us the contradiction.

In the case where \( x \) is in both, we begin by using \texttt{in\_split} to rewrite both lists to contain \( x \). We then use the fact that there are no duplicates in either list to show that \( x \) is not in \( l_1, l_2, l_1', \text{ or } l_2' \), and therefore the calls to remove will do nothing. This leaves us with a hypothesis that \( l_1 ++ l_2 = l_1' ++ l_2' \). To finish the proof, we destruct \( l_1 \) and \( l_1' \) to further compare the head of each list.

In the case where they are both empty, we arrive at a contradiction immediately, as this implies the head of both lists is \( x \) and therefore contradicts that \( a \neq a_0 \). In the case where they are both lists, doing inversion on our remove hypothesis gives us that the head of each list is equal again, also contradicting that \( a \neq a_0 \).

In the other two cases, we rewrite with the \texttt{in\_split} hypotheses into the \texttt{is\_mono} hypotheses. In both cases, we result in one statement that \( a \) comes before \( a_0 \) in the monomial, and one statement that \( a_0 \) comes before \( a \) in the monomial. With the help of \texttt{StronglySorted}, we are able to turn these into \( a < a_0 \) and \( a_0 < a \), which contradict each other to finish the proof.

Lemma \texttt{remove\_Sorted\_eq} : \( \forall x \ (l \ l':\text{mono}), \)
\begin{align*}
\text{is\_mono } l &\rightarrow \text{is\_mono } l' \rightarrow \\
\text{In } x \ l &\leftrightarrow \text{In } x \ l' \rightarrow \\
\text{remove var\_eq\_dec } x \ l = \text{remove var\_eq\_dec } x \ l' \rightarrow \\
l = l'.
\end{align*}

Proof.
\begin{verbatim}
intros x l l' Hl Hl' Hx Hrem.
 generalize dependent l'; induction l; induction l'; intros.
- auto.
- destruct (var\_eq\_dec x a) eqn:Heq.
  + rewrite e in Hx. exfalso. apply Hx. intuition.
  + simpl in Hrem. rewrite Heq in Hrem. inversion Hrem.
- destruct (var\_eq\_dec x a) eqn:Heq.
  + rewrite e in Hx. exfalso. apply Hx. intuition.
  + simpl in Hrem. rewrite Heq in Hrem. inversion Hrem.
- clear IHl'. destruct (var\_eq\_dec a a0).
  + rewrite e. f_equal. rewrite e in Hrem. simpl in Hrem.
\end{verbatim}

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apply mono_cons in Hl as Hl1. apply mono_cons in Hl' as Hl'1.
destruct (var_eq_dec x a0).

× apply IHl; auto. apply NoDup_VarSorted in Hl.
apply NoDup_cons_iff in Hl. rewrite e in Hl. rewrite \leftrightarrow e0 in Hl.
destruct Hl. split; intro. contradiction. apply NoDup_VarSorted in Hl'.
apply NoDup_cons_iff in Hl'. rewrite \leftrightarrow e0 in Hl'. destruct Hl'.
contradiction.

× inversion Hrem. apply IHl; auto. destruct Hx. split; intro. simpl in H.
rewrite e in H. destruct H; auto. rewrite H in n. contradiction.
simpl in H1. rewrite e in H1. destruct H1; auto. rewrite H1 in n.
contradiction.
+ destruct (in_dec var_eq_dec x (a::l)).

× apply Hx in i as i'. apply in_split in i. apply in_split in i'.
destruct i as [l1[l2 i]]. destruct i' as [l1'l2' i'].
pose (NoDup_VarSorted _ Hl). pose (NoDup_VarSorted _ Hl').
apply (NoDup_Ln_split _ _ _ i) in n0 as [].
apply (NoDup_Ln_split _ _ _ i') in n1 as [].
rewrite i in Hrem. rewrite i' in Hrem.
repeat rewrite remove_distr_app in Hrem. simpl in Hrem.
destruct (var_eq_dec x x); try contradiction.
repeat (rewrite not_Ln_remove in Hrem; auto). destruct l1; destruct l1';
simpl in i; simpl in i'; simpl in Hrem; inversion i; inversion i'.
- rewrite H4 in n. rewrite H6 in n. contradiction.
- rewrite H7 in Hl. rewrite i in Hl. rewrite Hrem in Hl.
  rewrite H6 in Hl'. assert (x < v). apply Sorted_inv in Hl as [].
  apply HdRel_inv in H8. auto. assert (v < x).
  apply StronglySorted in Hl'.
  apply StronglySorted_inv in Hl' as []. rewrite Forall_forall in H9.
  apply H9. intuition. apply lt_Transitive. apply lt_asymm in H8.
  contradiction.
- rewrite H7 in Hl'. rewrite i in Hl. rewrite \leftrightarrow Hrem in Hl'.
  rewrite H6 in Hl'. assert (n0 < x).
  apply Sorted_StronglySorted in Hl.
  apply StronglySorted_inv in Hl as []. rewrite Forall_forall in H8.
  apply H8. intuition. apply lt_Transitive. assert (x < n0).
  apply Sorted_inv in Hl' as []. apply HdRel_inv in H9; auto.
  apply lt_asymm in H8. contradiction.
- inversion Hrem. rewrite \leftrightarrow H4 in H8. rewrite \leftrightarrow H6 in H8.
  contradiction.
× assert (\neg in x (a0::l')). intro. apply n0. apply Hx. auto.
repeat (rewrite not_Ln_remove in Hrem; auto).

Qed.
Next is that if we map remove across a polynomial where every monomial contains \( x \), there will still be no duplicates at the end.

**Lemma NoDup_map_remove**: \( \forall x \, p, \)
\[
\text{is\_poly } p \rightarrow \\
(\forall m, \ln m \, p \rightarrow \ln x \, m) \rightarrow \\
\text{NoDup } (\text{map } (\text{remove } \text{var\_eq\_dec } x) \, p).
\]

**Proof**.

intros \( x \, p \, Hx \). induction \( p \); simpl; auto.
apply NoDup\_cons.
- intro. apply \text{in\_map\_iff} in \( H \). destruct \( H \) as \([y \, |] \). assert \((y = a)\).
  + apply \text{poly\_cons} in \( Hp \). destruct \( Hp \). unfold \text{is\_poly} in \( H1 \). destruct \( H1 \).
    apply \( H3 \) in \( H0 \) as \( H4 \). apply \((\text{remove\_Sorted\_eq } x)\); auto. split; intro.
    apply \( Hx \). intuition. apply \( Hx \). intuition.
  + rewrite \( H1 \) in \( H0 \). unfold \text{is\_poly} in \( Hp \). destruct \( Hp \).
    apply NoDup\_MonoSorted in \( H2 \) as \( H4 \). apply NoDup\_cons\_iff in \( H4 \) as \(| \).
    contradiction.
- apply \( IHp \).
  + apply \text{poly\_cons} in \( Hp \). apply \( Hp \).
  + intros \( m \, H \). apply \( Hx \). intuition.
Qed.

Building off that, if every monomial in a list does not contain some \( x \), then appending \( x \) to every monomial and calling \text{make\_mono} still will not create any duplicates.

**Lemma NoDup_map_app**: \( \forall x \, l, \)
\[
\text{is\_poly } l \rightarrow \\
(\forall m, \ln m \, l \rightarrow \neg \ln x \, m) \rightarrow \\
\text{NoDup } (\text{map } \text{make\_mono } (\text{map } (\text{fun } a \Rightarrow a \, ++ \, [x]) \, l)).
\]

**Proof**.

intros \( x \, l \, Hp \, Hin \). induction \( l \).
- simpl. auto.
- simpl. apply NoDup\_cons.
  + intros \( H \). rewrite \text{map\_map} in \( H \). apply \text{in\_map\_iff} in \( H \) as \([m \, |] \).
    assert \((a = m)\).
  \times apply \text{poly\_cons} in \( Hp \) as \(|\). apply \text{Permutation\_Sorted\_mono\_eq}.
    + apply \text{Permutation\_sort\_mono\_eq} in \( H \). rewrite \text{no\_nodup\_NoDup} in \( H \).
      rewrite \text{no\_nodup\_NoDup} in \( H \).
      ++ pose \((\text{Permutation\_cons\_append } m \, x)\).
      pose \((\text{Permutation\_cons\_append } a \, x)\).
      apply \( \text{Permutation\_trans } p \) in \( H \). apply \text{Permutation\_sym} in \( p0 \).
      apply \( \text{Permutation\_trans } H \) in \( p0 \).
      apply \text{Permutation\_cons\_inv} in \( p0 \). apply \text{Permutation\_sym}. auto.
      ++ apply \text{Permutation\_NoDup} with \((l:=x :: a)\).
      apply \text{Permutation\_cons\_append}.

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apply NoDup_cons. apply Hin. intuition. unfold is_mono in H2.
apply NoDup_VarSorted in H2. auto.

++ apply Permutation_NoDup with \( (l := x :: m) \).
apply Permutation_cons_append. apply NoDup_cons. apply Hin.
intuition. unfold is_poly in H1. destruct H1. apply H3 in H0.
unfold is_mono in H0. apply NoDup_VarSorted in H0. auto.
- unfold is_mono in H2. apply Sorted_VarSorted. auto.
- unfold is_poly in H1. destruct H1. apply H3 in H0.
apply Sorted_VarSorted. auto.
× rewrite ← H1 in H0. unfold is_poly in Hp. destruct Hp.
apply NoDupMonoSorted in H2. apply NoDup_cons_iff in H2 as \[].
contradiction.
+ apply IHl. apply poly_cons in Hp. apply Hp. intros m H. apply Hin.
intuition.
Qed.

This next lemma is relatively straightforward, and really just served to remove the calls to sort and nodup_cancel for convenience when simplifying a mulPP.

Lemma mulPP_Permutation : \( \forall \ x \ a0 \ l, \)
\[
is\_poly \ ((a0 :: l) \rightarrow
(\forall \ m, \text{ln } m \ (a0 :: l) \rightarrow \neg \text{ln } x \ m) \rightarrow
\text{Permutation} \ ((\text{mulPP } [[x]]) (a0 :: l))
((\text{make\_mono} \ ((a0 ++ [x]))) :: (\text{mulPP } [[x]]) l)).
\]

Proof.
intros x a0 l Hp Hx. unfold mulPP, distribute. simpl. unfold make_poly.
pose (MonoSort.Permut_sorted (nodup_cancel mono_eq_dec

(map make_mono \((a0 ++ [x]) :: \text{concat} \ ((\text{map } \ a \Rightarrow [a ++ [x]]) l))))).
apply Permutation_trans p. simpl map.
rewrite no_nodup_cancel_NoDup; clear p.
- apply perm_skip. rewrite ← PermutationMonoSort_r.
rewrite no_nodup_cancel_NoDup; auto. rewrite concat_map.
apply NoDup_map_app. apply poly_cons in Hp. apply Hp. intros m H. apply Hx.
intuition.
- rewrite ← map_cons. rewrite concat_map.
rewrite ← map_cons with \((f := \text{fun } a \Rightarrow a ++ [x])\).
apply NoDup_map_app; auto.
Qed.

Building off of the previous lemma, this one serves to remove the calls to make_poly entirely, and instead replace mulPP with just the map app. We can do this because we know that \( x \) is not in any of the monomials, so nodup_cancel will have no effect as we proved earlier.

Lemma mulPP_map_app_permutation : \( \forall \ (x:\text{var}) \ (l \ l' : \text{poly}), \)
is_poly l →
(∀ m, ln m l → ¬ ln x m) →
Permutation l l’ →
Permutation (mulPP [[x]] l) (map (fun a ⇒ (make_mono (a ++ [x]))) l’).

Proof.
intros x l l’ Hp H0. generalize dependent l’. induction l; induction l’.
- intros. unfold mulPP, distribute, make_poly, MonoSort.sort. simpl. auto.
- intros. apply Permutation_nil_cons in H0. contradiction.
- intros. apply Permutation_sym in H0. apply Permutation_nil_cons in H0.
  contradiction.
- intros. clear IHL’. destruct (mono_eq_dec a a0).
  + rewrite e in *. pose (mulPP.Permutation x a0 l Hp H).
    apply (Permutation_trans p). simpl. apply perm_skip. apply IHL.
    - clear p. apply poly_cons in Hp. apply Hp.
    - intros m Hinl. apply H. intuition.
    - apply Permutation_cons_inv in H0. auto.
  + apply Permutation_incl in H0 as H1. destruct H1.
    apply incl_cons_inv in H1 as H1. destruct H1;
    try (rewrite H1 in v; contradiction). apply in_split in H1.
    destruct H1 as [l1 l2]. rewrite H1 in H0.
    pose (Permutation_middle (a0 :: l1) l2 a). apply Permutation_sym in p.
    simpl in p. apply (Permutation_trans H0) in p.
    apply Permutation_cons_inv in p. rewrite H1. simpl. rewrite map_app.
    simpl. pose (Permutation_middle ((make_mono (a0 ++ [x])) :: map (fun a1 ⇒
      make_mono (a1 ++ [x])) l1)) (map (fun a1 ⇒ make_mono (a1 ++ [x])) l2)
      (make_mono (a++[x])).
    simpl in p0. simpl. apply Permutation_trans with (l’:=make_mono (a ++ [x])
      :: make_mono (a0 ++ [x]) :: map (fun a1 : list var ⇒ make_mono (a1 ++
      [x])) l1 ++ map (fun a1 : list var ⇒ make_mono (a1 ++ [x])) l2); auto.
    clear p0. rewrite (map_cons (fun a1 ⇒ make_mono (a1 ++ [x])) a0 (@app (list var)
      l1 l2))
    pose (mulPP.Permutation x a l Hp H). apply (Permutation_trans p0).
    apply perm_skip. apply IHL.
    - clear p0. apply poly_cons in Hp. apply Hp.
    - intros m Hinl. apply H. intuition.
    - apply p.

Qed.

Finally, we combine the lemmas in this section to prove that, if there is some polynomial p
that has x in every monomial, removing and then re-appending x to every monomial results
in a list that is a permutation of the original polynomial.

Lemma map_app_remove_Permutation : ∀ p x,
\[
\text{is\_poly } p \rightarrow \\
(\forall m, \text{ln } m \rightarrow \text{ln } x m) \rightarrow \\
\text{Permutation } p \ (\text{map } (\text{fun } a \Rightarrow (\text{make\_mono } (a + [x]))) \\
(\text{map } (\text{remove } \text{var\_eq\_dec } x) p)).
\]

Proof.
intro p x H H0. rewrite map_map. induction p; auto.
simpl. assert (\text{make\_mono } (@\text{app } \text{var} (\text{remove } \text{var\_eq\_dec } x) [x]) = a).
- unfold make\_mono. rewrite no\_nodup\_NoDup.
  + apply Permutation\_Sorted\_mono\_eq.
    \times apply Permutation\_trans \text{ with } (l' := \text{remove } \text{var\_eq\_dec } x a + [x]).
    apply Permutation\_sym. apply VarSort.Permutated\_sort.
    pose (in\_split x a). destruct e as [l1 [l2 e]]. apply H0. intuition.
    rewrite e. apply Permutation\_trans \text{ with }
    \times apply Permutation\_trans \text{ with } (l' := x :: \text{remove } \text{var\_eq\_dec } x (l1 ++ x :: l2)).
    apply Permutation\_sym. apply Permutation\_cons\_append.
    apply Permutation\_trans \text{ with } (l' := (x :: [l1 + l2])). apply perm\_skip.
    rewrite remove\_distr\_app. replace (x :: l2) with ([x] ++ l2); auto.
    rewrite remove\_distr\_app. simpl. destruct (var\_eq\_dec x x);
    try contradiction. rewrite app\_nil_l. repeat rewrite not\_ln\_remove;
    try apply Permutation\_refl; try (apply poly\_cons in H as []);
    unfold is\_mono in H1; apply NoDup\_VarSorted in H1; rewrite e in H1;
    apply NoDup\_remove\_2 in H1). intros x2. apply H1. intuition. intros x1.
    apply H1. intuition. apply Permutation\_middle.
  \times apply VarSort.LocallySorted\_sort.
  \times apply poly\_cons in H as []. unfold is\_mono in H1.
  apply Sorted\_VarSorted. auto.
  + apply Permutation\_NoDup with (l := (x :: \text{remove } \text{var\_eq\_dec } x a)).
    apply Permutation\_cons\_append. apply NoDup\_cons.
    apply remove\_ln. apply NoDup\_remove. apply poly\_cons in H as [].
    unfold is\_mono in H1. apply NoDup\_VarSorted. auto.
  - rewrite H1. apply perm\_skip. apply IHp.
    + apply poly\_cons in H. apply H.
    + intros m Hin. apply H0. intuition.
Qed.
Chapter 7

Library B_Unification.poly_unif

Require Import List.
Import ListNotations.
Require Import Arith.
Require Import Permutation.
Require Export poly.

7.1 Introduction

This section deals with defining substitutions and their properties using a polynomial representation. As with the inductive term representation, substitutions are just lists of replacements, where variables are swapped with polynomials instead of terms. Crucial to the proof of correctness in the following chapter, substitution is proven to distribute over polynomial addition and multiplication. Definitions are provided for unifier, unifiable, and properties relating multiple substitutions such as more general and composition.

7.2 Substitution Definitions

A substitution is defined as a list of replacements. A replacement is just a tuple of a variable and a polynomial.

Definition repl := prod var poly.
Definition subst := list repl.

Since the poly data type doesn’t enforce the properties of actual polynomials, the is_poly predicate is used to check if a term is in polynomial form. Likewise, the is_poly_subst predicate below verifies that every term in the range of the substitution is a polynomial.

Definition is_poly_subst (s : subst) : Prop :=
  \forall x p, \ln(x, p) s \rightarrow is_poly p.
The next three functions implement how substitutions are applied to terms. At the
top level, \texttt{substP} applies a substitution to a polynomial by calling \texttt{substM}
on each monomial. From there, \texttt{substV} is called on each variable. Because variables
and monomials are converted to polynomials, the process isn’t simplying mapping
application across the lists. \texttt{substM} and \texttt{substP} must multiply and add each
polynomial together respectively.

\begin{verbatim}
Fixpoint substV (s : subst) (x : var) : poly :=
  match s with
  | [] ⇒ [[x]]
  | (y, p) :: s' ⇒ if (x =? y) then p else (substV s' x)
  end.

Fixpoint substM (s : subst) (m : mono) : poly :=
  match m with
  | [] ⇒ [[]]
  | x :: m ⇒ mulPP (substV s x) (substM s m)
  end.

Definition substP (s : subst) (p : poly) : poly :=
  make_poly (concat (map (substM s) p)).
\end{verbatim}

Useful in later proofs is the ability to rewrite the unfolded definition of \texttt{substP}
as just the function call.

\begin{verbatim}
Lemma substP_refold : ∀ s p,
  make_poly (concat (map (substM s) p)) = substP s p.
Proof. auto. Qed.
\end{verbatim}

The following lemmas state that substitution applications always produce polynomials.
This fact is necessary for proving distribution and other properties of substitutions.

\begin{verbatim}
Lemma substV_is_poly : ∀ x s,
  is_poly_subst s →
  is_poly (substV s x).
Proof.
  intros x s H. unfold is_poly_subst in H. induction s; simpl; auto.
  destruct a eqn:Ha. destruct (x =? v).
  - apply (H v). intuition.
  - apply IHs. intros x0 p0 H0. apply (H x0). intuition.
Qed.

Lemma substM_is_poly : ∀ s m,
  is_poly (substM s m).
Proof.
  intros s m. unfold substM; destruct m; auto.
Qed.

Lemma substP_is_poly : ∀ s p,
  is_poly (substP s p).
\end{verbatim}
Proof.
intros. unfold substP. auto.
Qed.

Hint Resolve substP_is_poly substM_is_poly.

The lemma below states that a substitution applied to a variable in polynomial form is equivalent to the substitution applied to just the variable. This fact only holds when the substitution’s range consists of polynomials.

Lemma subst_var_eq : ∀ x s,
  is_poly subst s → substP s [[x]] = substV s x.
Proof.
  intros. simpl.
  apply (substV_is_poly x s) in H. unfold substP. simpl. rewrite app-nil_r. rewrite mulPP_1r; auto. rewrite no_make_poly; auto.
Qed.

The next two lemmas deal with simplifying substitutions where the first replacement tuple is useless for the given term. This is the case when the variable being replaced is not present in the term. It allows the replacement to be dropped from the substitution without changing the result.

Lemma substM_cons : ∀ x m,
  ¬ ln x m → ∀ p s, substM ((x, p) :: s) m = substM s m.
Proof.
  intros. induction m; auto. simpl. f_equal.
  - destruct (a =? x) eqn:H0; auto.
    symmetry in H0. apply beq_nat_eq in H0. exfalso.
    simpl in H. apply H. left. auto.
  - apply IHm. intro. apply H. right. auto.
Qed.

Lemma substP_cons : ∀ x p,
  (∀ m, ln m p → ¬ ln x m) → ∀ q s, substP ((x, q) :: s) p = substP s p.
Proof.
  intros. induction p; auto. unfold substP. simpl.
  repeat rewrite ← (make_poly_pointless_r_ (concat _)). f_equal. f_equal.
  - apply substM_cons. apply H. left. auto.
  - apply IHp. intros. apply H. right. auto.
Qed.

Substitutions applied to constants have no effect.

Lemma substP_1 : ∀ s,


\[ \text{substP } s \; [] = [] \].

Proof.
\begin{itemize}
  \item intros. unfold substP. simpl. auto.
\end{itemize}
Qed.

Lemma substP_0 : \( \forall s, \)  
\[ \text{substP } s \; [] = [] \].

Proof.
\begin{itemize}
  \item intros. unfold substP. simpl. auto.
\end{itemize}
Qed.

The identity substitution—the empty list—has no effect when applied to a term.

Lemma empty_substM : \( \forall m, \)  
\[ \text{is\_mono } m \rightarrow \]  
\[ \text{substM } [] \; m = [m] \].

Proof.
\begin{itemize}
  \item intros. induction \( m \); auto. simpl.
  \item apply mono_cons in \( H \) as \( H0 \).
  \item rewrite \( IHm \); auto.
  \item apply mulPP\_mono\_cons; auto.
\end{itemize}
Qed.

Lemma empty_substP : \( \forall p, \)  
\[ \text{is\_poly } p \rightarrow \]  
\[ \text{substP } [] \; p = p \].

Proof.
\begin{itemize}
  \item intros. induction \( p \); auto. unfold substP. simpl.
  \item apply poly_cons in \( H \) as \( H0 \). destruct \( H0 \).
  \item rewrite \( \leftarrow \) make\_poly\_pointless\_r. rewrite substP\_refold.
  \item rewrite \( IHp \); auto. rewrite empty_substM; auto.
  \item apply addPP\_poly\_cons; auto.
\end{itemize}
Qed.

### 7.3 Distribution Over Arithmetic Operators

Below is the statement and proof that substitution distributes over polynomial addition. Given a substitution \( s \) and two terms in polynomial form \( p \) and \( q \), it is shown that \( s(p+q) \downarrow_p = (s(p) + s(q)) \downarrow_p \). The proof relies heavily on facts about permutations proven in the list\_util library.

Lemma substP\_distr\_addPP : \( \forall p \; q \; s, \)  
\[ \text{is\_poly } p \rightarrow \]  
\[ \text{is\_poly } q \rightarrow \]  
\[ \text{substP } s \; (\text{addPP } p \; q) = \text{addPP } (\text{substP } s \; p) \; (\text{substP } s \; q) \].

Proof.
intros p q s Hp Hq. unfold substP, addPP.
apply Permutation_sort_eq. apply Permutation_trans with (l':=(nodup_cancel mono_eq_dec (map make_mono (concat (map (substM s) (p)))) ++ (nodup_cancel mono_eq_dec (map make_mono (concat (map (substM s) q)))))
apply nodup_cancel_Permutation. apply Permutation_map.
apply Permutation_concat. apply Permutation_map. unfold make_poly.
rewrite ← Permutation_MonoSortI. auto.
apply Permutation_sym. apply Permutation_trans with (l':=(nodup_cancel mono_eq_dec (map make_mono (concat (map (substM s) (p)))) ++ (nodup_cancel mono_eq_dec (map make_mono (concat (map (substM s) q)))))). apply nodup_cancel_Permutation.
apply Permutation_map. apply Permutation_app; unfold make_poly;
rewrite ← Permutation_MonoSortI; auto.
rewrite (no_map_make_mono ((nodup_cancel _ _) ++ (nodup_cancel _ _)}).
rewrite nodup_cancel_pointless. apply Permutation_trans with (l':=(nodup_cancel mono_eq_dec (nodup_cancel mono_eq_dec (map make_mono (concat (map (substM s) q)) ++ map make_mono (concat (map (substM s) p)))))
apply nodup-cancel_Permutation. apply Permutation_app_comm.
rewrite nodup_cancel_pointless. rewrite ← map_app. rewrite ← concat_app.
rewrite ← map_app. rewrite (no_map_make_mono (p++q)).
apply Permutation_trans with (l':=(nodup_cancel mono_eq_dec (map make_mono (concat (map (substM s) (p ++ q)))))
apply nodup_cancel_Permutation.
apply Permutation_map. apply Permutation_concat. apply Permutation_map.
apply Permutation_app_comm.
apply Permutation_sym. repeat rewrite List.concat_map.
repeat rewrite map_map. repeat rewrite substM_concat_map.
intros x; rewrite no_map_make_mono. apply NoDup_MonoSorted;
apply substM_is_poly.
intros m Hin. apply (substM_is_poly s x); auto.
intros m Hin. apply in_app_iff in Hin as ||; destruct Hp; destruct Hq; auto.
intros m Hin. apply in_app_iff in Hin as ||; apply nodup_cancel_in in H;
apply mono_in_map_make_mono in H; auto.
Qed.

The next six lemmas deal with proving that substitution distributes over polynomial multiplication. Given a substitution $s$ and two terms in polynomial form $p$ and $q$, it is shown that $s(p * q) \downarrow_p = (s(p) * s(q)) \downarrow_p$. The proof turns out to be much more difficult than the one for addition because the underlying arithmetic operation is more complex.

If two monomials are permutations (obviously not in monomial form), then applying any substitution to either will produce the same result. A weaker form that follows from this is that the results are permutations as well.

**Lemma substM_Permutation_eq :** \(\forall s m n,\)

**Permutation** \(m n \rightarrow\)
\text{substM} \; s \; m = \text{substM} \; s \; n.

\textbf{Proof.}

\begin{itemize}
  \item \text{intros} \; s \; m \; n \; H. \; \text{induction} \; H; \; \text{auto}.
  \item - \text{simpl.} \; \text{rewrite} \; \text{IHP}{}_{\text{Permutation}}. \; \text{auto}.
  \item - \text{simpl.} \; \text{rewrite} \; \text{mulPP\_comm.} \; \text{rewrite} \; \text{mulPP\_assoc.}
    \text{rewrite} \; (\text{mulPP\_comm} \; (\text{substM} \; s \; l)). \; \text{auto}.
  \item - \text{rewrite} \; \text{IHP}{}_{\text{Permutation1}}. \; \text{rewrite} \; \text{IHP}{}_{\text{Permutation2}}. \; \text{auto}.
\end{itemize}

\text{Qed.}

\textbf{Lemma substM\_Permutation :} \; \forall \; s \; m \; n,
\begin{itemize}
  \item \text{Permutation} \; m \; n \rightarrow
  \item \text{Permutation} \; (\text{substM} \; s \; m) \; (\text{substM} \; s \; n).
\end{itemize}

\textbf{Proof.}

\begin{itemize}
  \item \text{intros} \; s \; m \; n \; H. \; \text{rewrite} \; (\text{substM\_Permutation\_eq} \; s \; m \; n); \; \text{auto}.
\end{itemize}

\text{Qed.}

Adding duplicate variables to a monomial doesn’t change the result of applying a substitution. This is only true if the substitution’s range only has polynomials.

\textbf{Lemma substM\_nodup\_pointless :} \; \forall \; s \; m,
\begin{itemize}
  \item \text{is\_poly\_subst} \; s \rightarrow
  \item \text{substM} \; s \; (\text{nodup\_var\_eq\_dec} \; m) = \text{substM} \; s \; m.
\end{itemize}

\textbf{Proof.}

\begin{itemize}
  \item \text{intros} \; s \; m \; Hps. \; \text{induction} \; m; \; \text{auto}. \; \text{simpl.} \; \text{destruct in\_dec}.
  \item - \text{apply} \; \text{in\_split} \; \text{in} \; i. \; \text{destruct} \; i \; \text{as} \; [l1 \; l2 \; H]].
    \item assert \; (\text{Permutation} \; m \; (a \; :: \; l1 \; ++ \; l2)). \; \text{rewrite} \; H. \; \text{apply} \; \text{Permutation\_sym}.
  \item - \text{apply} \; \text{Permutation\_middle}.
  \item - \text{apply} \; \text{substM\_Permutation\_eq} \; \text{with} \; (s:=s) \; \text{in} \; H0. \; \text{rewrite} \; H0. \; \text{simpl}.
  \item - \text{rewrite} \; (\text{mulPP\_comm} \; (\text{substM} \; s \; l)). \; \text{rewrite} \; \text{mulPP\_comm}.
  \item - \text{rewrite} \; \text{mulPP\_assoc.} \; \text{rewrite} \; \text{mulPP\_p\_p.} \; \text{rewrite} \; \text{mulPP\_comm.} \; \text{rewrite} \; \text{IHM}.
  \item - \text{rewrite} \; H0. \; \text{simpl.} \; \text{auto}. \; \text{apply} \; \text{substV\_is\_poly.} \; \text{auto}.
  \item - \text{simpl.} \; \text{rewrite} \; \text{IHM}. \; \text{auto}.
\end{itemize}

\text{Qed.}

The idea behind the following two lemmas is that substitutions distribute over multiplication of a monomial and polynomial. The specifics of both are convoluted, yet easier to prove than distribution over two polynomials.

\textbf{Lemma substM\_distr\_mulMP :} \; \forall \; m \; n \; s,
\begin{itemize}
  \item \text{is\_poly\_subst} \; s \rightarrow
  \item \text{is\_mono} \; n \rightarrow
  \item \text{Permutation}
    \begin{itemize}
      \item \text{(nodup\_cancel\_mono\_eq\_dec} \; (\text{map\_make\_mono} \; (\text{substM} \; s \; (\text{make\_mono} \; (\text{make\_mono} \; (m \; ++ \; n)))))))
      \item \text{(nodup\_cancel\_mono\_eq\_dec} \; (\text{map\_make\_mono} \; (\text{concat} \; (\text{map} \; (\text{mulMP}'') \; (\text{map\_make\_mono} \; (\text{substM} \; s \; m))) \; (\text{map\_make\_mono} \; (\text{substM} \; s \; n))))))
    \end{itemize}
\end{itemize}
Proof.

intros m n s Hps H. rewrite (no_make_mono (make_mono (m ++ n))); auto.
repeat rewrite (no_map_make_mono (substM s _)); auto. apply Permutation_trans
with (l':= (nodup_cancel mono_eq_dec (substM s (nodup var_eq_dec
(m ++ n))))). apply nodup_cancel_Permutation. apply substM_Permutation.
unfold make_mono. rewrite ← Permutation_VarSort_l. auto.

induction m.

- simpl. pose (mulPP_1r (substM s n)). rewrite mulPP_comm in e.
  pose (substM_is_poly s n). apply e in i. rewrite mulPP_mulPP'' in i.
  unfold mulPP'' in i. rewrite ← no_make_poly in i; auto.
  apply Permutation_sort_eq in i. rewrite i. rewrite no_nodup_NoDup.
  rewrite no_map_make_mono. auto. intros m Hin. apply (substM_is_poly s n);
  auto. apply NoDup_VarSorted. auto.
- simpl substM at 2. apply Permutation_sort_eq. rewrite make_poly_refold.
  rewrite mulPP''''_refold. rewrite ← mulPP_mulPP'''' rewrite mulPP_assoc.
  repeat rewrite mulPP'''' rewrite mulPP'''' rewrite mulPP'''' rewrite mulPP''''
  rewrite substM_nodup_pointless; auto. simpl. rewrite mulPP_mulPP''''
  unfold mulPP'''' at 1. apply Permutation_sort_eq in IHm.
  rewrite make_poly_refold in IHm. rewrite mulPP''''_refold in IHm.
  rewrite no_nodup_cancel_NoDup in IHm. rewrite no_sortMonoSorted in IHm.
  rewrite ← substM_nodup_pointless; auto. rewrite IHm. unfold make_poly.
  apply Permutation_trans with (l':= (nodup_cancel mono_eq_dec (nodup_cancel
  mono_eq_dec (map make_mono (concat (map (mulPP'' substV s a))
  (mulPP'' substM s m) (substM s n)))))).
  apply nodup_cancel_Permutation. rewrite ← Permutation_MonoSort_l. auto.
  rewrite no_nodup_cancel_NoDup; auto.
  apply NoDup_nodup_cancel. apply substM_is_poly. apply NoDup_MonoSorted.
  apply substM_is_poly.
- intros m0 Hin. apply (substM_is_poly s n). auto.
- intros m0 Hin. apply (substM_is_poly s m). auto.
- intros m0 Hin. apply (substM_is_poly s (make_mono (m ++ n))). auto.
Qed.

Lemma map_substM_distr_map_mulMP : ∀ m p s,
  is_poly subst s →
  is_poly p →
  Permutation
  (nodup_cancel mono_eq_dec (map make_mono (concat (map (substM s) (map
  make_mono (mulPP'' p m)))))))
  (nodup_cancel mono_eq_dec (map make_mono (concat (map (mulPP'' (map
  make_mono (concat (map (substM s) p)))) (map make_mono (substM s m))))))

Proof.
intros m p s Hps H. unfold mulPP'' at 1. apply Permutation_trans with (l':=
Here is the formulation of substitution distributing over polynomial multiplication. Similar to the proof for addition, it is very dense and makes common use of permutation facts. Where it differs from that proof is that it relies on the commutativity of multiplication. The proof of distribution over addition didn’t need any properties of addition.

Lemma substP_distr_mulPP : ∀ p q s, substP s (mulPP p q) = mulPP (substP s p) (substP s q).

Proof.
intros p q s Hps H. repeat rewrite mulPP_mulPP''. unfold substP, mulPP''.
apply Permutation_sort_eq. apply Permutation_trans with (l' := (nodup_cancel mono_eq_dec (map make_mono (concat (map make_mono (concat (map make_mono (map (mulMP'' (map make_mono (concat (map make_mono (map (mulMP'' p) q))))))))))))).
apply nodup-cancel_Permutation. apply Permutation_map.
apply Permutation_concat. apply Permutation_map. unfold make_poly.
rewrite ← Permutation_MonoSort_l. auto.
apply Permutation_sym. apply Permutation_trans with \( \ell' := \text{nodup-cancel mono_eq_dec} (\text{map make_mono} (\text{concat} (\text{map} (\text{mulMP''}\ (\text{make_poly} (\text{concat} (\text{map} (\text{substM}\ s)\ p)))))) (\text{nodup-cancel mono_eq_dec} (\text{map make_mono} (\text{concat} (\text{map} (\text{substM}\ s)\ q)))))).\) apply nodup-cancel_Permutation.
apply Permutation_map. apply Permutation_concat. apply Permutation_map.
unfold make_poly. rewrite ← Permutation_MonoSort_l. auto.
apply Permutation_trans with \( \ell' := \text{nodup-cancel mono_eq_dec} (\text{map make_mono} (\text{concat} (\text{map} (\text{mulMP''}\ (\text{make_poly} (\text{concat} (\text{map} (\text{substM}\ s)\ p)))))) (\text{nodup-cancel mono_eq_dec} (\text{map make_mono} (\text{concat} (\text{map} (\text{substM}\ s)\ q)))))).\) apply nodup-cancel_concat_map. intros x. rewrite no_map_make_mono.
unfold mulMP". apply NoDup_MonoSorted. apply make_poly_is_poly.
intros m Hin. apply mono_in_make_poly in Hin; auto.
apply Permutation_sort_eq. rewrite make_poly_refold. rewrite mulPP''''_refold.
rewrite ← mulPP-mulPP"'''. rewrite mulPP_mulPP"'''.
apply Permutation_sort_eq. apply Permutation_trans with \( \ell' := \text{nodup-cancel mono_eq_dec} (\text{map make_mono} (\text{concat} (\text{map} (\text{mulMP''}\ (\text{make_poly} (\text{concat} (\text{map} (\text{substM}\ s)\ p)))))) (\text{nodup-cancel mono_eq_dec} (\text{map make_mono} (\text{concat} (\text{map} (\text{substM}\ s)\ q)))))).\) apply nodup-cancel_concat_map. intros x. rewrite no_map_make_mono.
unfold mulMP". apply NoDup_MonoSorted. apply make_poly_is_poly.
intros m Hin. apply mono_in_make_poly in Hin; auto.
apply Permutation_sort_eq. rewrite make_poly_refold. rewrite mulPP''''_refold.
rewrite ← mulPP-mulPP"'''. rewrite mulPP_mulPP"'''.
apply Permutation_sort_eq. apply Permutation_trans with \( \ell' := \text{nodup-cancel mono_eq_dec} (\text{map make_mono} (\text{concat} (\text{map} (\text{substM}\ s)\ (\text{map make_mono} (\text{concat} (\text{map} (\text{mulMP''}\ p)\ q))))))).\) repeat rewrite (List.concat_map make_mono (map (mulMP" _))). repeat rewrite (map_map _ (map make_mono)).
apply nodup-cancel_concat_map. intros x. rewrite no_map_make_mono. unfold mulMP". apply NoDup_MonoSorted; apply substM_is_poly.
intros m Hin; apply (substM_is_poly s x); auto.
induction q; auto. simpl. repeat rewrite map_app. repeat rewrite concat_app.
repeat rewrite map_app. repeat rewrite ← (nodup_cancel_pointless (map _ _)).
repeat rewrite ← (nodup_cancel_pointless_r_ (map _ _)).
apply nodup_cancel_Permutation. apply Permutation_app.
apply map_substM_distr_map_mulMP; auto. apply IHq.
Qed.

7.4 Unifiable Definitions

The following six definitions are all predicate functions that verify some property about substitutions or polynomials.

A unifier for a given polynomial \( p \) is a substitution \( s \) such that \( s(p) \downarrow_p = 0 \). This definition also includes that the range of the substitution only contain terms in polynomial form.

**Definition unifier** \( (s : \text{subst}) \ (p : \text{poly}) : \text{Prop} := \)

\[ \text{is\_poly\_subst} \ s \land \text{substP} \ s \ p = []. \]

A polynomial \( p \) is unifiable if there exists a unifier for \( p \).

**Definition unifiable** \( (p : \text{poly}) : \text{Prop} := \)

\[ \exists \ s, \ \text{unifier} \ s \ p. \]

A substitution \( u \) is a composition of two substitutions \( s \) and \( t \) if \( u(x) \downarrow_p = t(s(x)) \downarrow_p \) for every variable \( x \). The lemma \text{subst\_comp\_eq\_poly} below extends this definition from variables to polynomials.

**Definition subst\_comp\_eq** \( (s t u : \text{subst}) : \text{Prop} := \)

\[ \forall \ x, \ \text{substP} \ t (\text{substP} \ s \ [[x]]) = \text{substP} \ u \ [[x]]. \]

A substitution \( s \) is more general than a substitution \( t \) if there exists a third substitution \( u \) such that \( t \) is a composition of \( u \) and \( s \).

**Definition more\_general** \( (s t : \text{subst}) : \text{Prop} := \)

\[ \exists \ u, \ \text{subst\_comp\_eq} \ s \ u \ t. \]

Given a polynomial \( p \), a substitution \( s \) is the most general unifier of \( p \) if \( s \) is more general than every unifier of \( p \).

**Definition mgu** \( (s : \text{subst}) \ (p : \text{poly}) : \text{Prop} := \)

\[ \text{unifier} \ s \ p \land \forall \ t, \ \text{unifier} \ t \ p \rightarrow \ \text{more\_general} \ s \ t. \]

Given a polynomial \( p \), a substitution \( s \) is a reproductive unifier of \( p \) if \( t \) is a composition of itself and \( s \) for every unifier \( t \) of \( p \). This property is similar but stronger than most general because the substitution that composes with \( s \) is restricted to \( t \), whereas in most general it can be any substitution.

**Definition reprod\_unif** \( (s : \text{subst}) \ (p : \text{poly}) : \text{Prop} := \)

\[ \text{unifier} \ s \ p \land \]
∀ t, unifier t p → subst_comp_eq s t t.

Because the notion of most general is weaker than reproductive, it can be proven to logically follow as shown below. Any unifier that is reproductive is also most general.

Lemma reprod_is_mgu : ∀ p s, reprod_unif s p → mgu s p.

Proof.
unfold mgu, reprod_unif, more_general, subst_comp_eq.
intros p s [].
split; auto.
intros.
exists t.
intros.
apply H0; auto.
Qed.

As stated earlier, substitution composition can be extended to polynomials. This comes from the implicit fact that if two substitutions agree on all variables then they agree on all terms.

Lemma subst_comp_eq_poly : ∀ s t u, is_poly_subst s → is_poly_subst t → is_poly_subst u → (∀ x, substP t (substP s [[x]]) = substP u [[x]]) → ∀ p, substP t (substP s p) = substP u p.

Proof.
intros. induction p; auto. simpl. unfold substP at 2. simpl.
rewrite ← make_poly_pointless_r. rewrite addPP_refold.
rewrite substP_distr_addPP; auto. unfold substP at 3. simpl.
rewrite ← make_poly_pointless_r. rewrite addPP_refold. f_equal.
- induction a; auto. simpl. rewrite substP_distr_mulPP; auto. f_equal; auto.
  + rewrite ← subst_var_eq; auto. rewrite ← subst_var_eq; auto.
  + apply substV_is_poly; auto.
- rewrite substP_refold. apply IHp.
Qed.

The last lemmas of this section state that the identity substitution is a reproductive unifier of the constant zero. Therefore it is also most general.

Lemma empty_unifier : unifier [] []

Proof.
unfold unifier, is_poly_subst. split; auto.
intros. inversion \( H \).
Qed.

Lemma empty_reprod_unif : reprod_unif [] [].
Proof.
  unfold reprod_unif, more_general, subst_comp_eq.
  split; auto. apply empty_unifier.
Qed.

Lemma empty_mgu : mgu [] [].
Proof.
  apply reprod_is_mgu. apply empty_reprod_unif.
Qed.
Chapter 8

Library B_Unification.sve

Require Import List.
Import ListNotations.
Require Import Arith.
Require Import Permutation.
Require Export poly_unif.

8.1 Introduction

Here we implement the algorithm for successive variable elimination. The basic idea is to remove a variable from the problem, solve that simpler problem, and build a solution from the simpler solution. The algorithm is recursive, so variables are removed and problems are generated until we are left with either of two problems; \( 1 \approx_B 0 \) or \( 0 \approx_B 0 \). In the former case, the whole original problem is not unifiable. In the latter case, the problem is solved without any need to substitute since there are no variables. From here, we begin the process of building up substitutions until we reach the original problem.

8.2 Eliminating Variables

This section deals with the problem of removing a variable \( x \) from a term \( t \). The first thing to notice is that \( t \) can be written in polynomial form \( t \downarrow_p \). This polynomial is just a set of monomials, and each monomial a set of variables. We can now separate the polynomials into two sets \( qx \) and \( r \). The term \( qx \) will be the set of monomials in \( t \downarrow_p \) that contain the variable \( x \). The term \( q \), or the quotient, is \( qx \) with the \( x \) removed from each monomial. The term \( r \), or the remainder, will be the monomials in \( t \downarrow_p \) that do not contain \( x \). The original term can then be written as \( x \ast q + r \).

Implementing this procedure is pretty straightforward. We define a function \( \text{div}\_by\_var \) that produces two polynomials given a polynomial \( p \) and a variable \( x \) to eliminate from it.
The first step is dividing \( p \) into \( qx \) and \( r \) which is performed using a partition over \( p \) with the predicate \texttt{has_var}. The second step is to remove \( x \) from \( qx \) using the helper \texttt{elim_var}.

The function \texttt{has_var} determines whether a variable appears in a monomial.

**Definition** \texttt{has_var} \((x : \text{var}) := \exists b \ (\texttt{beq_nat} x)\).

The function \texttt{elim_var} removes a variable from each monomial in a polynomial. It is possible that this leaves the term not in polynomial form so it is then repaired with \texttt{make_poly}.

**Definition** \texttt{elim_var} \((x : \text{var}) (p : \text{poly}) : \text{poly} := \text{make_poly} (\text{map} (\text{remove var_eq_dec} x) p)\).

The function \texttt{div_by_var} produces a quotient \( q \) and remainder \( r \) from a polynomial \( p \) and variable \( x \) such that \( p \approx_B x \ast q + r \) and \( x \) does not occur in \( r \).

**Definition** \texttt{div_by_var} \((x : \text{var}) (p : \text{poly}) : \text{prod} \text{ poly poly} := \langle \text{let} (qx, r) := \text{partition} (\texttt{has_var} x) p \text{ in} \ (\texttt{elim_var} x qx, r) \rangle\).

We would also like to prove some lemmas about variable elimination that will be helpful in proving the full algorithm correct later. The main lemma below is \texttt{div_eq}, which just asserts that after eliminating \( x \) from \( p \) into \( q \) and \( r \) the term can be put back together as in \( p \approx_B x \ast q + r \). This fact turns out to be rather hard to prove and needs the help of 10 or so subsidiary lemmas.

After eliminating a variable \( x \) from a polynomial to produce \( r \), \( x \) does not occur in \( r \).

**Lemma** \texttt{elim_var_not_in_rem} : \( \forall \ x \ p \ r, \ \texttt{elim_var} x \ p = r \rightarrow (\forall m, \text{ln} m \ r \rightarrow \neg \text{ln} x m) \).

**Proof**.

\begin{verbatim}
intros.
unfold \texttt{elim_var} in \( H \).
unfold \texttt{make_poly} in \( H \).
rewrite \( \leftarrow H \) in \( H0 \).
apply \texttt{ln_sorted} in \( H0 \).
apply \texttt{nodup_cancel_in} in \( H0 \).
rewrite \texttt{map_map} in \( H0 \).
apply \texttt{in_map_iff} in \( H0 \) as \[ n \ [ ] \].
rewrite \( \leftarrow H0 \).
intro.
rewrite \texttt{make_mono_in} in \( H2 \).
apply \texttt{remove_in} in \( H2 \).
auto.
Qed.
\end{verbatim}

Eliminating a variable from a polynomial produces a term in polynomial form.

**Lemma** \texttt{elim_var_is_poly} : \( \forall \ x \ p, \ \texttt{is_poly} (\texttt{elim_var} x p) \).
Proof.

intros.
unfold elim_var.
apply make_poly_is_poly.
Qed.

Hint Resolve elim_var_is_poly.

The next four lemmas deal with the following scenario: Let \( p \) be a term in polynomial form, \( x \) be a variable that occurs in each monomial of \( p \), and \( r = \text{elim}_\text{var} \, x \, p \).

The term \( r \) is a permutation of removing \( x \) from \( p \). Another way of looking at this statement is when \( \text{elim}_\text{var} \) repairs the term produced from removing a variable it only sorts that term.

Lemma elim_var_map_remove_Permutation : \( \forall \, p \, x, \)
\( \text{is}_\text{poly} \, p \rightarrow \)
\( (\forall \, m, \text{In} \, m \, p \rightarrow \text{In} \, x \, m) \rightarrow \)
\( \text{Permutation} \,(\text{elim}_\text{var} \, x \, p) \,(\text{map} \,(\text{remove} \, \text{var_eq_dec} \, x) \, p) \).  
Proof.

intros \( p \, x \, H \, H0 \). destruct \( p \) as \( \| a \, p \| \).
- simpl. unfold elim_var, make_poly, MonoSort.sort. auto.
- simpl. unfold elim_var. simpl. unfold make_poly.
  rewrite ← Permutation_MonoSort_l. rewrite unsorted_poly; auto.
  + rewrite ← map_cons. apply NoDup_map_remove; auto.
  + apply poly_cons in \( H \). intros \( m \, Hin \). destruct \( Hin \).
    × rewrite ← \( H1 \). apply remove_is_mono. apply \( H \).
    × apply in_map_iff in \( H1 \) as \( y \| \| \). rewrite ← \( H1 \). apply remove_is_mono.
    destruct \( H \). unfold is_poly in \( H \). destruct \( H \). apply \( H4 \). auto.

Qed.

The term \((x \ast r) \downarrow_P \) is a permutation of the result of removing \( x \) from \( p \), appending \( x \) to the end of each monomial, and repairing each monomial. The proof relies on the mulPP_map_app_permutation lemma from the poly library, which has a simpler goal but does much of the heavy lifting.

Lemma rebuild_map_permutation : \( \forall \, p \, x, \)
\( \text{is}_\text{poly} \, p \rightarrow \)
\( (\forall \, m, \text{In} \, m \, p \rightarrow \text{In} \, x \, m) \rightarrow \)
\( \text{Permutation} \,(\text{mulPP} \,(\,[x]\,) \,(\text{elim}_\text{var} \, x \, p))\)
\( \,(\text{map} \,(\text{fun} \, a \Rightarrow \text{make}_\text{mono} \,(a \,\text{++} \,[x])))\)
\( \,(\text{map} \,(\text{remove} \, \text{var_eq_dec} \, x) \, p)) \).

Proof.

intros \( p \, x \, H \, H0 \). apply mulPP_map_app_permutation; auto.
- apply (elim_var_not_in_rem x p); auto.
- apply elim_var_map_remove_Permutation; auto.

Qed.
The term $p$ is a permutation of $(x * r) \downarrow_p$. Proof of this fact relies on the lengthy \text{map\_app\_remove\_Permutation} lemma from \text{poly}.

\textbf{Lemma elim\_var\_permutation}: $\forall p x$, is\_poly $p \rightarrow$ $(\forall m, \ln m p \rightarrow \ln x m) \rightarrow$ $\text{Permutation} \ p (\text{mulPP} \ [[x]] \ (\text{elim\_var} \ x \ p))$.

Proof.

intros $p x H H0$. pose (rebuild\_map\_permutation $p x H H0$).
apply Permutation\_sym in $p0$.
pose (map\_app\_remove\_Permutation $p x H H0$).
apply (Permutation\_trans $p1 \ p0$).
Qed.

Finally, $p = (x * r) \downarrow_p$.

\textbf{Lemma elim\_var\_mul}: $\forall x p$, is\_poly $p \rightarrow$ $(\forall m, \ln m p \rightarrow \ln x m) \rightarrow$ $p = \text{mulPP} \ [[x]] \ (\text{elim\_var} \ x \ p)$.

Proof.

intros. apply Permutation\_Sorted\_eq.
- apply elim\_var\_permutation; auto.
- unfold is\_poly in $H$. apply Sorted\_MonoSorted. apply $H$.
- pose (mulPP\_is\_poly $[[x]] \ (\text{elim\_var} \ x \ p))$. unfold is\_poly in $i$.
  - apply Sorted\_MonoSorted. apply $i$.
Qed.

The function \text{has\_var} is an equivalent boolean version of the $\ln$ predicate.

\textbf{Lemma has\_var\_eq\_in}: $\forall x m$, \text{has\_var} $x m = \text{true} \leftrightarrow \ln x m$.

Proof.

intros.
unfold has\_var.
rewrite existsb\_exists.
split; intros.
- destruct $H$ as $[x0 \ ||]$.
  - apply Nat\_eqb\_eq in $H0$.
    - rewrite $H0$. apply $H$.
  - $\exists x$. rewrite Nat\_eqb\_eq. auto.
Qed.

Let a polynomial $p$ be partitioned by \text{has\_var} $x$ into two sets $qx$ and $r$. Obviously, every monomial in $qx$ contains $x$ and no monomial in $r$ contains $x$.

\textbf{Lemma part\_var\_eq\_in}: $\forall x p qx r$, partition (has\_var $x$) $p = (qx, r) \rightarrow$
\((\forall m, \ln m \cdot qx \rightarrow \ln x \cdot m) \land \\
(\forall m, \ln m \cdot r \rightarrow \neg \ln x \cdot m)\).

Proof.

intros.
split; intros.
- apply part_fst_true with \((a:=m)\) in \(H\).
  + apply has_var_eq_in. apply \(H\).
  + apply \(H_0\).
- apply part_snd_false with \((a:=m)\) in \(H\).
  + rewrite \leftrightarrow has_var_eq_in. rewrite \(H\). auto.
  + apply \(H_0\).
Qed.

The function \(\text{div by var}\) produces two terms both in polynomial form.

Lemma \(\text{div is poly}\) :
\[
\forall x p q r, \\
\text{is poly } p \rightarrow \\
\text{div by var } x p = (q, r) \\
\text{is poly } q \land \text{is poly } r.
\]

Proof.

intros.
unfold \(\text{div by var}\) in \(H_0\).
destruct \((\text{partition (has var } x) p)\) eqn:Hp.
apply \((\text{part is poly } \_ \_ \_ H)\) in Hpart as Hp.
destruct \(Hp\) as [Hpl Hpr].
injection \(H_0\). intros Hr Hq.
rewrite \(Hr\) in Hpr.
apply part_var_eq_in in Hpart as [Hin Hout].
split.
- rewrite \leftrightarrow Hq; auto.
- apply Hpr.
Qed.

As explained earlier, given a polynomial \(p\) decomposed into a variable \(x\), a quotient \(q\), and a remainder \(r\), \(\text{div eq}\) asserts that \(p = (x \cdot q + r) \downarrow p\).

Lemma \(\text{div eq}\) :
\[
\forall x p q r, \\
\text{is poly } p \rightarrow \\
\text{div by var } x p = (q, r) \rightarrow \\
p = \text{addPP (mulPP [[x]] } q) r.
\]

Proof.

intros \(x p q r HP HD\).
assert \((HE := HD)\).
unfold \(\text{div by var}\) in \(HE\).
destruct \((\text{partition (has var } x))\) as \([qx r0]\) eqn:Hq.
injection HE. intros Hr Hq.
assert (HIH: ∀ m, ln m qx → ln x m). intros.
apply has_var_eq_in.
apply (part_fst_true _ _ _ _ _ _ Hqr _ _ H).
assert (is_poly q ∧ is_poly r) as [HPq HPr].
apply (div_is_poly _ _ _ _ _ _ HP HD).
assert (is_poly qx ∧ is_poly r0) as [HPqx HPr0].
rewrite ← Hq.
rewrite ← (elim_var_mul x qx HPqx HIH).
apply (part_add_eq (has_var x) _ _ _ HP).
rewrite ← Hr.
apply Hqr.
Qed.

Given a variable x, div_by_var produces two polynomials neither of which contain x.

Lemma div_var_not_in_qr : ∀ x p q r,
div_by_var x p = (q, r) →
((∀ m, ln m q → ¬ ln x m) ∧
(∀ m, ln m r → ¬ ln x m)).
Proof.
    intros.
    unfold div_by_var in H.
    assert (∃ qxr, qxr = partition (has_var x) p) as [[q r0] Hqxr]. eauto.
    rewrite ← Hqxr in H.
    injection H. intros Hr Hq.
    split.
    - apply (elim_var_not_in_rem _ _ _ Hq).
    - rewrite Hr in Hqxr.
        symmetry in Hqxr.
        intros. intro.
        apply has_var_eq_in in H1.
        apply Bool.negb_false_iff in H1.
        revert H1.
        apply Bool.eq_true_false_abs.
        apply Bool.negb_true_iff.
        revert m H0.
        apply (part_snd_false _ _ _ _ Hqxr).
    Qed.

This helper function build_poly is used to construct p' = ((q + 1) * r) ↓_p given the two polynomials q and r as input.

Definition build_poly (q r : poly) : poly :=
mulPP (addPP [] q) r.

The function build_poly produces a term in polynomial form.

Lemma build_poly_is_poly : \forall q r,
is_poly (build_poly q r).

Proof.
unfold build_poly. auto.
Qed.

Hint Resolve build_poly_is_poly.

The second main lemma about variable elimination is below. Given that a term p has been decomposed into the form \( (x \cdot q + r) \downarrow_p \), we can define \( p' = ((q + 1) \cdot r) \downarrow_p \). The lemma div_build_unif states that any unifier of \( p \approx_B 0 \) is also a unifier of \( p' \approx_B 0 \). Much of this proof relies on the axioms of polynomial arithmetic.

Lemma div_build_unif : \forall x p q r s,
is_poly p \rightarrow
div_by_var x p = (q, r) \rightarrow
unifier s p \rightarrow
unifier s (build_poly q r).

Proof.
unfold build_poly, unifier.
intros x p q r s HPp HD [Hps Hsp0].
apply (div_eq _ _ _ HPp) in HD as Hp.
assert (\exists q1, q1 = addPP [] q) as [q1 Hq1]. eauto.
assert (\exists sp, sp = substP s p) as [sp Hsp]. eauto.
assert (\exists sq1, sq1 = substP s q1) as [sq1 Hsq1]. eauto.
rewrite ← (mulPP_0 (substP s q1)).
rewrite ← Hsp0.
rewrite Hp, Hq1.
rewrite ← substP_distr_mulPP; auto.
f_equal.
apply (div_is_poly x p q r HPp) in HD.
destruct HD as [HPq HPr].
rewrite mulPP_addPP_1; auto.
Qed.

Given a polynomial p and a variable x, div_by_var produces two polynomials q and r that have no more variables than p has. Obviously, q and r don’t contain x either.

Lemma incl_div : \forall x p q r xs,
is_poly p \rightarrow
div_by_var x p = (q, r) \rightarrow
incl (vars p) (x :: xs) \rightarrow
incl (vars q) xs \land incl (vars r) xs.
Proof.

intros. assert (Hdiv:=H0). unfold div_by_var in H0.
destruct partition as [qx r0] eqn:Hpart. apply partition_Permutation in Hpart.
apply Permutation_incl in Hpart as []. inversion H0. clear H2.
assert (incl (vars q) (vars p)). unfold incl, vars in *. intros a Hin.
apply nodup_ln. apply nodup_ln in Hin. apply ln_concat_exists in Hin.
destruct Hin as [m ||]. rewrite ← H5 in H2. unfold elim_var in H2.
apply ln_sorted in H2. apply nodup_cancel_in in H2. rewrite map_map in H2.
apply in_map_iff in H2. destruct H2 as [mx ||]. rewrite ← H2 in H4.
rewrite make_mono_ln in H4. apply ln_remove in H4. apply ln_concat_exists.
exists mx. split; auto. apply H3. intuition.
assert (incl (vars r) (vars p)). rewrite H6 in H3. unfold incl, vars in *.
intros a Hin. apply nodup_ln. apply nodup_ln in Hin.
apply ln_concat_exists in Hin. destruct Hin as [l ||].
apply ln_concat_exists. ∃ l. split; auto. apply H3. intuition.
split.
- rewrite H5. apply incl_tran with (n:=(x::xs)) in H2; auto.
  apply incl_not_in in H2; auto. apply div_var_not_in_qr in Hdiv as [Hq _].
  apply in_mono_in_vars in Hq. auto.
- apply incl_tran with (n:=(x::xs)) in H4; auto.
  apply incl_not_in in H4; auto. apply div_var_not_in_qr in Hdiv as [_. Hr].
  apply in_mono_in_vars in Hr. auto.
Qed.

Given a term $p$ decomposed into the form $(x \ast q + r) \downarrow P$, then the polynomial $p' = ((q + 1) \ast r) \downarrow P$ has no more variables than $p$ and does not contain $x$.

Lemma div_vars : \( \forall x xs p q r, \)
is_poly p \rightarrow
incl (vars p) (x :: xs) \rightarrow
div_by_var x p = (q, r) \rightarrow
incl (vars (build_poly q r)) xs.
Proof.

intros x xs p q r Hincl Hdiv. unfold build_poly.
apply div_var_not_in_qr in Hdiv as Hin. destruct Hin as [Hinq Hinr].
apply in_mono_in_vars in Hinq. apply in_mono_in_vars in Hinr.
apply incl_vars_mulPP. apply (incl_div _ _ _ _ H Hdiv) in Hincl. split.
- apply incl_vars_addPP; auto. apply div_is_poly in Hdiv as []; auto. split.
  + unfold vars. simpl. unfold incl. intros a [].
  + apply Hincl.
- apply Hincl.
Qed.

Hint Resolve div_vars.
8.3 Building Substitutions

This section handles how a solution is built from subproblem solutions. Given that term \( p \) decomposed into \((x * q + r) \downarrow_P\) and \( p' = ((q + 1) * r) \downarrow_P\), the lemma `reprod_build_subst` states that if some substitution \( \sigma \) is a reproductive unifier of \( p' \approx_B 0 \), then we can build a substitution \( \sigma' \) which is a reproductive unifier of \( p \approx_B 0 \). The way \( \sigma' \) is built from \( \sigma \) is defined in `build_subst`. Another replacement is added to \( \sigma \) of the form \( \{ x \mapsto (x * (\sigma(q) + 1) + \sigma(r)) \downarrow_P \} \) to construct \( \sigma' \).

```plaintext
Definition build_subst (s : subst) (x : var) (q r : poly) : subst :=
  let q1 := addPP [[]] q in
  let q1s := substP s q1 in
  let rs := substP s r in
  let xs := (x, addPP (mulPP [[x]] q1s) rs) in
  xs :: s.
```

The function `build_subst` produces a substitution whose range only contains polynomials.

Lemma `build_subst_is_poly` : \( \forall s x q r, \)
```
is_poly_subst s \rightarrow
is_poly_subst (build_subst s x q r).
```

Proof.
```
unfold build_subst.
unfold is_poly_subst.
intros.
destruct H0.
- inversion H0. auto.
- apply (H x0). auto.
Qed.
```

Given that term \( p \) decomposed into \((x * q + r) \downarrow_P\), \( p' = ((q + 1) * r) \downarrow_P\), and \( \sigma \) is a reproductive unifier of \( p' \approx_B 0 \), then the substitution \( \sigma' \) built from \( \sigma \) unifies \( p \approx_B 0 \).

Lemma `build_subst_is_unif` : \( \forall x p q r s, \)
```
is_poly p \rightarrow
div_by_var x p = (q, r) \rightarrow
reprod_unif s (build_poly q r) \rightarrow
unifier (build_subst s x q r) p.
```

Proof.
```
unfold reprod_unif, unifier.
intros x p q r s Hpoly Hdiv [[Hps Hunif] Hreprod].
assert (is_poly_subst (build_subst s x q r)).
  apply build_subst_is_poly; auto.
split; auto.
unfold build_poly in Hunif.
```
assert (Hnqr := Hdiv).
apply div_var_not_in_qr in Hnqr.
destruct Hnqr as [Hnq Hnr].
assert (HpolyQR := Hdiv).
apply div_is_poly in HpolyQR as [HpolyQ HpolyR]; auto.
apply div_eq in Hdiv; auto.
rewrite Hdiv.
rewrite substP_distr_addPP; auto.
rewrite substP_distr_mulPP; auto.
unfold build_subst.
rewrite (substP_cons _ _ Hnq).
rewrite (substP_cons _ _ Hnr).
assert (Hsx: (substP
((x,
  addPP
    (mulPP [[x]])
  (addPP s (addPP [[]] q)))))
((x])) = (addPP
  (mulPP [[x]])
  (addPP s (addPP [[]] q)))).
unfold substP. simpl.
rewrite ← beq_nat_refl.
rewrite no_make_poly; auto.
rewrite Hsx.
rewrite substP_distr_addPP; auto.
rewrite substP_1.
rewrite mulPP_distr_addPPr; auto.
rewrite mulPP_1r; auto.
rewrite mulPP_distr_addPP; auto.
rewrite mulPP_distr_addPP; auto.
rewrite mulPP_assoc.
rewrite mulPP_p_p; auto.
rewrite addPP_p_p; auto.
rewrite addPP_0; auto.
rewrite substP_distr_mulPP; auto.
rewrite substP_distr_addPP; auto.
rewrite (mulPP_1r r) at 2; auto.
rewrite mulPP_comm; auto.
rewrite (mulPP_comm r [[]]); auto.
rewrite ← mulPP_distr_addPP; auto.
rewrite addPP_comm; auto.
Qed.

Given that term \( p \) decomposed into \( (x \cdot q + r) \downarrow_P, \) \( p' = ((q + 1) \cdot r) \downarrow_P, \) and \( \sigma \) is a reproductive unifier of \( p' \approx_B 0 \), then the substitution \( \sigma' \) built from \( \sigma \) is reproductive with regards to unifiers of \( p \approx_B 0 \).

Lemma build_subst_is_reprod : \( \forall x p q r s, \)
    is_poly p →
    div_by_var x p = (q, r) →
    reproducunif s (build_poly q r) →
    \( \forall t, \) unifier t p →
    substcomp_eq (build_subst s x q r) t t.
Proof.
  unfold reproducunif.
  intros x p q r s HpolyP Hdiv [[HpsS HunifS] Hsub_comp] t HunifT.
  assert (HunifT' := HunifT).
  destruct HunifT as [HpsT HunifT].
  apply (div_build_unif _ _ _ _ HpolyP Hdiv) in HunifT'.
  unfold substcomp_eq in *.
  intros y.
  destruct (y =? x) eqn:Hyx.
  - unfold build_subst.
    assert (H: (substP ((x, addPP (mulPP [[x]] (substP s (addPP [[]] q))))
        (substP s r)) :: s) [[y]]) =
        (addPP (mulPP [[x]] (substP s (addPP [[]] q)))) (substP s r)).
    unfold substP. simpl.
    rewrite Hyx.
    rewrite mulPP_lr; auto. rewrite app_nil_r.
    rewrite no_make_poly; auto.
    rewrite H.
    rewrite substP_distr_addPP; auto.
    rewrite substP_distr_mulPP; auto.
    pose (div_is_poly _ _ _ _ HpolyP Hdiv); destruct a.
    rewrite substP_distr_addPP; auto.
    rewrite substP_distr_addPP; auto.
    rewrite substP_s1.
    assert (Hdiv2 := Hdiv).
    apply div_eq in Hdiv; auto.
    apply div_is_poly in Hdiv2 as [HpolyQ HpolyR]; auto.
    rewrite (substcomp_eq_poly s t t); auto.
    rewrite (substcomp_eq_poly s t t); auto.
rewrite mulPP_comm; auto.
rewrite mulPP_distr_addPP; auto.
rewrite mulPP_comm; auto.
rewrite mulPP_1r; auto.
rewrite (addPP_comm (substP t [][]) _) auto.
rewrite addPP_assoc; auto.
rewrite (addPP_comm (substP t [][]) _) auto.
rewrite ← addPP_assoc; auto.
rewrite ← substP_distr_mulPP; auto.
rewrite ← substP_distr_addPP; auto.
rewrite mulPP_comm; auto.
rewrite ← Hdiv.
unfold unifier in HunifT.
rewrite HunifT.
rewrite addPP_0; auto.
apply beq_nat_true in Hyx.
rewrite Hyx.
reflexivity.
- unfold build_subst.
  rewrite substP_cons; auto.
  intros.
  inversion H; auto.
  rewrite ← H0.
  simpl. intro.
  destruct H1; auto.
  apply Nat.eqb_eq in H1.
  rewrite Hyx in H1.
  inversion H1.
Qed.

Given that term p decomposed into \((x \cdot q + r) \downarrow_P, p' = ((q + 1) \cdot r) \downarrow_P\), and a reproductive unifier \(\sigma\) of \(p' \approx_B 0\), then the substitution \(\sigma'\) built from \(\sigma\) is a reproductive unifier \(p \approx_B 0\) based on the previous two lemmas.

**Lemma reprod_build_subst** : \(\forall x\ p\ q\ r\ s,\)
- is_poly \(p \rightarrow\)
- div_by_var \(x\ p\ =\ (q, \ r) \rightarrow\)
- reprod_unif \(s\ (\text{build} - \text{poly} q\ r) \rightarrow\)
- reprod_unif \(\text{build} - \text{subst} s\ x\ q\ r\) \(p\).

**Proof.**
- intros. unfold reprod_unif. split.
  - apply build_subst_is_unif; auto.
  - apply build_subst_is_reprod; auto.
Qed.
8.4 Recursive Algorithm

Now we define the actual algorithm of successive variable elimination. Built using five helper functions, the definition is not too difficult to construct or understand. The general idea, as mentioned before, is to remove one variable at a time, creating simpler problems. Once the simplest problem has been reached, to which the solution is already known, every solution to each subproblem can be built from the solution to the successive subproblem. Formally, given the polynomials \( p = (xq + r) \downarrow_P \) and \( p' = ((q + 1)r) \downarrow_P \), the solution to \( p \approx_B 0 \) is built from the solution to \( p' \approx_B 0 \). If \( \sigma \) solves \( p' \approx_B 0 \), then \( \sigma \cup \{ x \mapsto (x \cdot (\sigma(q) + 1) + \sigma(r)) \downarrow_P \} \) solves \( p \approx_B 0 \).

The function \( \text{sve} \) is the final result, but it is \( \text{sveVars} \) which actually has all of the meat. Due to Coq’s rigid type system, every recursive function must be obviously terminating. This means that one of the arguments must decrease with each nested call. It turns out that Coq’s type checker is unable to deduce that continually building polynomials from the quotient and remainder of previous ones will eventually result in 0 or 1. So instead we add a fuel argument that explicitly decreases per recursive call. We use the set of variables in the polynomial for this purpose, since each subsequent call has at least one less variable.

Fixpoint \( \text{sveVars} \, (\text{varlist} : \text{list var}) \, (p : \text{poly}) : \text{option subst} \) :=

\[
\begin{align*}
\text{match varlist with} \\
\mid [] & \Rightarrow \\
\mid [] & \Rightarrow \text{Some []} \\
\mid & \Rightarrow \text{None} \\
\end{align*}
\]

\[
\begin{align*}
\mid x :: xs & \Rightarrow \\
\text{let} \ (q, r) := \text{div_by_var} \ x \ p \ \text{in} \\
\text{let} \ p' := (\text{build_poly} \ q \ r) \ \text{in} \\
\text{match} \ \text{sveVars} \ xs \ p' \ \text{with} \\
\mid & \Rightarrow \text{None} \\
\mid \text{Some} \ s & \Rightarrow \text{Some} (\text{build subst} \ s \ x \ q \ r) \\
\end{align*}
\]

end.

The function \( \text{sve} \) simply calls \( \text{sveVars} \) with an initial fuel of \( \text{vars} \ p \).

Definition \( \text{sve} \, (p : \text{poly}) : \text{option subst} := \text{sveVars} \, (\text{vars} \ p) \, p \).

8.5 Correctness

Finally, we must show that this algorithm is correct. As discussed in the beginning, the correctness of a unification algorithm is proven for two cases. If the algorithm produces a solution for a problem, then the solution must be most general. If the algorithm produces
no solution, then the problem must be not unifiable. These statements have been formalized in
the theorem sve_correct with the help of the predicates mgu and unifiable as defined in the
library poly_unif. The two cases of the proof are handled separately by the lemmas
sveVars_some and sveVars_none.

If sveVars produces a substitution σ, then the range of σ only contains polynomials.

Lemma sveVars_poly_subst : ∀ xs p,
    incl (vars p) xs →
    is_poly p →
    ∀ s, sveVars xs p = Some s →
        is_poly_subst s.

Proof.
  induction xs as [|| x xs]; intros.
  simpl in H1. destruct p; inversion H1. unfold is_poly_subst.
  intros x p ||.
  - intros.
    assert (exists qr, div_by_var x p = qr) as [[q r] Hqr]. eauto.
    simpl in H1.
    rewrite Hqr in H1.
    destruct (sveVars xs (build_poly q r)) eqn:Hs0; inversion H1.
    apply IHxs in Hs0; eauto.
    apply build_subst_is_poly; auto.
Qed.

If sveVars produces a substitution σ for the polynomial p, then σ is a most general unifier
of p ≈ B 0.

Lemma sveVars_some : ∀ (xs : list var) (p : poly),
    NoDup xs →
    incl (vars p) xs →
    is_poly p →
    ∀ s, sveVars xs p = Some s →
        mgu s p.

Proof.
  intros xs p Hdup H0 s H1.
  apply reprod_is_mgu.
  revert xs p Hdup H0 s H1.
  induction xs as || x xs.
  - intros. simpl in H1. destruct p; inversion H1.
    apply empty_reprod_unif.
  - intros.
    assert (exists qr, div_by_var x p = qr) as [[q r] Hqr]. eauto.
    simpl in H1.
    rewrite Hqr in H1.

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destruct (sveVars xs (build_poly q r)) eqn:Hs0; inversion H1.
apply NoDup_cons_iff in Hdup as Hnin. destruct Hnin as [Hnin Hdup0].
apply sveVars_poly_subst in Hs0 as HpsS0; eauto.
apply IHxs in Hs0; eauto.
apply reprod_build_subst; auto.
Qed.

If sveVars does not produce a substitution for the polynomial \( p \), then the problem \( p \equiv_B 0 \) is not unifiable.

Lemma sveVars_None : \( \forall (xs : list \text{var}) (p : \text{poly}), \)
\( \begin{align*}
&\text{NoDup} \; xs \rightarrow \\
&\text{incl} \; (\text{vars} \; p) \; xs \rightarrow \\
&\text{is_poly} \; p \rightarrow \\
&sveVars \; xs \; p = \text{None} \rightarrow \\
&\neg \text{unifiable} \; p.
\end{align*} \)
Proof.
induction \( xs \) as \( ||x\;xs\).\n- intros \( p \) Hdup H H0 H1. simpl in H1. destruct \( p \); inversion H1. intro.
  unfold unifiable in H2. destruct H2. unfold unifier in H2.
  apply incl_nil in H. apply no_vars_is_ground in H; auto.
  destruct H; inversion H.
  rewrite H4 in H2.
  rewrite H5 in H2.
  rewrite substP_1 in H2.
  inversion H2. inversion H6.
- intros \( p \) Hdup H H0 H1.
  assert (\( \exists qr, \text{div_by_var} \; x \; p = qr \)) as [\( \text{eq} \; Hqr \)]. eauto.
  simpl in H1.
  rewrite Hqr in H1.
  destruct (sveVars xs (build_poly q r)) eqn:Hs0; inversion H1.
  apply NoDup_cons_iff in Hdup as Hnin. destruct Hnin as [Hnin Hdup0].
  apply IHxs in Hs0; eauto.
  unfold not, unifiable in *.
  intros.
  apply Hs0.
  destruct H2 as [\( s \) \( Hu \)].
  \( \exists s. \)
  apply (div_build_unif \( x \; p \)); auto.
Qed.

Hint Resolve NoDup_vars incl_refl.

If sveVars produces a substitution \( \sigma \) for the polynomial \( p \), then \( \sigma \) is a most general unifier of \( p \equiv_B 0 \). Otherwise, \( p \equiv_B 0 \) is not unifiable.
Lemma sveVars_correct : \( \forall (p : \text{poly}), \)
\[
  \text{is\_poly } p \rightarrow \\
  \text{match} \text{sveVars} (\text{vars } p) \text{ } \text{p} \text{ with} \\
  | \text{Some } s \Rightarrow \text{mgu } s \text{ } p \\
  | \text{None } \Rightarrow \neg \text{unifiable } p \\
\]
end.
Proof.
intros.
destruct (sveVars (vars p) p) eqn: Hsve.
- apply (sveVars_some (vars p)); auto.
- apply (sveVars_none (vars p)); auto.
Qed.

If sve produces a substitution \( \sigma \) for the polynomial \( p \), then \( \sigma \) is a most general unifier of \( p \approx_B 0 \). Otherwise, \( p \approx_B 0 \) is not unifiable.

Theorem sve_correct : \( \forall (p : \text{poly}), \)
\[
  \text{is\_poly } p \rightarrow \\
  \text{match} \text{sve } p \text{ with} \\
  | \text{Some } s \Rightarrow \text{mgu } s \text{ } p \\
  | \text{None } \Rightarrow \neg \text{unifiable } p \\
\]
end.
Proof.
intros.
apply sveVars_correct.
auto.
Qed.


